

- Classically, simplicial/cubical sets model spaces/ ∞ -groupoids
- Simplicial sets more convenient
 - products respect homotopy types
 - Nerve functor from categories
- How "well" do $\hat{\Delta}, \hat{\square}, \hat{\Theta}, \hat{\Omega}, \hat{\Delta}_s$, & other cubical sets model spaces?
- \mathcal{A} is a (weak/local/-strict) test category when $\hat{\mathcal{A}}$ models the homotopy theory of spaces with certain properties

Homotopy Theory of Categories

- A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is a weak equivalence of categories if $NF: N\mathcal{C} \rightarrow N\mathcal{D}$ is a weak equivalence of simplicial sets
- A natural transformation $\mathcal{C} \xrightarrow{\text{zigzag}} \mathcal{D}$ induces a homotopy $N\mathcal{C} \times \Delta^1 \rightarrow N\mathcal{D}$ from F to G
- Any functor $F: \mathcal{C} \rightarrow \mathcal{C}$ with a zigzag $F \Rightarrow \dots \Leftarrow \text{id}$ to the identity is a w.e.
- \mathcal{C} is contractible if $\mathcal{C} \rightarrow *$ is a w.e. (eg if \mathcal{C} has a terminal object)
- (Quillen Theorem A) For $F: \mathcal{C} \rightarrow \mathcal{D}$ if all $F_{/d}$ are contractible (doubtful) F is a w.e.
- (Thomason) The homotopy theory of categories is equivalent to spaces

Homotopy in $\hat{\mathcal{A}}$

- Consider $i_A: A \rightarrow \text{Cat}: a \mapsto A/a$

$A \xrightarrow{i_A} \text{Cat}$
 $\downarrow \text{slab} \quad \uparrow e_1$
 $A \xrightarrow{e_1} \hat{A}$ is a left Kan extension

$e_1: \hat{A} \rightarrow \text{Cat}$ sends X to its category of elements

Ex If $A = \square$, and $X = \square^2$, $e_1(X) =$

 $\cong \square / \square^2$

- $f: X \rightarrow Y$ is a weak equivalence in $\hat{\mathcal{A}}$ if $e_1(f): e_1(X) \rightarrow e_1(Y)$ is a w.e. in Cat

X is contractible if $\text{el}(X)$ is contractible

Ex Each representable a in \hat{A} is contractible as $\text{el}(a) \cong A/a$, which has a terminal object

Ex For $*$ terminal in \hat{A} , $\text{el}(*) \cong A$, so $*$ is contractible in \hat{A} iff A is contractible in Cat

— i_A also determines a functor $i_A^*: \text{Cat} \rightarrow \hat{A}: C \mapsto (a \mapsto \text{Fun}(A/a, C))$

Test Categories

- A is a weak test category if $\text{el}(i_A^*(C)) \rightarrow C$ is a w.e. for all C (\hat{A} and Cat have equivalent homotopy categories Hot)
- A is a test category if A is contractible and each A/a is a weak test cat
- A is a strict test category if furthermore $\hat{A} \rightarrow \text{Hot}$ preserves products

Ex Δ is a strict test category

Δ_s is a weak test category

— A separated interval in \hat{A} is a pullback square

$$\begin{array}{ccc} \emptyset & \rightarrow & * \\ \downarrow & & \downarrow \\ * & \rightarrow & I \end{array}$$

Ex

$$\begin{array}{ccc} \emptyset & \rightarrow & \square^0 \cong * \\ \downarrow & & \downarrow \delta_{00} \\ \square_0 & \xrightarrow{\delta_{10}} & \square_1 \end{array}$$

is a separated interval in $\hat{\square}_a$ if $\delta_{00} \in a$ and $\delta_{10} \notin a$

$$\begin{array}{ccc} \cdot & \longrightarrow & \cdot \\ \uparrow \delta_{0a} & & \uparrow \end{array}$$

$$\begin{array}{ccc} \cdot & \xleftarrow{\quad} & \cdot \\ \uparrow \delta_{0a} & & \uparrow \end{array}$$

— (Grothendieck) A is a strict test category iff

- \hat{A} has a separated contractible interval
- $a \times b$ is contractible in \hat{A} for all $a, b \in A$

Some Cube Categories are strict test

— Any Cartesian cube category $(\square_a, w, \gamma, \delta \in a)$ is a strict test category

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- Has a separated contractible interval
- $a \times b$ is representable, hence contractible

— (Maltsiniotis) Any cube category with degeneracies and either type of connections ($\sigma, \delta_{\sigma}, \epsilon_a$) is a strict test category

