# Contract Design Under Uncertainty

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# Modern Algorithms in Society

• Interactions with **self-interested** individuals

• In AGT, we take participants' incentives into account



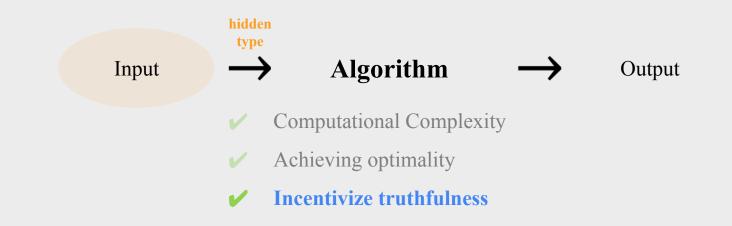
# Social Media Marketing

- Paid Ads in the social platform feed
- Influencer Marketing- a brand hires popular users



#### Paid Ads

 $\circ$  Algorithmic **auction** - advertisers bid  $\rightarrow$  allocation and payments



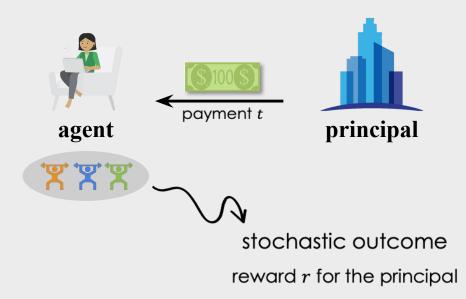
# Influencer Marketing

• The algorithm determines a **contract** to incentivize effort



# The Principal-Agent Problem [GH83]

- Moral hazard the agent's actions cannot be observed
- Objective: a contract maximizing expected rewards minus payment  $\mathbb{E}[r-t]$



# Research Agenda

Generalizations of the classic model:

- **Personalization** for participants from diverse population (with types)
- Multilateral contracts involving multiple principals or agents
- The need for simple contracts

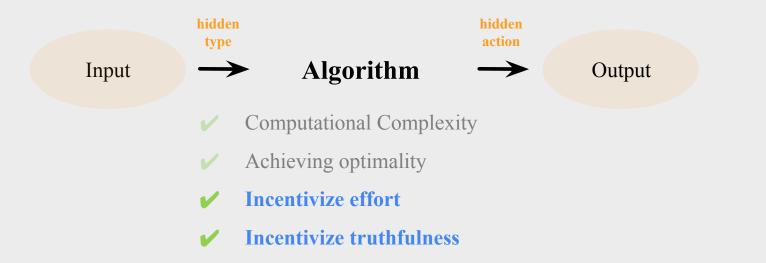
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Generalizations of the classic model:

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- Multilateral contracts involving multiple principals or agents

• The need for simple contracts

### Research Agenda



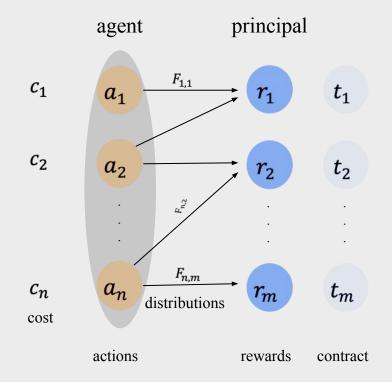
## Outline

- Single-parameter model of types [ADT.C. EC'21]
  - Motivated by single-parameter auction design
- Characterization of the design space [ADT.C. EC'21]
- Counter-intuitive and undesirable properties of optimal contracts [ADLT.C. EC'23]
- Linear contracts (aka commission-based) are near-optimal [ADLT.C. EC'23]

Recent works on contracts with types. Myerson )1982(, Guruganesh et al. )2021), Castiglioni et al. )2021), Gottlieb and Moreira )2022(, Casto-Pire et al. )2022(.

