

# **The Role of Side Information in Communication and Data Compression**

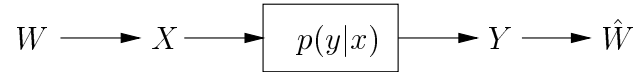
Tom Cover  
Stanford University

Combination of joint work with Mung Chiang, Arak Sutivong, Young-Han Kim, and David Julian.

# Introduction

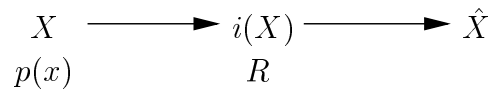
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## Channel



$$C = \max_{p(x)} I(X; Y)$$

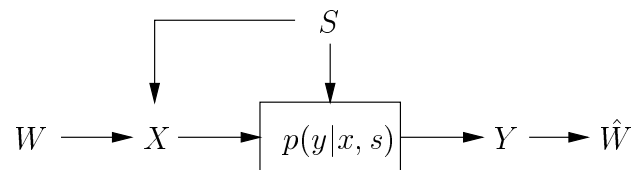
## Data Compression



$$R = \min_{p(\hat{x}|x)} I(X; \hat{X})$$

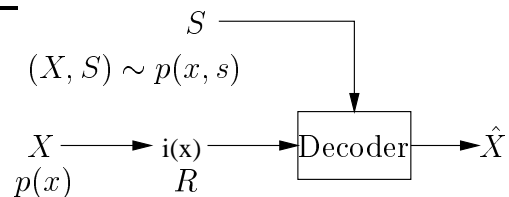
Now consider state information:

## Channel



$$C \stackrel{??}{=} \max_{p(x|s)} I(X; Y|S)$$

## Data Compression

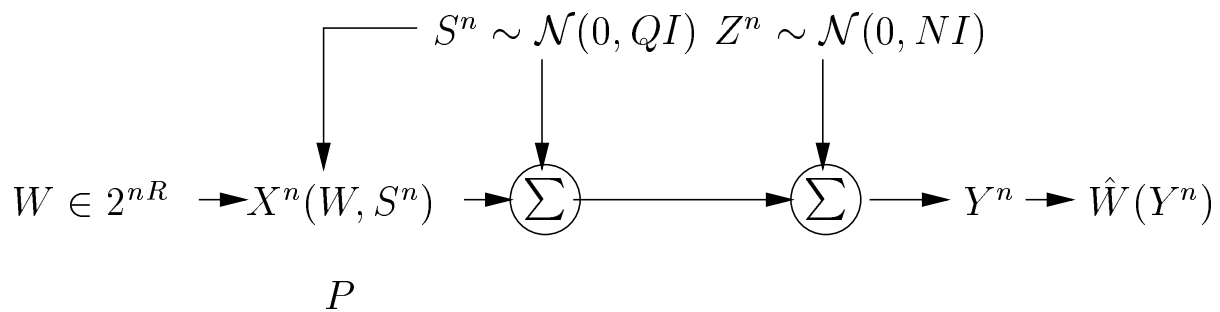


$$R \stackrel{??}{=} \min_{p(\hat{x}|x, s)} I(X; \hat{X}|S)$$

# Example 1

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## Gaussian



$$C = \frac{1}{2} \log\left(1 + \frac{P}{N}\right)$$

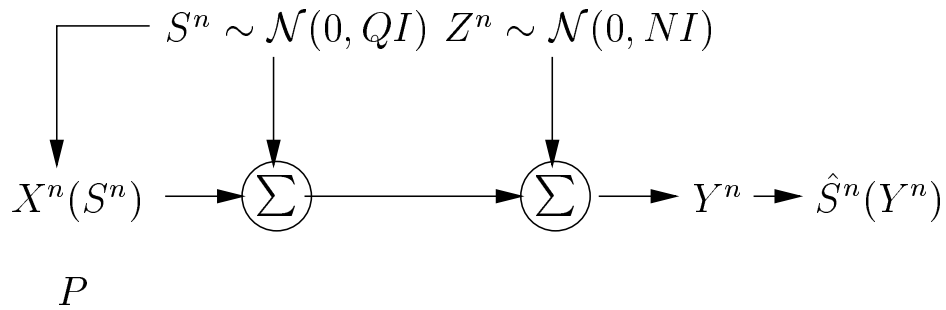
(Costa, Writing on Dirty Paper)

$$E(S - \hat{S})^2 = Q \frac{P + N}{Q + P + N}$$

## Example 2

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### State estimation



Let  $X = \sqrt{\frac{P}{Q}}S$  (amplifying state).