Proof

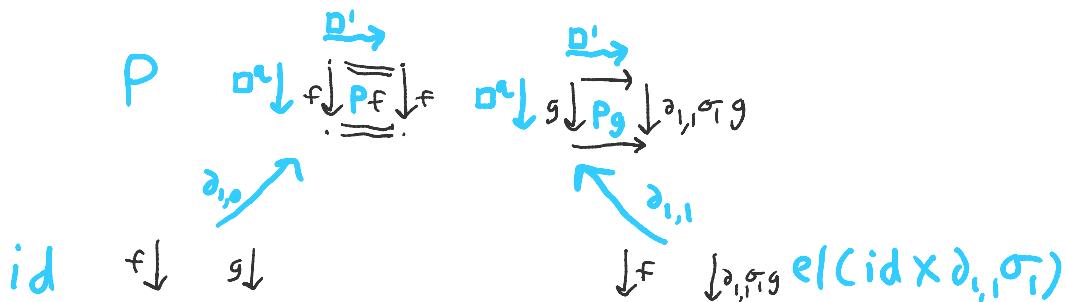
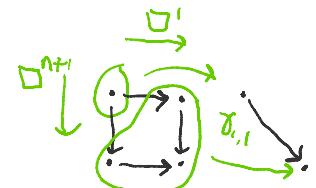
- Has a separated contractible interval
- Will show $\square^m \times \square^n \xrightarrow{\text{id} \times \sigma_1} \square^m \times \square^{n-1} \xrightarrow{\text{id} \times \sigma_1} \dots \xrightarrow{\text{id} \times \sigma_1} \square^m \times \square^0 \cong \square^m$ is a w.e.
- by showing $e1(\square^m \times \square^{n+1}) \xrightarrow{e1(\text{id} \times \sigma_1)} e1(\square^m \times \square^n)$ is a homotopy equivalence in Cat
- homotopy inverse $\xleftarrow{e1(\text{id} \times \delta_{1,1})}$
- Compose to identity on $e1(\square^m \times \square^n)$
- Suffices to find homotopy from id to $e1(\text{id} \times \delta_{1,1}, \sigma_1)$ in $e1(\square^m \times \square^{n+1})$

$e1(\square^m \times \square^{n+1})$ has objects $(\square^a \xrightarrow{f} \square^m, \square^a \xrightarrow{g} \square^{n+1})$

For any f, g , define $P_f : \square^{a+1} \xrightarrow{\sigma_1} \square^a \xrightarrow{f} \square^m$

$P_g : \square^{a+1} \xrightarrow{\text{id}_{\square^a} \otimes g} \square^{n+2} \xrightarrow{\delta_{1,1}} \square^{n+1}$

$P : e1(\square^m \times \square^{n+1}) \rightarrow e1(\square^m \times \square^{n+1})$



The remaining cube categories

— $i_A^* : \text{Cat} \rightarrow \widehat{A}$ is not the only nerve functor

Any $i : A \rightarrow \text{Cat}$ defines $i^* : \text{Cat} \rightarrow \widehat{A} : \mathcal{C} \mapsto (\mathcal{C} \mapsto \text{Fun}(i(\mathcal{C}), \mathcal{C}))$

— Let $\mathcal{D} = \dots \rightarrow \dots$ in Cat

- (Grothendieck) If A is contractible and $i:A \rightarrow \text{Cat}$ is s.t.
- $i(a)$ has a terminal object for all a
 - $\text{el}(\Delta^{\infty} i^*(\Delta))$ is contractible
 - then A is a test category

— (Cisinski) All cube categories are test categories

- Proof
- \square_a is contractible when $a \in a$
 - Let $i(\square^n) = \Delta^n$ (or $(\Delta^n)^n$ if there are reversals)
 - Δ^n has a terminal object
 - will show $\text{el}(\square^n \times i^*(\Delta)) \xrightarrow{\text{pr}} \text{el}(\square^n) \cong \square_{/\square^n}$ is a homotopy equivalence

$\text{el}(\square^n \times i^*(\Delta))$ has objects $(\square^q \xrightarrow{f} \square^n, \Delta^q \xrightarrow{g} \Delta)$

homotopy inverse $\square_a / \square^n \xrightarrow{\text{const}} \text{el}(\square^n \times i^*(\Delta))$

$\square^q \xrightarrow{f} \square^n \mapsto (f, \text{const}: \Delta^q \rightarrow \Delta)$

$\text{pr} \cdot \text{const} = \text{id}$, want homotopy from id to $\text{const} \cdot \text{pr}$

Define $P(f, g) = (\square^{q+1} \xrightarrow{\sigma_i} \square^q \xrightarrow{f} \square^n, \Delta^{q+1} \cong \Delta \times \Delta^q \xrightarrow{\Delta} \Delta)$

$$P \quad f \downarrow \underline{\underline{Pf}} \downarrow f \quad \Delta^q \downarrow g \downarrow \underline{\underline{Pg}} \downarrow \text{const}_i \quad \begin{array}{l} \Delta \downarrow \\ (0, x) \mapsto g(x) \\ (1, x) \mapsto 1 \end{array}$$

$\text{id } f, g$

$f, \text{const}, \text{const pr}$

$\text{id} \Rightarrow P \Leftarrow \text{const} \cdot \text{pr} \quad \square$

— (Maltsiniotis, Buchholtz-Morehouse)

$$\square_{\partial\alpha}, \square_{\partial\alpha\gamma}, \square_{\partial\alpha\beta}, \square_{\partial\alpha\gamma\beta}$$

$\square_{\partial\alpha}, \square_{\partial\alpha\gamma}, \square_{\partial\alpha\beta}, \square_{\partial\alpha\gamma\beta}$ are not strict test categories

Proof In the case of $\square_{\partial\alpha}$, \square' contains $\overset{\circ}{\square} \xrightarrow{x} \overset{\circ}{\square}$

$\text{el}(\square' \times \square')$ is (by Theorem A) weakly equivalent to the full subcategory on the elements

view category on the elements

$$\{(0,0), (0,1), (1,0), (1,1)\} = (\square' \times \square')$$

$$\{(\sigma_0, x), (\sigma_1, x), (x, \sigma_0), (x, \sigma_1), (x, x)\} = (\square' \times \square')$$

$$\{(\sigma_1 x, \sigma_2 x), (\sigma_2 x, \sigma_1 x)\} \subseteq (\square' \times \square')$$