#### The Model

 $\circ$  The classic principal-agent model by Grossman-Hart (1983)



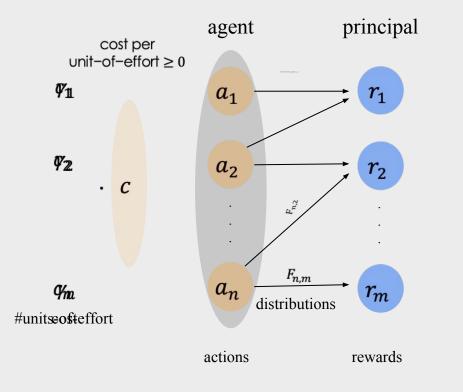
**Notation.**  $T_i = \mathbb{E}_{j \sim F_i}[t_j], R_i = \mathbb{E}_{j \sim F_i}[r_j]$ 

• Agent's action  $i^*$  maximizes  $T_i - c_i$ 

• Principal's revenue  $R_{i^*} - T_{i^*}$ 

#### The Model

• Type is drawn from G with density g supported on  $C = [\underline{c}, \overline{c}]$ 



#### The Model

• A contract (x, t) consists of two mappings:

• Payment rule  $t: C \rightarrow \mathbb{R}^m$  from types to a paymet scheme

• Allocation rule  $x: C \rightarrow [n]$  from types to an action recommendation

Notation.  $T_i^c = \mathbb{E}_{j \sim F_i}[t(c)_j]$ 

• Agent *c* report  $\hat{c}$  and action  $i^*(c)$  maximize  $T_{i^*(c)}^{\hat{c}} - \gamma_{i^*(c)}c$ 

• A contract (x, t) is incentive compatible (IC) if  $c = \hat{c}$  and  $x(c) = i^*(c)$ 

• Principal's contract maximizes  $\mathbb{E}_{c\sim G}[R_{x(c)} - T_{x(c)}^{c}]$  s.t. IC

• Welfare sum of utilities  $\mathbb{E}_{c \sim G}[R_{i^*(c)} - \gamma_{i^*(c)}c]$ 

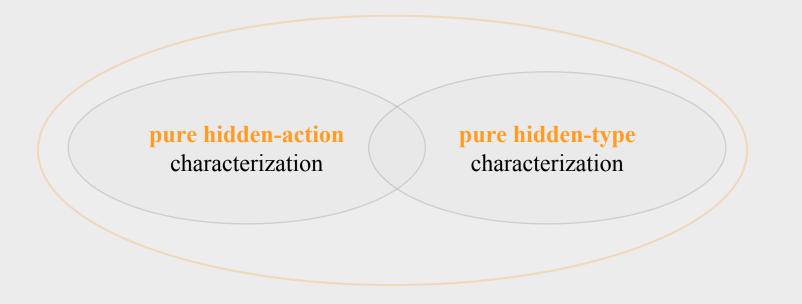
**Definition.** Allocation rule x is **implementable** if exists payment rule t s.t. (x, t) is IC

Q. What do implementable allocation rules look like?

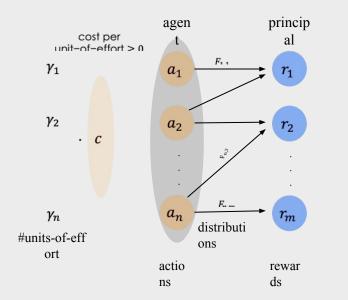
pure hidden-action characterization pure hidden-type characterization

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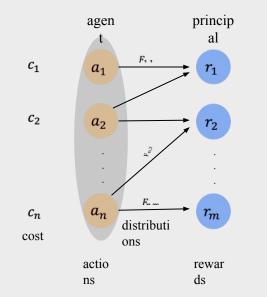
#### Pure hidden action



#### Pure hidden action

• x is **implementable**  $\Leftrightarrow$  exists no **deviation scheme**  $\lambda_k$  s.t. (1) dominant distribution  $\sum_k \lambda_k F_k \ge F_x$ 

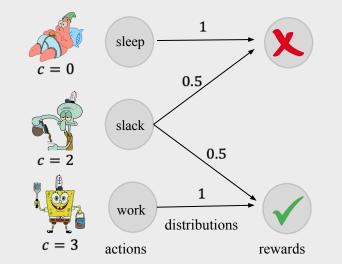
(2) strictly lower cost  $\sum_k \lambda_k c_k < c_x$ 



Pure hidden action

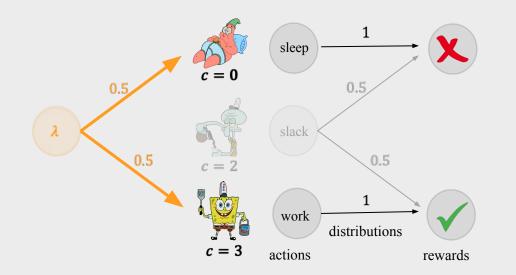
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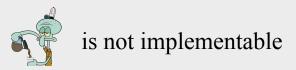
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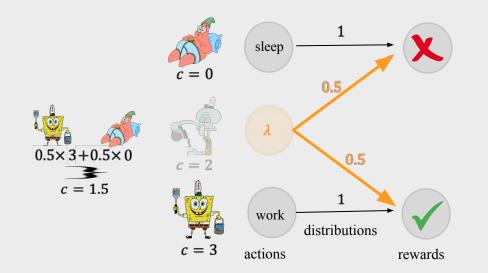
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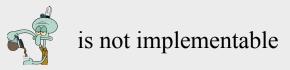