Then

$$\begin{aligned}
 Y &= X + S + Z \\
 &= \left( 1 + \sqrt{\frac{P}{Q}} \right) S + Z.
 \end{aligned}$$

$$E(S - \hat{S})^2 = Q \frac{N}{\left( \sqrt{Q} + \sqrt{P} \right)^2 + N}$$

## Example 3

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### 5-card channel

- 52 card deck
- Deal 5 cards
- Hold one back
- View other 4
- Infer the missing card

### Example

$$\begin{array}{cccccc} & 5C & 2H & 5H & JH & 10S \\ i = & 5 & 28 & 31 & 37 & 49 \end{array}$$

Hold out: 2H

Send: 5C 5H JH 10S

Infer missing card: 2H

### Conflict of state with info

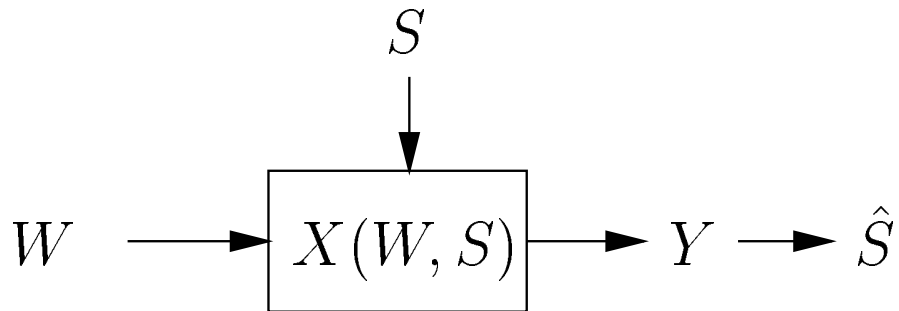
What you say conflicts with how you say it.

It's what you don't do that conveys the information.

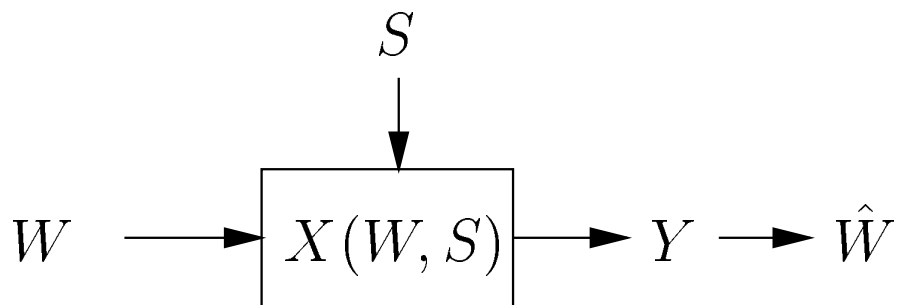
## Channel with known state

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$$\begin{array}{ccccc}
 S & \longrightarrow & X = Y & \longrightarrow & \hat{S} \\
 5 \text{ cards} & & 4 \text{ cards} & & 5 \text{ cards}
 \end{array}$$



Pure state info



Pure info

Conflict of  $\hat{S}$  and  $\hat{W}$ .

## 5-card channel: pure information

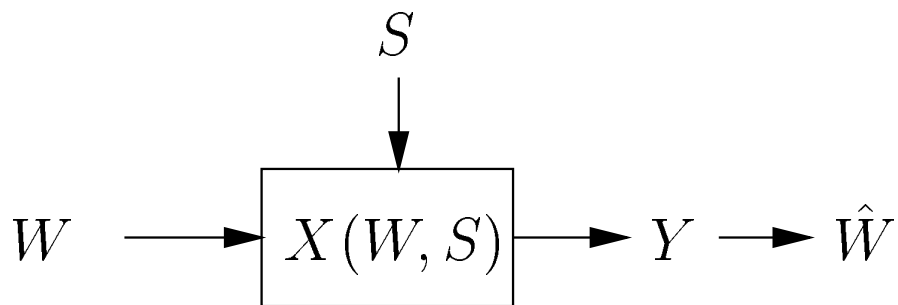
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$m = 52$ -card deck

$S = 5$  cards

$X = 4$ -card subset

$Y = X$



Claim Capacity  $C = \log 5$  bits

Pinsker, Gel'fand, El Gamal, Heegard  $C = I(U; Y) - I(U; S)$ .

