Foundations and Frontiers of Probabilistic Proofs (Summer 2021) Worksheet A.8: Linear-Size IOPs for Circuits Date: 2021.08.05

Problem. (Univariate sumcheck for additive subgroups) We saw how to design a sumcheck protocol for univariate polynomials  $g \in \mathbb{F}[X]$  when summing over *multiplicative* subgroups H of F with  $deg(g) < |H|$ , using the identity

$$
\sum_{a \in H} g(a) = |H|g(0) .
$$

In this problem we will design a univariate sumcheck protocol over *additive* subgroups H.

**Problem 1.** Using the fact that for every univariate polynomial  $g \in \mathbb{F}[X]$  with deg(g) < |H|, letting  $\beta$  be the coefficient of  $X^{|H|-1}$  in g, it holds that

$$
\sum_{a \in H} g(a) = \beta \cdot \sum_{a \in H} a^{|H|-1}
$$

,

design an efficient univariate sumcheck protocol for additive subgroups. Let  $v_H(X)$  be the vanishing polynomial for H. You may assume that  $\sum_{a\in H} a^{|H|-1}$ , and  $v_H(\gamma)$  for any  $\gamma \in \mathbb{F}$ , can be computed in time  $\mathsf{polylog}(|H|)$ .

**Problem 2.** In this question we will prove the identity used above. Let  $\mathbb{F}$  be a field of characteristic p (the prime field  $\mathbb{F}_p$  is a subfield of  $\mathbb{F}$ ). The *derivative* of a function  $f: \mathbb{F} \to \mathbb{F}$  in direction  $a \in \mathbb{F}$  is  $\Delta_a(f)(x) := \sum_{b \in \mathbb{F}_p} f(x + ba)$ . For  $a_1, \ldots, a_k \in \mathbb{F}$  we inductively define  $\Delta_{a_1, \ldots, a_k}(f) :=$  $\Delta_{a_1}(\Delta_{a_2,...,a_k}(f)).$ 

- 1. Let  $a_1, \ldots, a_k \in \mathbb{F}$  be a basis for H over  $\mathbb{F}_p$ . Prove that  $\Delta_{a_1, \ldots, a_k}(f)(a_0) = \sum_{a \in H} f(a_0 + a)$ .
- 2. Write the p-ary expansion of an integer  $c \in \mathbb{N}$  as  $\sum_{i\geq 0} c_i p^i$  for  $0 \leq c_i < p$ . Define the sum of the "p-ary digits" of c as  $ds_p(c) \coloneqq \sum_{i \geq 0} c_i$ . For a polynomial  $g(X) = \sum_{j \geq 0} \alpha_j X^j$ , define  $ds_p(g) := \max(\{ds_p(j) : \alpha_j \neq 0\} \cup \{-1\}).$  Prove that for every  $a \in \mathbb{F}$  it holds that

$$
\mathsf{ds}_p(\Delta_a(g)) \le \max\{\mathsf{ds}_p(g) - (p-1), -1\} .
$$

You may use the following facts:

- (a)  $b^p = b$  for all  $b \in \mathbb{F}_p$ .
- (b) For  $c, d \in \mathbb{N}$ , let  $c = \sum_i c_i p^i$  and  $d = \sum_i d_i p^i$  be their p-ary expansions. If there exists  $i \geq 0$  such that  $d_i > c_i$  then  $\binom{c_i}{d_i}$  $\binom{c}{d} \equiv 0 \pmod{p}.$

Hint: consider a single monomial  $X<sup>c</sup>$ , and apply the binomial theorem.

3. Prove that for every univariate polynomial  $g \in \mathbb{F}[X]$  with  $deg(g) < |H|$ , letting  $\beta$  be the coefficient of  $X^{|H|-1}$  in g, it holds that

$$
\sum_{a \in H} g(a) = \beta \cdot \sum_{a \in H} a^{|H|-1}
$$

.

**Problem 3. (Bonus)** Show that  $\sum_{a \in H} a^{|H|-1} = \prod_{a \in H \setminus \{0\}} a$  $\sum_{a \in H} a^{|H|-1} = \prod_{a \in H \setminus \{0\}} a$  $\sum_{a \in H} a^{|H|-1} = \prod_{a \in H \setminus \{0\}} a$ . (*Hint: use Newton's identities.*<sup>1</sup>) The right-hand side is equal to the coefficient of the linear term of  $v_H(X)$ , which can be computed in time  $O(\log^2|H|)$ .

<span id="page-1-0"></span> $^1$ [https://en.wikipedia.org/wiki/Newton's\\_identities](https://en.wikipedia.org/wiki/Newton)