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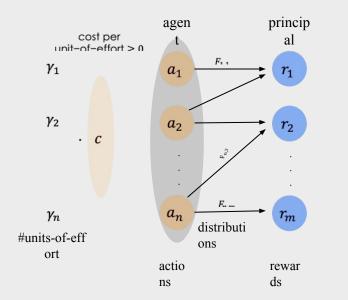
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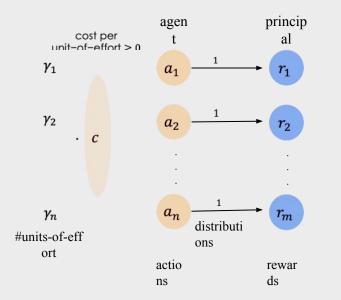
• LP duality approach

Pure hidden type



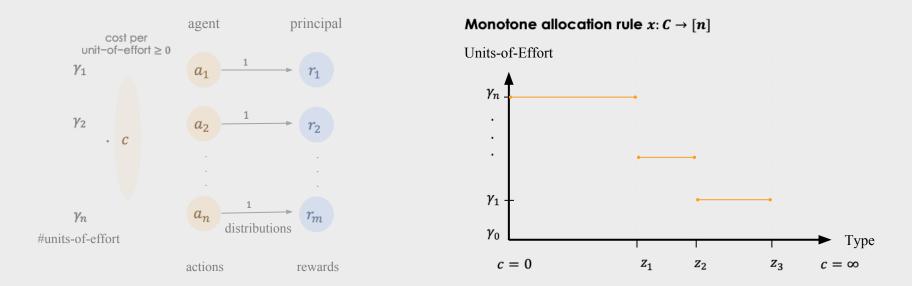
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• x is implementable  $\Leftrightarrow$  x is monotone [Myerson 1981]

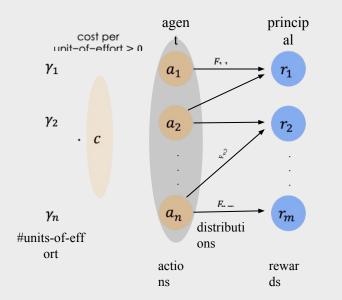


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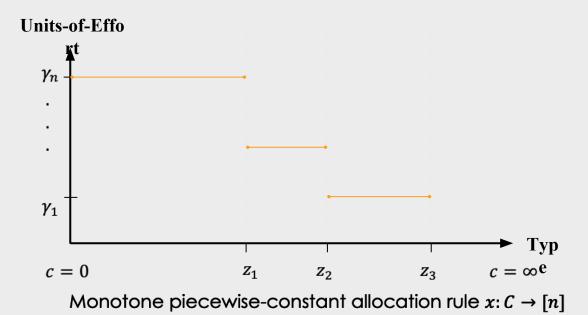


Hidden type and hidden action



#### Hidden type and hidden action

#### **Proposition [ADT EC'21].** If x is implementable, it is monotone



**Theorem [ADT EC'21].** x implementable  $\Leftrightarrow$  exists no deviation scheme  $\lambda_{(z,k)}$  s.t.

- (1) dominant sum of distributions  $\sum_{z,k} \lambda_{(z,k)} F_k \ge \sum_z F_{x(z)}$
- (2) strictly lower joint cost  $\sum_{z,k} \lambda_{(z,k)} \gamma_k z < \sum_z \gamma_{x(z)} z$

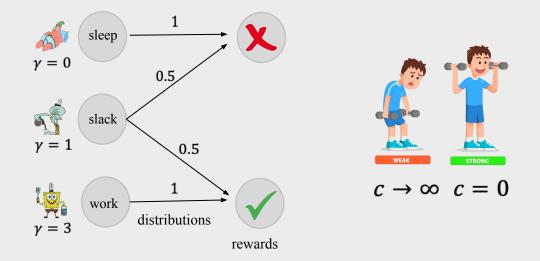
**Corollary** [ADT EC'21]. Optimal contract is **polytime computable** for const #actions

Hardness for constant #actions in the multi-parameter model [Guruganesh-Schneider-Wang'21]