## 5-card channel: state estimation

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How large a deck  $m$ ?

$C = \log 5$  bits

$4!$  orderings of 4 cards

(5-card hand)  $\rightarrow (5)4! = 120$  actions

Can determine 120 missing cards.

$$m = 124$$

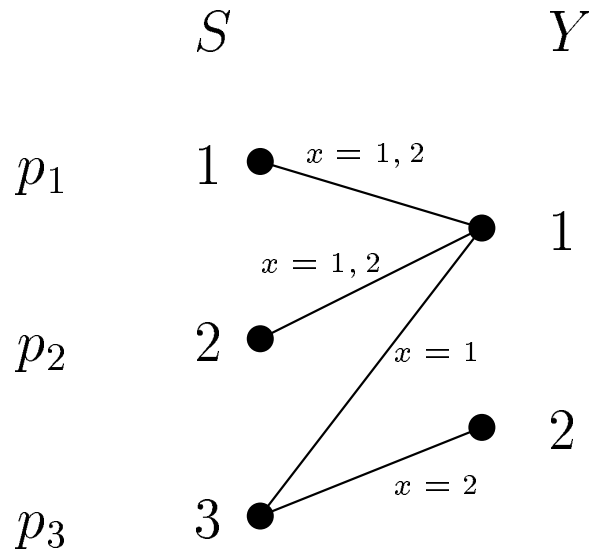


## Example 4

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### Deterministic channel

$$y = g(x,s)$$



### Pure info

$$C_{10} = E_S \log |\mathcal{Y}_S| = \max_{p(x|s)} H(Y|S) = p_3$$

$$\alpha^* = \frac{1}{2}$$

### State estimation Decrease in entropy of $S$ :

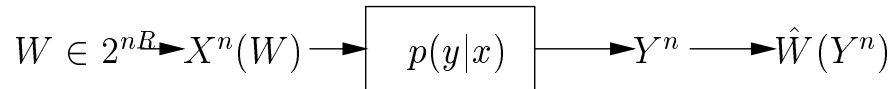
$$\Delta H(S) = H(S) - (H(S|Y) - H(Y|S)) = H(Y)$$

$$\max_{p(x|s)} H(Y) = \max_{\alpha} H(\alpha p_3) = \begin{cases} 1, & p_3 \geq \frac{1}{2} \\ H(p_3), & p_3 \leq \frac{1}{2} \end{cases}$$

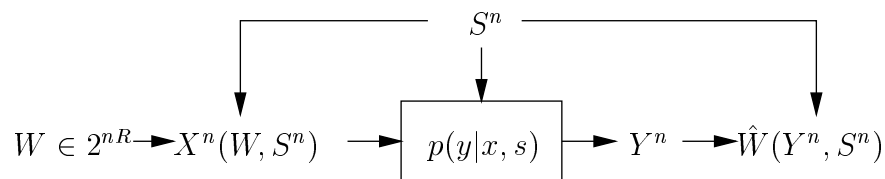
$$\alpha^* = \frac{1}{2p_3}$$

# Channel capacity with side information

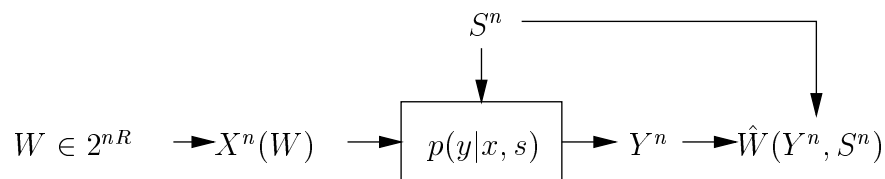
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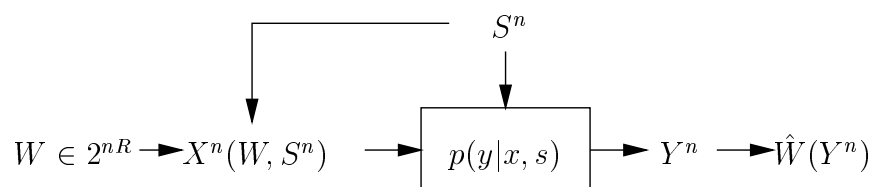
$$C_{00} = \max_{p(x)} I(X; Y)$$



