

**(Univariate sumcheck for additive subgroups)** We saw how to design a sumcheck protocol for univariate polynomials  $g \in \mathbb{F}[X]$  when summing over *multiplicative* subgroups  $H$  of  $\mathbb{F}$  with  $\deg(g) < |H|$ , using the identity

$$\sum_{a \in H} g(a) = |H|g(0) .$$

In this problem we will design a univariate sumcheck protocol over *additive* subgroups  $H$ .

**Problem 1.** Using the fact that for every univariate polynomial  $g \in \mathbb{F}[X]$  with  $\deg(g) < |H|$ , letting  $\beta$  be the coefficient of  $X^{|H|-1}$  in  $g$ , it holds that

$$\sum_{a \in H} g(a) = \beta \cdot \sum_{a \in H} a^{|H|-1} ,$$

design an efficient univariate sumcheck protocol for additive subgroups. Let  $v_H(X)$  be the vanishing polynomial for  $H$ . You may assume that  $\sum_{a \in H} a^{|H|-1}$ , and  $v_H(\gamma)$  for any  $\gamma \in \mathbb{F}$ , can be computed in time  $\text{polylog}(|H|)$ .

**Problem 2.** In this question we will prove the identity used above. Let  $\mathbb{F}$  be a field of characteristic  $p$  (the prime field  $\mathbb{F}_p$  is a subfield of  $\mathbb{F}$ ). The *derivative* of a function  $f: \mathbb{F} \rightarrow \mathbb{F}$  in direction  $a \in \mathbb{F}$  is  $\Delta_a(f)(x) := \sum_{b \in \mathbb{F}_p} f(x + ba)$ . For  $a_1, \dots, a_k \in \mathbb{F}$  we inductively define  $\Delta_{a_1, \dots, a_k}(f) := \Delta_{a_1}(\Delta_{a_2, \dots, a_k}(f))$ .

1. Let  $a_1, \dots, a_k \in \mathbb{F}$  be a basis for  $H$  over  $\mathbb{F}_p$ . Prove that  $\Delta_{a_1, \dots, a_k}(f)(a_0) = \sum_{a \in H} f(a_0 + a)$ .
2. Write the  $p$ -ary expansion of an integer  $c \in \mathbb{N}$  as  $\sum_{i \geq 0} c_i p^i$  for  $0 \leq c_i < p$ . Define the sum of the “ $p$ -ary digits” of  $c$  as  $\text{ds}_p(c) := \sum_{i \geq 0} c_i$ . For a polynomial  $g(X) = \sum_{j \geq 0} \alpha_j X^j$ , define  $\text{ds}_p(g) := \max(\{\text{ds}_p(j) : \alpha_j \neq 0\} \cup \{-1\})$ . Prove that for every  $\alpha \in \mathbb{F}$  it holds that

$$\text{ds}_p(\Delta_a(g)) \leq \max\{\text{ds}_p(g) - (p - 1), -1\} .$$

You may use the following facts:

- (a)  $b^p = b$  for all  $b \in \mathbb{F}_p$ .
- (b) For  $c, d \in \mathbb{N}$ , let  $c = \sum_i c_i p^i$  and  $d = \sum_i d_i p^i$  be their  $p$ -ary expansions. If there exists  $i \geq 0$  such that  $d_i > c_i$  then  $\binom{c}{d} \equiv 0 \pmod{p}$ .

*Hint: consider a single monomial  $X^c$ , and apply the binomial theorem.*

3. Prove that for every univariate polynomial  $g \in \mathbb{F}[X]$  with  $\deg(g) < |H|$ , letting  $\beta$  be the coefficient of  $X^{|H|-1}$  in  $g$ , it holds that

$$\sum_{a \in H} g(a) = \beta \cdot \sum_{a \in H} a^{|H|-1} .$$

**Problem 3.** Show that  $\sum_{a \in H} a^{|H|-1} = \prod_{a \in H \setminus \{0\}} a$ . (*Hint: use Newton’s identities.*<sup>1</sup>) The right-

<sup>1</sup>[https://en.wikipedia.org/wiki/Newton's\\_identities](https://en.wikipedia.org/wiki/Newton's_identities)

hand side is equal to the coefficient of the linear term of  $v_H(X)$ , which can be computed in time  $O(\log^2 |H|)$ .