## SCHOOL CHOICE AND FAIRNESS PSET

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**Problem 1.** Give an example of an economy where the Boston Mechanism is not strategyproof.

Problem 2. The goal of this problem is to prove Theorem 1. This says that for a regular market (see Definition 4 on page 12), any set of budget probabilities can be implemented by DA with quotas and priorities. More formally, Given a market M and an arbitrary budget set probability matrix  $y \in Y^M$ , define the priority distribution  $G_t$  for each segment t as follows. For an agent i, sample a set  $S_i \subseteq J$  according to the probability vector  $y_t$  (with probability  $y_{tS_i}$ ). Define the agent's priority score for item j as

$$
\pi_{ij} = \mathbb{1}_{(j \in S_i)} + \delta_i
$$

, where  $\delta_i \sim \text{Uniform}(0, 1)$ . Define quota  $q_i$  as the left hand side of (11), which is the mass of agents assigned to j under budget set probabilities y. If M is regular, then the DA mechanism with priority distribution G and budget quota vector  $q$  implements  $y$ .

- a Let z be an *n*-dimensional vector of all 1's. Show that z is a fixed point of the DA operator (see A.2), ie.  $DA(z) = z$ .
- b Let X be the priority-based allocation mechanism associating each priority  $\pi \in \Pi$  with the budget set

$$
B_{\pi}^X = \{ j \in J : j = 0 \text{ or } \pi_j \ge z_j \}
$$

Show that for any agent segment t and any  $S \subseteq J$  this set is equal to S with probability  $y_{tS}$ 

c Let  $z' = z^{DA(M,G,q)}$  be the DA cutoff. This is the minimum element of the lattice of fixed points,  $\{z' \in \Pi : DA(z') = z'\}.$  Define X' to be DA mechanism with priority distribution G and quota q. This mechanism associates each priority  $\pi \in \Pi$  with the budget set  $B_{\pi}^{X'}$ . Show that the budget set probabilities for the DA mechanism and our mechanism form part b are the same, ie.  $y^X = y^{X'}$ 

Problem 3. We will show that in Example 1 of Section 5, the neighborhood assignment plan maximizes utilitarian welfare among all priority-based allocation mechanisms if and only if the following inequality holds:

$$
E[u_0 - \alpha | u_0 \ge \alpha] \ge E[max(u_0, u_1) - \alpha | max(u_0, u_1) \ge \alpha],
$$

where  $u_0 \sim F_0$ ,  $u_1 \sim F_1$  and  $\alpha \sim H$  and the random variables are independent. On the other hand, the open enrollment plan (RSD) is optimal if and only if the above inequality is reversed.

- a Consider the LP setup of the priority allocation problem from lecture, and define the sets  $S_0 =$  $\{0\}, S_1 = \{0, 1\}, S_2 = \{0, 2\}, S_3 = \{0, 1, 2\}.$  Show that without loss of generality it suffices to consider feasible solutions  $y \in Y^M$  that are symmetric in the sense that  $y_{1S_3} = y_{2S_3}, y_{1S_0} = y_{2S_0}, y_{1S_1} =$  $y_{2S_2}, y_{1S_1} = y_{2S_2}$ , and  $y_{1S_2} = y_{2S_1}$ .
- b Rewrite the LP using these symmetry conditions, with decision variable  $z_k$  replacing  $y_{1S_k}$  for  $k =$ 1, 2, 3. and  $c = c_1 = c_2$  denoting the common capacity for each school.
- c Show that in any optimal solution z we must have  $z_2 = 0$ .
- d Plugging in  $z_2 = 0$  yields that the feasible region is now a triangle. What are the three endpoints of that triangle?
- e Show that the vertex  $(0, 0, 0)$  is never optimal, and translate the conditions where each of the non-zero vertices are optimal into the statement of the proposition.

Problem 4. Nguyen et al discussion: Do we agree with the issues pointed out by this paper about school choice? Is there anything we'd add or de-emphasize?

Problem 5. Hitzig discussion: What were your main takeaways of Hitzig's argument? Do you agree with it?