$$C_{11} = \max_{p(x|s)} I(X; Y | S)$$



$$C_{01} = \max_{p(x)} I(X; Y | S)$$

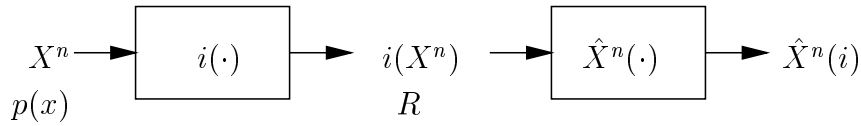


$$C_{10} = \max_{p(u, x|s)} (I(U; Y) - I(U; S))$$

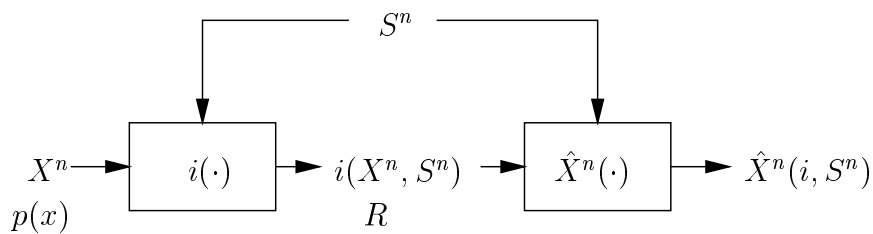
Gel'fand, Pinsker, El Gamal, and Heegard

# Rate-distortion with side information

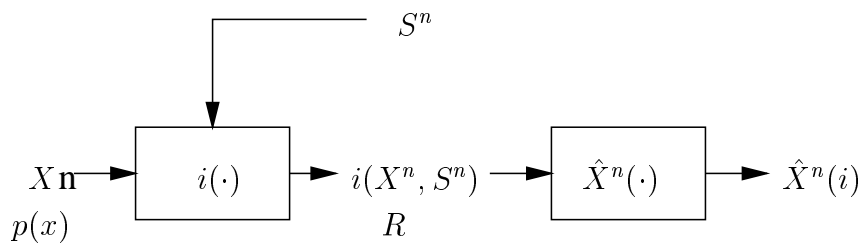
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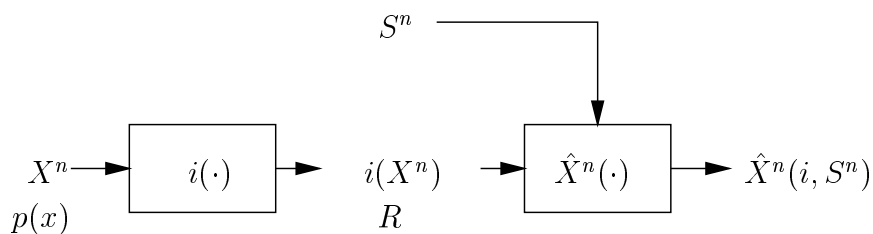
$$R_{00} = \min_{p(\hat{x}|x)} I(X; \hat{X})$$



$$R_{11} = \min_{p(\hat{x}|x,s)} I(X; \hat{X} | S)$$



$$R_{10} = \min_{p(\hat{x}|x)} I(X; \hat{X})$$

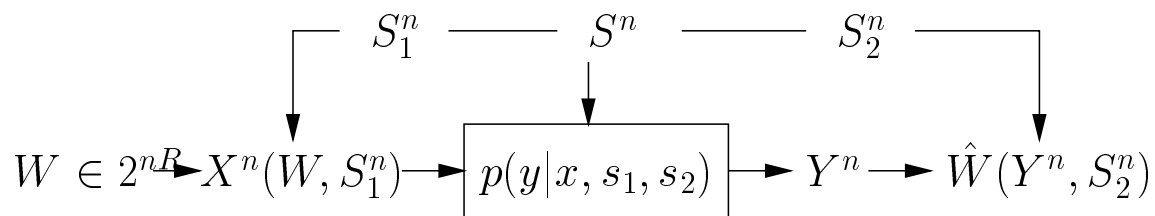


$$R_{01} = \min_{p(u|x)p(\hat{x}|u,s)} (I(U; X) - I(U; S))$$

Wyner and Ziv

# Generalized capacity with state information

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Channel:  $p(y|x, s_1, s_2)$