#### Example

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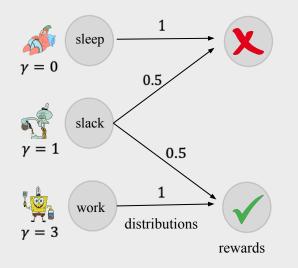
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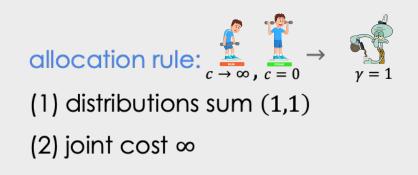


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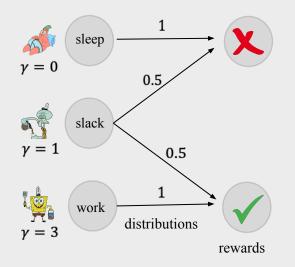




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allocation rule:  $(c, q) = 0 \rightarrow (c, q) \rightarrow (c, q)$ 

# Optimal Contracts and Their Issues

- Informational requirements, extensive analysis, etc.
- Unintuitive, e.g., non-monotonicity in rewards [DRT EC'19]

**Theorem [ADLT EC'23].** In the single dimensional typed model

- Large menu-size complexity
- Revenue **non-monotonicity** w.r.t type distribution

# Simple Contracts

◦ In a linear contract, the principal offers a fixed share  $\alpha \in [0,1]$  of the rewards



"It is probably the great robustness of linear rules based on aggregates that accounts for their popularity. That

point is not made as effectively as we would like by our model; we suspect

that it cannot be made effectively in any traditional Bayesian model." [Milgrom and Holmstrom 1987]

# Simple Contracts

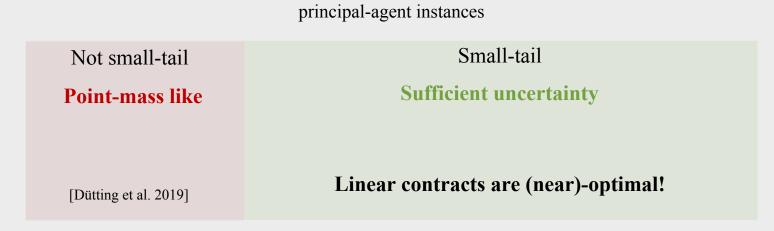
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- Robustness of linear contracts. Carroll )2015(, Dütting et al. )2019(, Yu and Kong )2020(, Dai and Toikka )2022(, Walton and Carroll )2022(
- Approximation of linear contracts. Dütting et al. )2019(, Castiglioni et al. )2021(, Guruganesh et al. )2021(

# Near-Optimality of Linear Contracts

- o  $\theta(n)$  separation for point-mass distributions [DRT EC'19]
  - o Boundary case
- o Approximately optimal with sufficient uncertainty
  - o The small-tail assumption



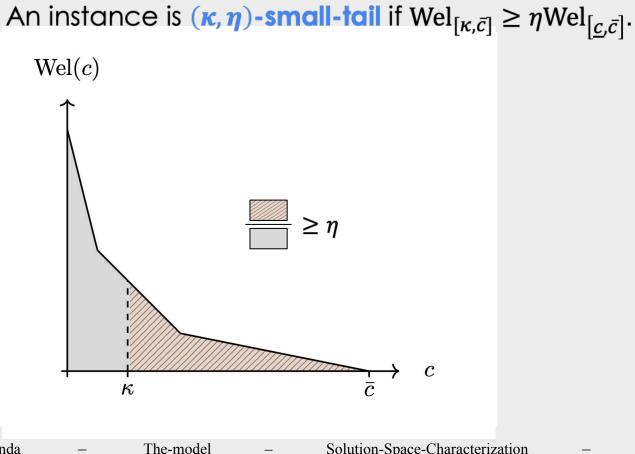
#### Near-Optimality of Linear Contracts

**Theorem [ADLT EC'23]**. **Revenue** benchmark:

- **3-approximation** for normal  $\mathcal{N}(\mu, \sigma^2)$  truncated at c = 0 with  $\sigma \ge 5\eta/2\sqrt{2}$
- **2-approximation** for **uniform**  $U[0, \bar{c}]$ 
  - **Optimal** when  $i^*(r, \bar{c}) = 0$
- o 2-approximation for decreasing densities (e.g., exponential)
- Constant approximation w.r.t optimal welfare benchmark [ADLT EC'23]

