

Jesús De Loera

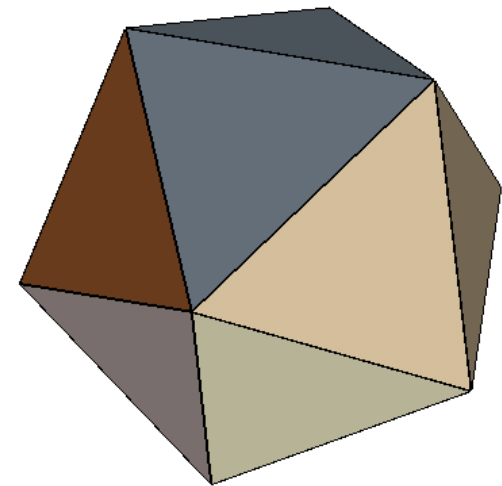
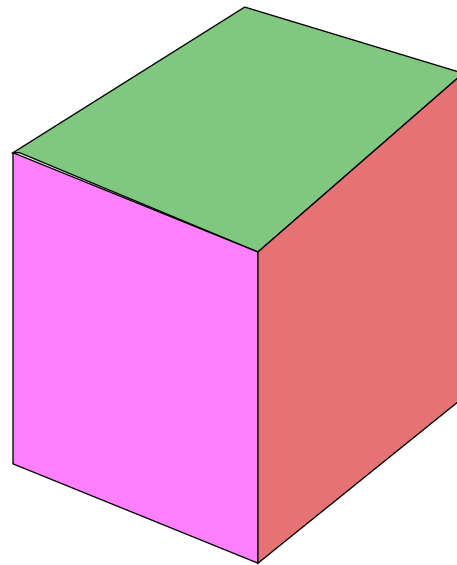
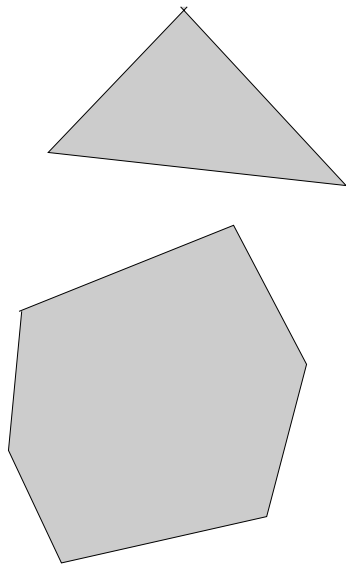
# Easy-to-State but Hard to Solve: Challenges in Polyhedral Computation

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# What is a Convex Polytope?