State:  $(S_{1i}, S_{2i}) \sim \text{i.i.d. } p(s_1, s_2)$

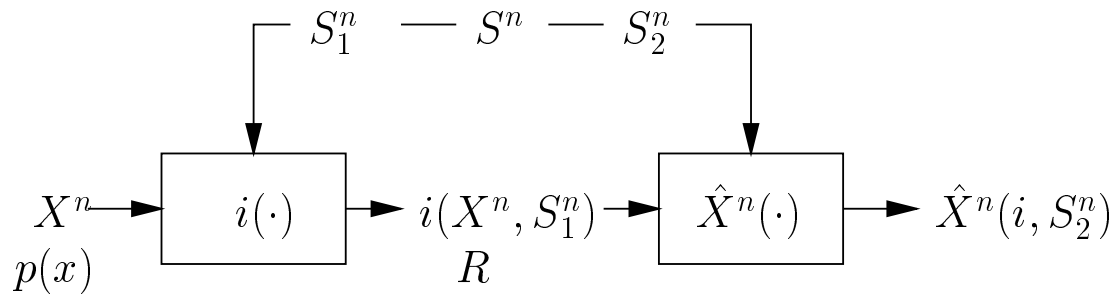
## Theorem 1\*

$$C = \max_{p(u,x|s_1)} (I(U; Y, S_2) - I(U; S_1))$$

\* M. Chiang

# Generalized rate-distortion with side information

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Source:  $(X_i, S_{1i}, S_{2i}) \sim \text{i.i.d. } p(x, s_1, s_2)$

Distortion:  $d(x, \hat{x})$

## Theorem 2

$$R = \min_{p(u|x, s_1)p(\hat{x}|u, s_2)} (I(U; X, S_1) - I(U; S_2))$$

# Duality

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## Channel capacity

$$C_{S_1 S_2} = \max_{p(u|s_1)p(x|u,s_1)} (I(U; Y, S_2) - I(U; S_1))$$

## Rate-distortion function

$$R_{S_1 S_2}(D) = \min_{p(u|x,s_1)p(\hat{x}|u,s_2): Ed(X, \hat{X}) \leq D} (I(U; X, S_1) - I(U; S_2))$$

Duality:

- Isomorphism
- Complementarity

$$C \Leftrightarrow R(D)$$

maximization  $\Leftrightarrow$  minimization

$$C_{00} \Leftrightarrow R_{00}$$

$$C_{11} \Leftrightarrow R_{11}$$

$$C_{10} \Leftrightarrow R_{01}$$

Y (received symbol)  $\Leftrightarrow$  X (source)

S (state)  $\Leftrightarrow$  S (state)

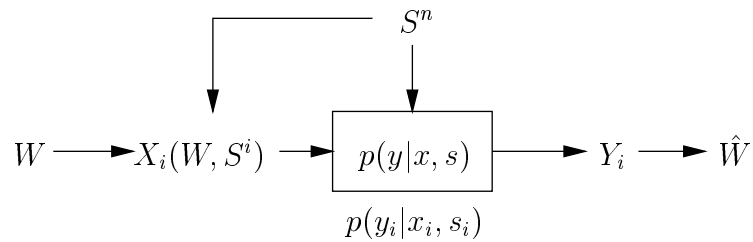
U (auxiliary)  $\Leftrightarrow$  U (auxiliary)

X (transmitted symbol)  $\Leftrightarrow$   $\hat{X}$  (estimation)

Note: Pradhan, Chou, and Ramchandran have some recent results on functional duality between  $C$  and  $R$ .

## Causal side information (Shannon)

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$$C = \max_{p(t)} I(T; Y), \quad T \in \{f : \mathcal{S} \rightarrow \mathcal{X}\}$$

Can recast as

$$C = \max_{p(x(\cdot))} I(X(\cdot); Y)$$

(Shannon, Channels with state information at the transmitter, *IBM J. Res. Dev.*, 1958, pp 289-293.)

