

Problem 1. (PCP to extremely laconic MA) Consider a language $\mathcal{L} \in \text{PCP}[\epsilon_c, \epsilon_s, \Sigma, l, q, r]$. Let $\text{PCP} = (P, V)$ be a PCP for the language \mathcal{L} . Recall the following PCP to laconic one-round MA transformation from lecture:

- $\mathcal{P}_{\text{MA}}(\mathbf{x})$:
 1. Compute $\pi \leftarrow P(\mathbf{x})$.
 2. Sample random $\mathcal{Q} \leftarrow \binom{[l]}{q}$.
 3. Send $\Pi := (\mathcal{Q}, \pi[\mathcal{Q}])$.
- $\mathcal{V}_{\text{MA}}(\mathbf{x}, \Pi)$:
 1. Sample $\rho \leftarrow \{0, 1\}^r$.
 2. Parse Π as $(\mathcal{Q}, a \in \Sigma^q)$.
 3. Check that $V(\mathbf{x}; \rho) = 1$ by answering queries $i \in \mathcal{Q}$ with $a[i]$. Reject if V queries anything outside of \mathcal{Q} .

The above transformation gives us a MA protocol with completeness error $1 - (1 - \epsilon_c) \cdot 2^{-q \cdot \log l}$, soundness error ϵ_s , and prover communication $q \cdot (\log l + \log |\Sigma|)$.

We are not happy about the prover communication and want to improve the above transformation. Consider the following transformation and argue why the resulting MA protocol has soundness error 1.

- $\mathcal{P}_{\text{MA}}(\mathbf{x})$:
 1. Compute $\pi \leftarrow P(\mathbf{x})$.
 2. Sample random $\mathcal{Q} \leftarrow \binom{[l]}{q}$.
 3. Send $\Pi := \pi[\mathcal{Q}]$.
- $\mathcal{V}_{\text{MA}}(\mathbf{x}, \Pi)$:
 1. Sample $\rho \leftarrow \{0, 1\}^r$.
 2. Check that $V(\mathbf{x}; \rho) = 1$ by answering the i -th query made by V with $\Pi[i]$.

Problem 2. (PCP to laconic MA, a different approach) Consider a language $\mathcal{L} \in \text{PCP}[\epsilon_c, \epsilon_s, \Sigma, l, q, r]$. Design a one-round MA protocol for the language \mathcal{L} with prover communication $r + q \cdot \log |\Sigma|$. What is the completeness error and soundness error for your MA protocol?