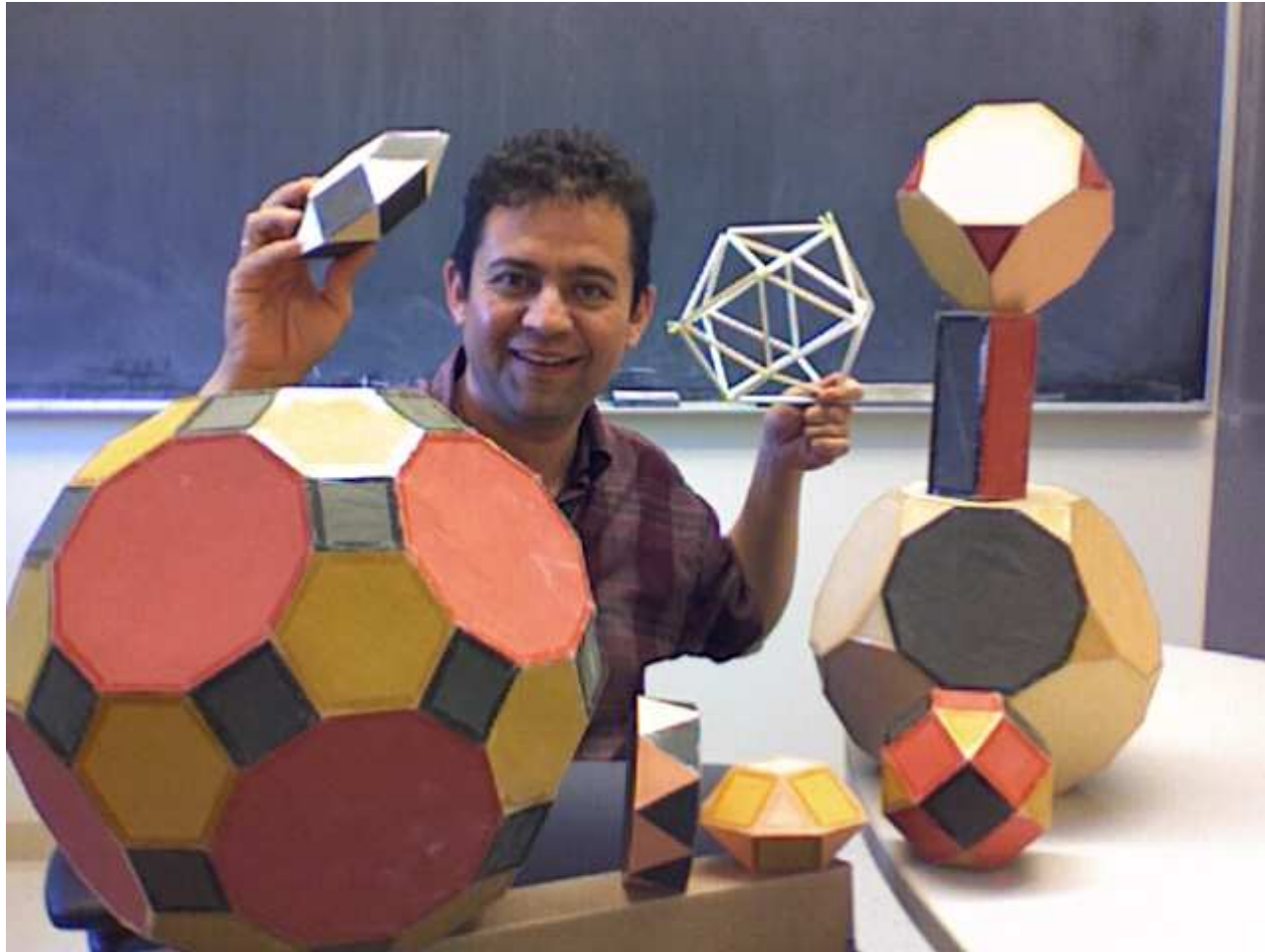
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**Well, something like these...**



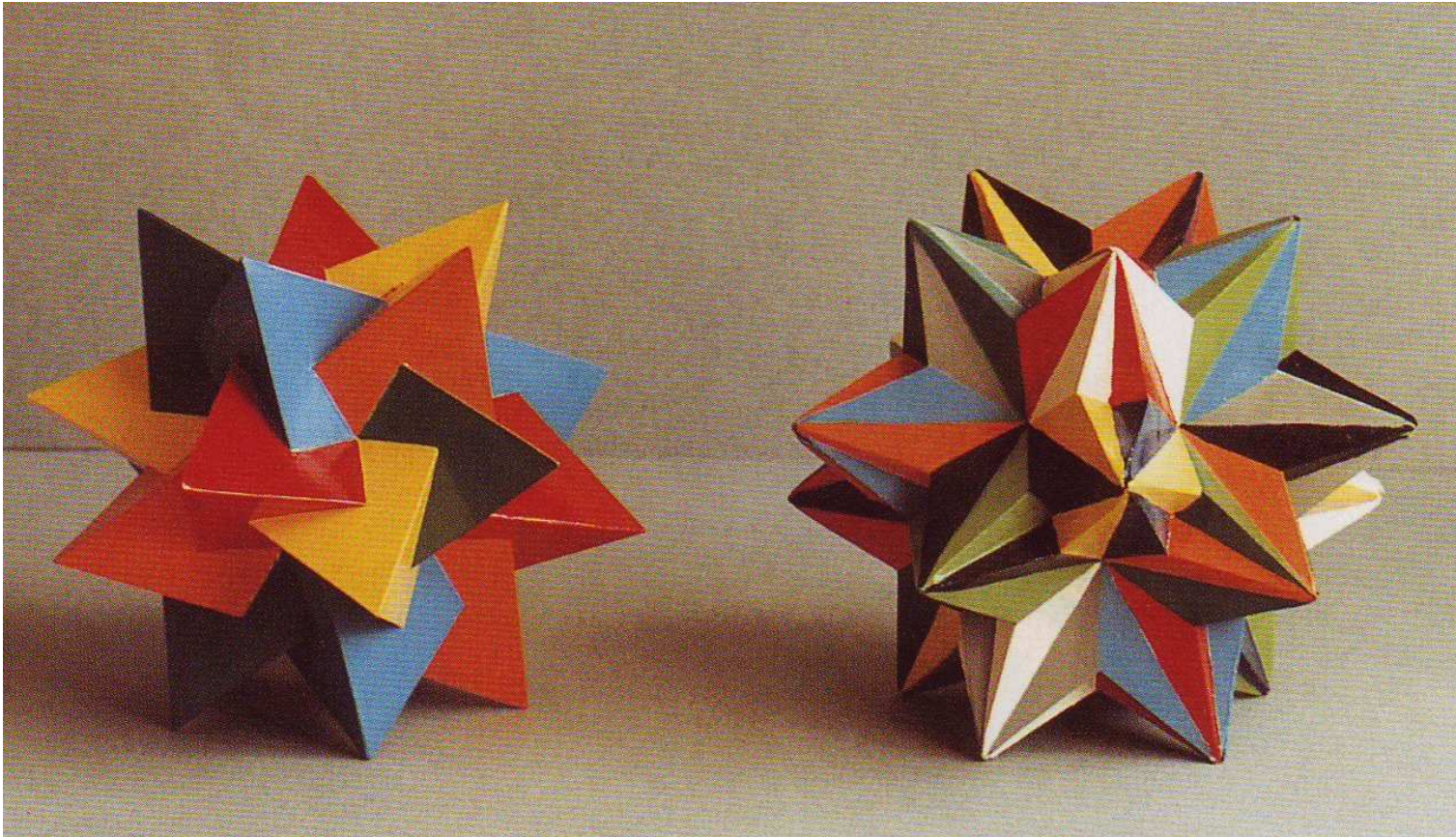
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or like these