## New characterization

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### Shannon (Causal)

$$C_{10} = \max_{p(x(\cdot))} I(X(\cdot); Y)$$

Can recast (YHK) as

$$C_{10} = \max_{\substack{p(u) p(x|u, s) \\ x = f(u, s)}} (I(U; Y) - I(U; S))$$

Generalized capacity with two-sided state information:

### Noncausal

$$C = \max_{p(u|s_1) p(x|u, s_1)} (I(U; Y, S_2) - I(U; S_1))$$

### Causal

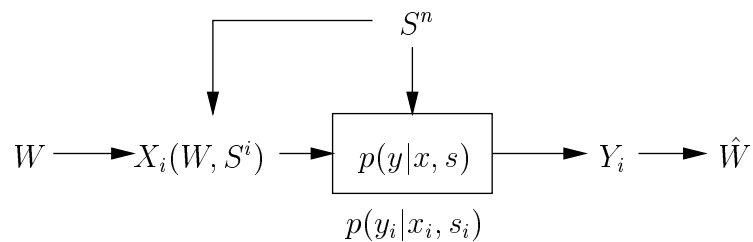
$$C = \max_{p(u) p(x|u, s_1)} (I(U; Y, S_2) - I(U; S_1))$$

(Young-Han Kim)



## Converse for Shannon causal side information channel

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$$\begin{aligned}
 C_{10} &= \max_{p(u)p(x|u,s)} (I(U; Y) - I(U; S)) \\
 &= \max_{p(u)p(x|u,s)} I(U; Y)
 \end{aligned}$$

Converse:

Given code  $X_i(W, S^i)$ , with  $\Pr\{\hat{W} \neq W\} \rightarrow 0$ .

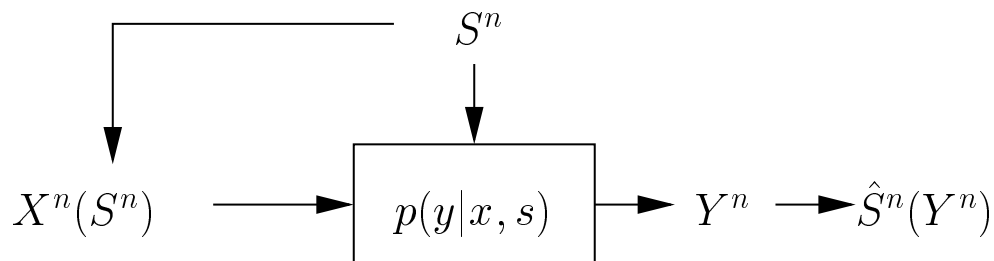
Then

$$\begin{aligned}
 nR &= H(W) \\
 &\leq I(W; Y^n) + n\epsilon_n \\
 &= H(Y^n) - H(Y^n|W) + n\epsilon_n \\
 &\leq \sum_{i=1}^n H(Y_i) - H(Y_i | \underbrace{W, Y^{i-1}}_{U_i}) + n\epsilon_n \\
 &= \sum_{i=1}^n I(U_i; Y_i) + n\epsilon_n \\
 &\leq nC + n\epsilon.
 \end{aligned}$$

## Sending state information

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Deterministic channel. “Amplify” state information.



Let  $y = g(x, s)$  (i.e., a deterministic function).

Note: Capacity  $C = \max_{p(x|s)} H(Y|S)$  bits

Reduction in uncertainty about  $S$ :

Original uncertainty:  $H(S)$

Uncertainty after observing  $Y$ :  $H(S|Y)$

Uncertainty after refinement (using  $C$ ):  $H(S|Y) - H(Y|S)$

Note: Reduction in uncertainty  $\Delta H(S) = \max_{p(x|s)} H(Y)$

(Arak Sutivong)

# Concluding remarks

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## Channel with noncausal state information

$$C_{S_1 S_2} = \max_{p(u|s_1)p(x|u,s_1)} (I(U; Y, S_2) - I(U; S_1))$$

## Rate-distortion function

$$R_{S_1 S_2} = \min_{p(u|x,s_1)p(\hat{x}|u,s_2)} (I(U; X, S_1) - I(U; S_2))$$

## Channel with causal state information

$$C_{S_1 S_2} = \max_{p(u)p(x|u,s_1)} (I(U; Y, S_2) - I(U; S_1))$$

## Amplification of state

$$\Delta H(S) = \max_{p(x|s)} H(Y)$$