

Another comb. char. of triang's:

THM: [de Loera et al.]

$\sigma \neq \tau \in \binom{\mathcal{A}}{d+1}$ is a triang of $\text{conv } \mathcal{A}$ iff

(CP) For every interior facet of a simplex $\sigma \in T$ there are two simplices

$$\sigma, \sigma' \in T, \sigma \neq \sigma', \tau \subset \sigma, \tau \subset \sigma'$$

s.t.

$$\chi(\tau, \sigma \setminus \tau) = -\chi(\tau, \sigma' \setminus \tau) \quad \text{"cocircuit Property"}$$

(EP) $\text{conv } \mathcal{A}$ has a simplicial facet or there is a point p in the interior of $\text{conv } \mathcal{A}$ in general position that is contained in $\text{conv } \sigma$ for exactly one simplex $\sigma \in T$. "Extension Property"

Rem.: (EP) can also be checked by chirotope values.

Rem.: This one is implemented in TOPCOM.

Obs.: ▷ Check triang. prop. ✓

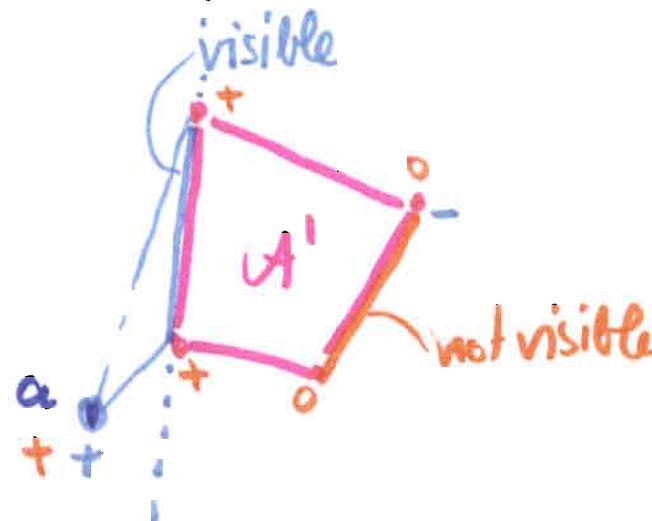
Visible facets and cocircuits:

Prop.: Let $\mathcal{U}' \subsetneq \mathcal{U} \subset \mathbb{R}^d$ and $a \in \mathcal{U} \setminus \mathcal{U}'$.

A fault a is visible by \mathcal{U}' iff the cocircuit defined by the facet has different signs on a and on at least one other point in \mathcal{U}' .
 ("+" "-" or "-" "+", not "+" "0" etc.)

Cor.: Placing triang. can be computed via cocircuit signatures

Example:



Obs.: ▷ Find one triang. ✓

BFS vs. Equivariant BFS (EBFS):

L10

Alg. BFS

Input: $G = (V, E), v_0 \in V, \Gamma: V \rightarrow 2^V$ oracle

Output: $|V| = v_0$ for brevity

1. $\text{cnt} \leftarrow 1, \mathcal{G} \leftarrow \{(v_0, \Gamma(v_0))\}, W \leftarrow \emptyset$
2. While $\mathcal{G} \neq \emptyset$
3. For all $v \in \mathcal{G}$
4. For all $w \in \Gamma(v)$ unmarked
5. If $w \in \mathcal{G} \cup W$
6. mark $w \in \Gamma(v)$
7. else
8. mark $w \in \Gamma(v)$
9. mark $v \in \Gamma(w)$
10. $W \leftarrow W \cup \{w\}, \text{cnt} \leftarrow \text{cnt} + 1$
11. endif
12. endfor
13. endfor
14. $\mathcal{G} \leftarrow W, W \leftarrow \emptyset$
15. endwhile
16. return cnt

Alg. EBFS

Input: $G = (V, E), v_0 \in V, \Gamma: V \rightarrow 2^V$ oracle, $\Pi(G)$ oracle

Output: $|V|$

1. $\text{cnt} \leftarrow 1, \mathcal{G} \leftarrow \{v_0\}, W \leftarrow \emptyset$
2. While $\mathcal{G} \neq \emptyset$
3. For all $v \in \mathcal{G}$
4. For all $w \in \Gamma(v)$ unmarked
5. If $\exists \pi \in \Pi(\mathcal{G}): \Pi(w) \in \mathcal{G} \cup W$
6. mark $\Pi(v)$ in $\Gamma(\Pi(w))$
7. mark w in $\Gamma(v)$
8. else
9. mark w in $\Gamma(v)$
10. mark v in $\Gamma(w)$
11. $W \leftarrow W \cup \{w\}, \text{cnt} \leftarrow \text{cnt} + 1$
12. endif
13. endfor
14. endfor
15. $\mathcal{G} \leftarrow W, W \leftarrow \emptyset$
16. endwhile
17. return cnt

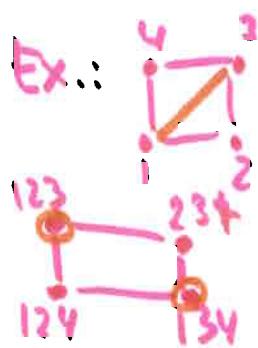
Find comp. of Rippgr.

Obs.: Can be applied to flip graph: symm. in $V \xrightarrow{\sim}$ symm. in G_1 .

Triang's and independent sets

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Def.: The conflict graph on the set of all d -simplices in Δ is the graph $G^d = (V, E)$ with

Ex.: 

$$V := \left\{ \sigma \in \binom{\Delta}{d+1} : \dim \sigma = d \right\}$$

$$E := \left\{ (\sigma, \sigma') : \sigma \text{ and } \sigma' \text{ do not intersect properly} \right\}$$

$$\Leftrightarrow \text{conv} \sigma \cap \text{conv} \sigma' \neq \text{conv}(\sigma \cap \sigma')$$

$$\Leftrightarrow \exists \text{ circuit } (C_+, C_-) \text{ with } C_+ \subseteq \sigma, C_- \subseteq \sigma'$$

Prop.: Every triang. of Δ is a maximal independent set in G^d .

Example: dim 1 and 2: triang's $\xrightarrow{1 \text{ to } d}$ max. indep. sets.

dim 3: \exists partial triang's that cannot be extended.

Obs.: Enumeration of all max. indep. sets yields all non-extendable partial triang's. Checking (UP) \Rightarrow all triang's.

D Find all triang's. ✓ (e.g., backtracking DFS)

Rem.: Do not know to use symmetries here.