# The Small-Tail Assumption

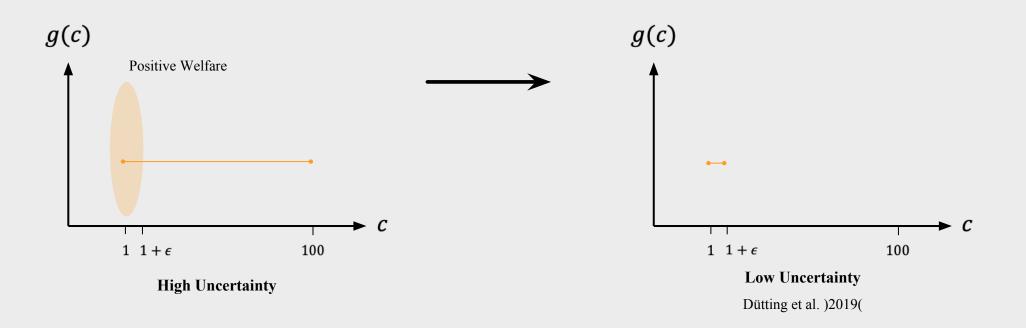
**Definition [ADLT EC'23]**. Let  $\kappa \in [\underline{c}, \overline{c}], \eta \in [0,1]$ .



Agenda

# The Small-Tail Assumption

Depends on the entire principal-agent setting



#### Universal Approximation Guarantee

**Theorem [ADLT EC'23].** Let  $q \in (0,1)$  and  $G(c_q) = q$ . If for  $\alpha, \eta \in (0,1)$  the settings is

 $(\frac{c_q}{\alpha}, \eta)$ -small-tail then linear contract  $\alpha$  is at least  $(1 - \alpha)\eta q$  of the optimal welfare

# Slowly-Increasing Distributions

• Applies to any CDF and captures its **rate of increase** 

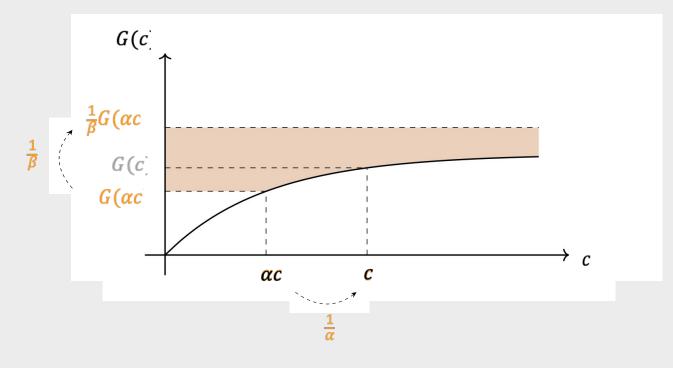
• Parametric approximation of linear contracts

• Any distribution is slowly-increasing for some parameters

#### **Slowly-Increasing Distributions**

**Definition [ADLT EC'23].** Let  $\alpha, \beta \in (0,1)$ , and  $\kappa \in [\underline{c}, \overline{c}]$ .

A distribution G is  $(\alpha, \beta, \kappa)$ -slowly-increasing if  $G(c) \le \frac{1}{\beta}G(\alpha c) \ \forall \kappa \le c$ 



# Approximation for Slowly Increasing

**Theorem [ADLT EC'23].** Let  $\alpha, \beta, \eta \in (0,1)$ , and  $\kappa \in [\frac{c}{\alpha}, \overline{c}]$ .

Under  $(\alpha, \beta, \kappa)$ -slowly-increasing and  $(\kappa, \eta)$ -small-tail

