

Probabilistic tools in nonlinear Dvoretzky Theory (Assaf Naor)

Dvoretzky's thm

$\forall k \in \mathbb{N}, \forall D > 1, \exists n = n(k, D) \text{ s.t. } \forall n\text{-dim}$
 normed space $X \exists \text{ linear subspace } Y \text{ s.t.}$

- 1) $\dim Y = k$
- 2) $C_2(Y) \leq D$

where $C_2(M) = \underset{\substack{\uparrow \\ (M, d) \text{ metric space}}}{\text{smallest } D \text{ s.t. } \exists f: M \rightarrow \ell_2 \text{ with}}$
 $d(x, y) \leq \|f(x) - f(y)\| \leq Dd(x, y) \quad \forall x, y \in M$

(Later, we will also discuss $C_1(M)$, the analogous
 thing for ℓ_1).

Bourgain-Figiel-Milman (1986)

$\forall n \in \mathbb{N}, \forall D > 1$, what is the largest $m = m(n, D) \in \mathbb{N}$
 such that any n -point metric space $(X, d) \exists Y \subseteq X$ s.t.
 1) $|Y| \geq m$
 2) $C_2(Y) \leq D$?

Tao 2006

$\forall \alpha > 0 \quad \forall D > 1$, what is the largest $\beta = \beta(\alpha, D)$
 such that in any compact metric space $X \exists$ closed
 $Y \subseteq X$ s.t. $\dim_H(Y) \geq \alpha$

- 1) $\dim_H Y \geq \beta$
- 2) $C_2(Y) \leq D$?

Mendel, N 2011 : Answer to Tao's problem

2012? : Some more extensions.

2006: Answer to BFM: $\forall \varepsilon \in (0, 1) \quad \forall n\text{-point } (X, d)$
 $\exists Y \subseteq X$ s.t. 1) $|Y| \geq n^{1-\varepsilon}$
 2) $C_2(Y) \leq \frac{100}{\varepsilon}$.

Bartal, Linial, Mendel, N (2002):

The estimate is sharp: $\forall n \exists$ n -point (X_n, d)
s.t. if $Y \subseteq X_n$, $|Y| \geq n^{1-\varepsilon}$ then $C_2(Y) \geq c/\varepsilon$.

N, Tao (2010): $D > 2$. Let Θ be the solution
to $\frac{2}{D} = (1-\Theta)^{\Theta/(1-\Theta)}$. Then any n -point
metric space X has $Y \subseteq X$ s.t.
1) $|Y| \geq n^\Theta$
2) $C_2(Y) \leq D$.

$$\begin{aligned}\text{Note: } D \rightarrow \infty &\Rightarrow \Theta \sim 1 - \frac{2\varepsilon}{D} \\ D = 2 + \varepsilon &\Rightarrow \Theta \sim \varepsilon / \log(1/\varepsilon)\end{aligned}$$

This Θ is the best for the method.

Majorizing measures then is a cor of (2012?).

N-Tao 2010: application to Hardy-Littlewood
inequality.

How to prove lower bounds of the BLMN type?

Consider $(\{0,1\}^n, \|\cdot\|_1)$.

Enflo 69: $C_2(\{0,1\}^n) = \sqrt{n}$.

Fourier-analytic gives the same bound.

What about $S \subseteq \{0,1\}^n$, $|S|$ big?

BLMN: $C_2(S) \geq \sqrt{n / \log 2^n / |S|}$.

(So $|S| = 2^{n(1-\varepsilon)} \Rightarrow C_2(S) \geq \min \left\{ \frac{1}{\sqrt{\varepsilon}}, \sqrt{n} \right\}$.)

$\exists S \subseteq \{0,1\}^n$, $|S| \geq 2^{n(1-\varepsilon)}$ and
 $C_2(S) \leq \sqrt{\frac{\log \frac{1}{\varepsilon}}{\varepsilon}}$

We will do a simpler case that contains all of the ideas:
 let G be an n -vertex k -regular graph with girth g .
 (i.e. balls of radius $\leq g/4$ are isometric to a tree).

Bourgain: If T_n is the complete binary tree of depth n then $C_2(T_n) \approx \sqrt{\log n}$
 $\Rightarrow C_2(G) \geq \sqrt{\log g}$.

Linial, Mayenz, N (2001): $C_2(G) \geq \sqrt{g}$ using Markov type,
 because random walks diverge at linear speed.

④ Thm: $S \subseteq G \Rightarrow C_2(S) \geq \sqrt{\frac{g}{1 + \log(\frac{n}{|S|})}}$.

Open: are there graphs with girth $\rightarrow \infty$ s.t. $C_2(G) = O(\sqrt{g})$

Ostrovskii (2011): $\exists G_n$ with girth $\rightarrow \infty$ s.t. $C_1(G_n) = O(1)$.

These are related to the question of whether graphs with large girth contain expanders.

Proof of ④: If G has average degree k & girth g then $C_2(G) \geq (1 - 2/k) \sqrt{g}$.



Want to turn S into something with many edges.

General fact: If H is a graph on $\{1, \dots, n\}$.

$$d = d_1 \geq d_2 \geq \dots \geq d_n < 0$$

$$\text{then } \forall S \subseteq H, \quad 2E(S) \geq d|S|^2/n + d_n|S|$$

large girth tells us something about this.

Let $G^{(m)}$ be the distance- m graph, $m < \frac{g}{2}$.
 $A_{G^{(m)}}$ = adjacency mat.

$$A_{G^{(m)}} = P_m^k (A_G)$$

↑ Geronimus polynomials

$$P_0^k(x) = 1$$

$$P_1^k(x) = x$$

$$P_2^k(x) = x^2 - k$$

$$P_m^k(x) = x P_{m-1}^k(x) - (k-1) P_{m-2}^k(x)$$

$$\lambda_n(G^{(m)}) \geq \min_{x \in \mathbb{R}} P_m^k(x)$$

↑ very crude!

$$P_m^k(2\sqrt{k+1} \cos \theta) = (k-1)^{\frac{m}{2}-1} \frac{(k-1) \sin((m+1)\theta) \sin((n+1)\theta)}{\sin \theta}$$

~~break~~

Now do a random walk on $G^{(m)} \cap S$ of length t , with $mt < \frac{g}{2}$, and optimize the parameters. \square .