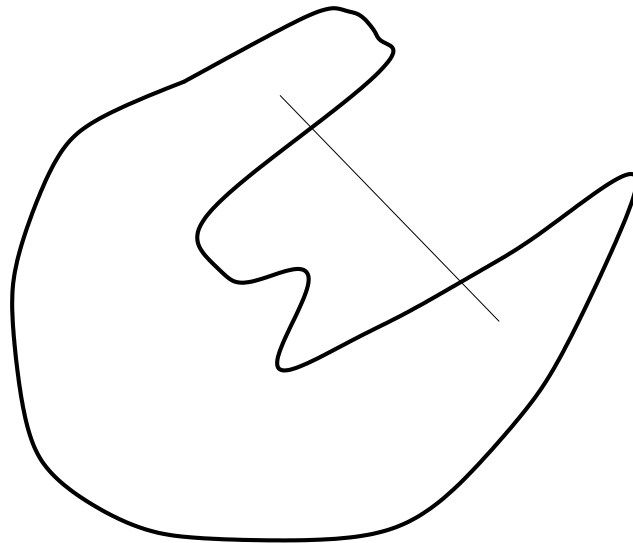
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**But NOT quite like these!**

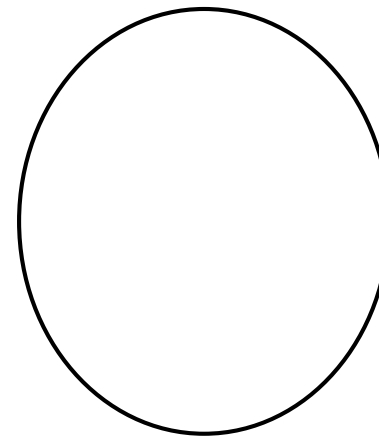


## A definition PLEASE!

The word **CONVEX** stands for sets that contain any line segment joining two of its points:



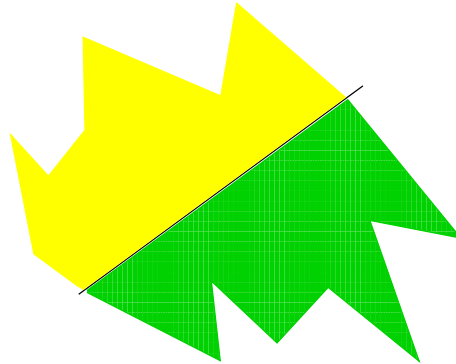
NOT CONVEX



CONVEX

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A (hyper)plane divides spaces into two *half-spaces*. Half-spaces are convex sets! Intersection of convex sets is a convex set!