linear contract  $\alpha$  is  $(1 - \alpha)\beta\eta$  of the optimal welfare

#### Proof Idea for Slowly Increasing

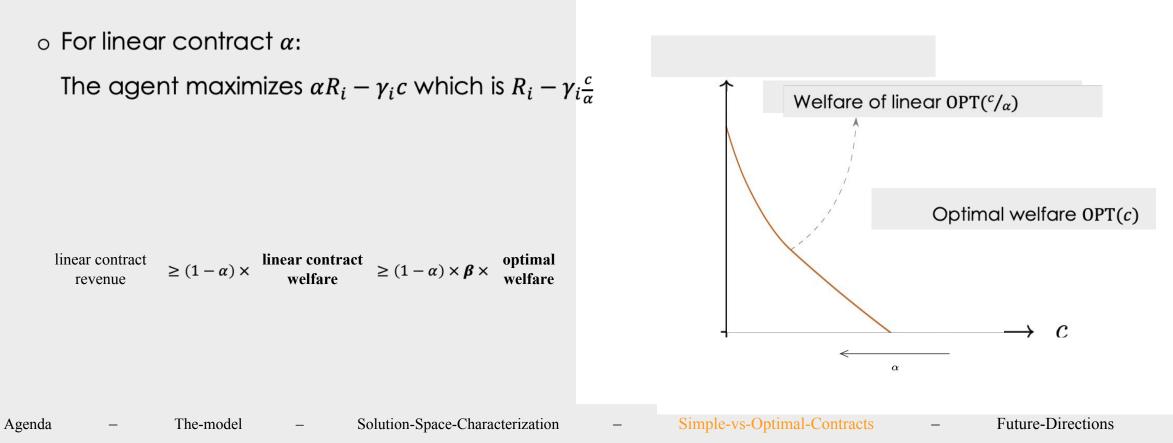
**Step 1.** Revenue of linear contract  $\alpha$  is at least  $1 - \alpha$  of its welfare

$$\begin{split} \text{revenue of linear contract} & \text{welfare of linear contract} \\ \mathbb{E}_{c\sim G}[(1-\alpha)R_{x(c)}] & \geq (1-\alpha) \times \\ \mathbb{E}_{c\sim G}[R_{x(c)}-\gamma_{x(c)}c] \end{split}$$

### Proof Idea for Slowly Increasing

**Step 2.** welfare of linear contract  $\alpha$  is at least  $\beta$  of the optimal welfare

 $\circ$  When maximizing welfare, the agent maximizes  $R_i - \gamma_i c$ 



# Summary and Future Directions

- Single-parameter model of types
- Characterization of the design space
- Counter-intuitive and **undesirable** properties of **optimal contracts**
- Linear contracts are near-optimal

#### **Future directions:**

- Other forms of simple contracts that are near-optimal
- Contracts that involve multiple agents
- $\circ$  Other models for incentivizing effort

#### !Thank You

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### Research Agenda

Generalizations of the classic model:

- Personalization for participants from diverse population (with types)
- Multilateral contracts involving multiple principals or agents
  - Bad welfare performance for independent contracts.

Identify unique welfare-maximizing for interdependent contracts [ALST EC'21,OR'23]

• The need for simple contracts

# Research Agenda

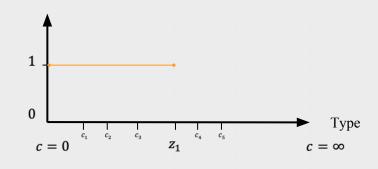
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#### Pure hidden type

• x is implementable  $\Leftrightarrow$  x is monotone [Myerson 1981]

Units-of-Effort



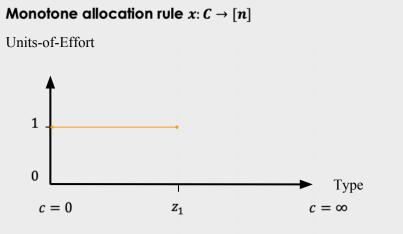
Notation.  $S = \{c \mid x(c) = 1\}$ Existence and the integrated service  $\sum_{c} \lambda_{c,1} = \sum_{c \in S} 1$ (2) lower joint cost  $\sum_{c} \lambda_{c,1} c < \sum_{c \in S} c$ 

#### TODO!!!

#### Fix citations

#### Absent hidden action

• x is implementable  $\Leftrightarrow$  x is monotone [Myerson 1981]



o **allocation rule:** 
$$x(c) = \begin{cases} 1, c \in [0, z_1), \\ 0, c \in [z_1, ∞) \end{cases}$$

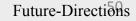
• payment rule: 
$$p(c) = \begin{cases} z_1, & c \in [0, z_1), \\ 0, & c \in [z_1, \infty) \end{cases}$$

• **utility** of serving:  $z_1 - c$ , not serving: 0

# Principals with Types

Agenda

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# Common Agency

- Multiple principals hire a single agent
- What if the **principals** have different interests?

Interests form principals' privately-known types



### TODO

•We study welfare-maximizing contracts with hidden principal types

•Such contracts incentivize the most rewarding action in aggregate



# Common Agency

#### Theorem [Alon Lavi Shamash & T.-C. EC'21].

The bad:

• In some settings welfare maximization is **impossible** 

The good:

• The class of settings with possibility is **large and identifiable** 

- For this class, welfare maximization is **tractable**
- The contracts are **unique** and make economic sense

# Summary

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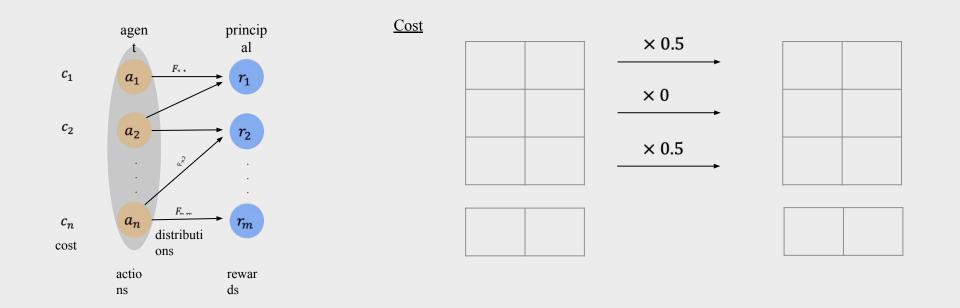
### Some more work

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#### Absent hidden type

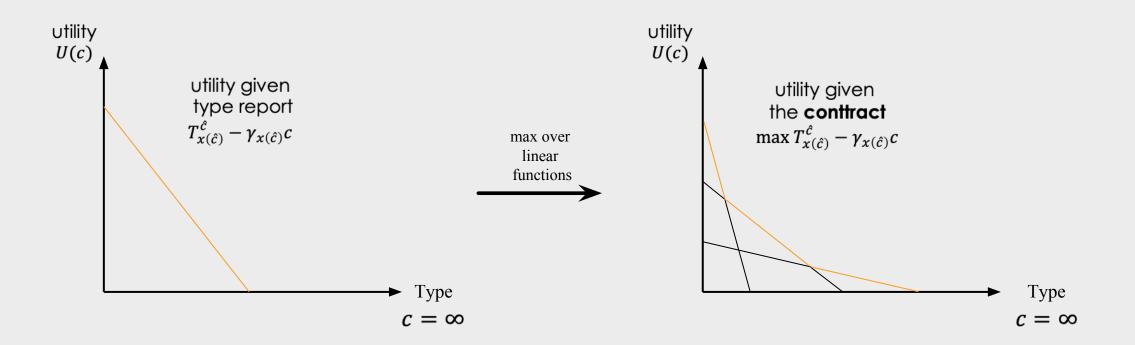
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○ x is implementable  $\Leftrightarrow$  [GH 1983]



### Intuition: Monotonicity Isn't Sufficient

Because

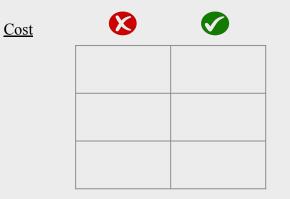


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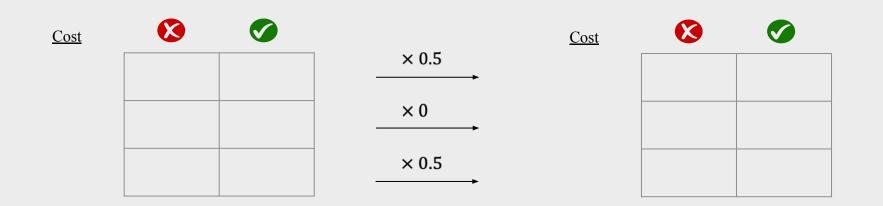


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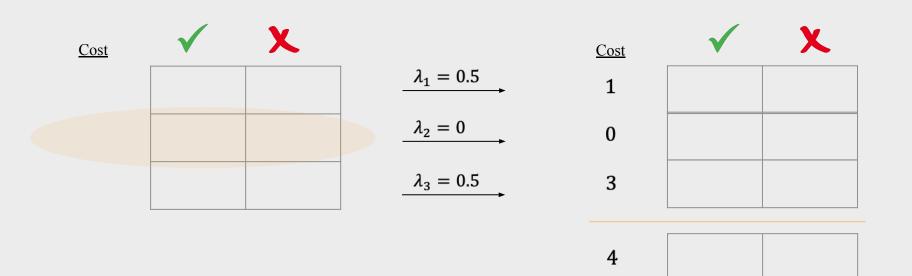


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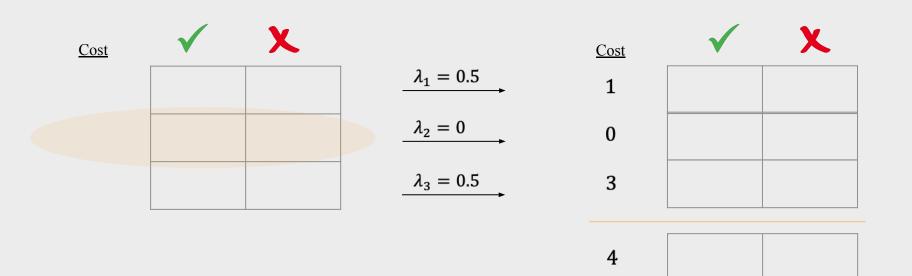
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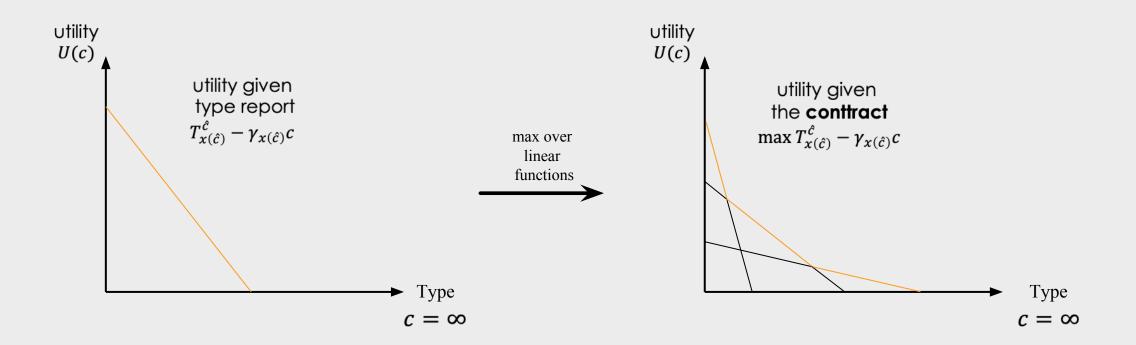
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### Intuition: Monotonicity Isn't Sufficient

#### No hidden-action

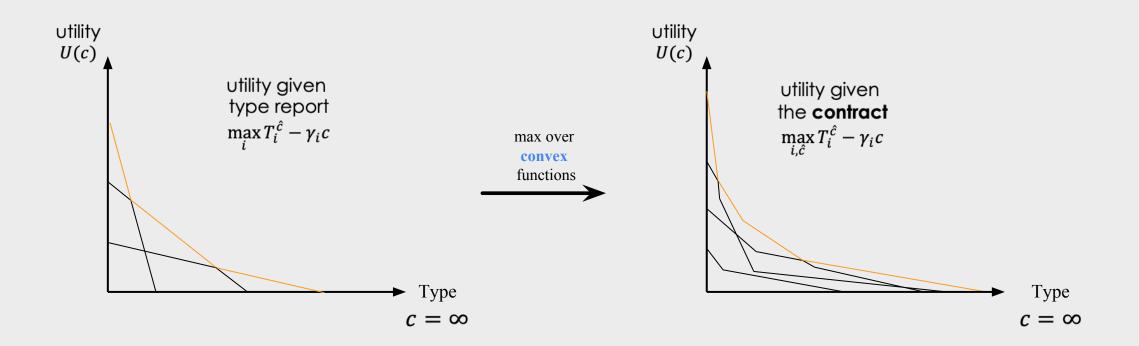
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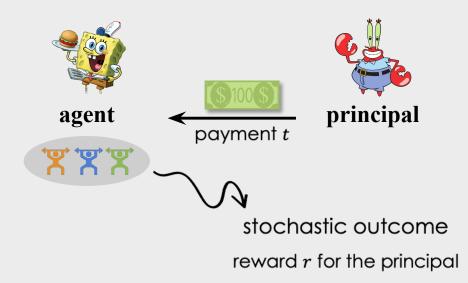
#### With hidden-action

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# The Principal-Agent Problem [GH83]

- Moral hazard the agent's actions cannot be observed
- Objective: a contract maximizing expected rewards minus payment  $\mathbb{E}[r-t]$



# Contract Theory Approach

**Proposition [Alon Dütting and T.-C. EC'21].** The following LP computes if x is **implementable** 

$$\begin{array}{ll} \max & 0 \\ \text{s.t.} & F_{x(z)}t^{z} - \gamma_{x(z)}z \geq F_{k}t^{z'} - \gamma_{k}z \; \forall z, z' \in C, k \in [n] \\ & t_{j}^{z} \geq 0 \end{array}$$

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