Formally a half-space is a *linear inequality*:

$$a_1x_1 + a_2x_2 + \dots + a_dx_d \leq b$$

**Definition:** A **polytope** is a bounded subset of Euclidean space that results as the intersection of finitely many half-spaces.

## An algebraic formulation for polytopes

A polytope has also an algebraic representation as the set of solutions of a system of linear inequalities:

$$a_{1,1}x_1 + a_{1,2}x_2 + \dots + a_{1,d}x_d \leq b_1$$

$$a_{2,1}x_1 + a_{2,2}x_2 + \dots + a_{2,d}x_d \leq b_2$$

⋮

$$a_{k,1}x_1 + a_{k,2}x_2 + \dots + a_{k,d}x_d \leq b_k$$

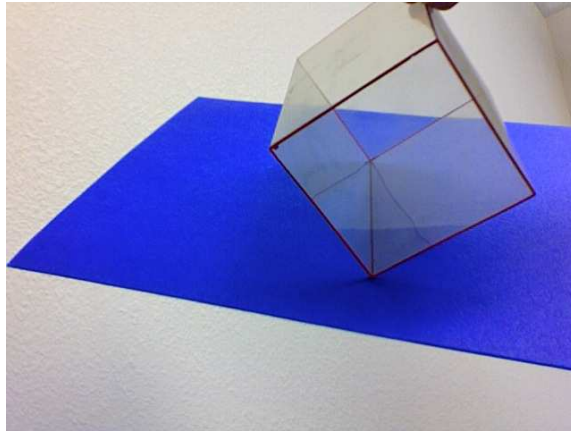
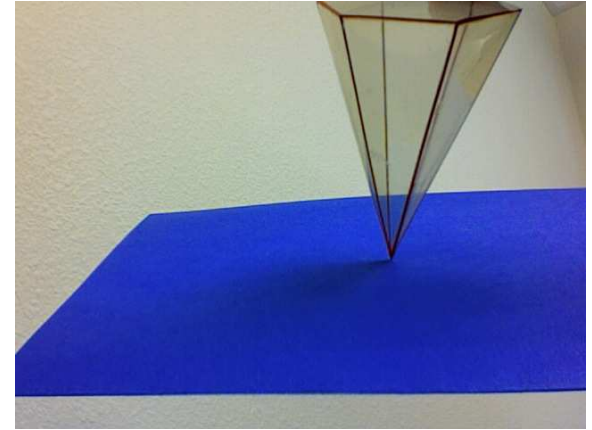
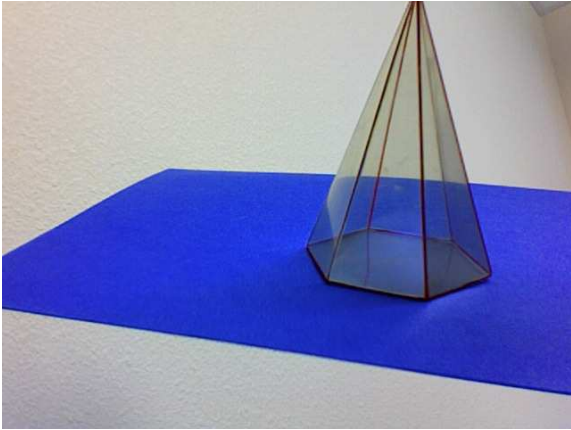
**Note:** This includes the possibility of using some linear equalities as well as inequalities!!



# Problem 1:

# NUMBERS OF FACES

# Faces of Polytopes



## Some Numeric Properties of Polyhedra



- **Euler's formula**  $V - E + F = 2$ , relates the number of vertices  $V$ , edges  $E$ , and facets  $F$  of a 3-dimensional polytope.

Given a convex 3-polytope  $P$ , if  $f_i(P)$  the number of  $i$ -dimensional faces. There is one vector  $(f_0(P), f_1(P), f_2(P))$ . that counts faces, the  **$f$ -vector** of  $P$ .

- **Theorem** (Steinitz 1906) A vector of non-negative integers  $(f_0(P), f_1(P), f_2(P)) \in \mathbb{Z}^3$  is a the  $f$ -vector of a 3-dimensional polytope if and only if

1.  $f_0(P) - f_1(P) + f_2(P) = 2$
2.  $2f_1(P) \geq 3f_0(P)$
3.  $2f_1(P) \geq 3f_2(P)$

- **OPEN PROBLEM 1:** Can one find similar conditions characterizing  $f$ -vectors of 4-dimensional polytopes?

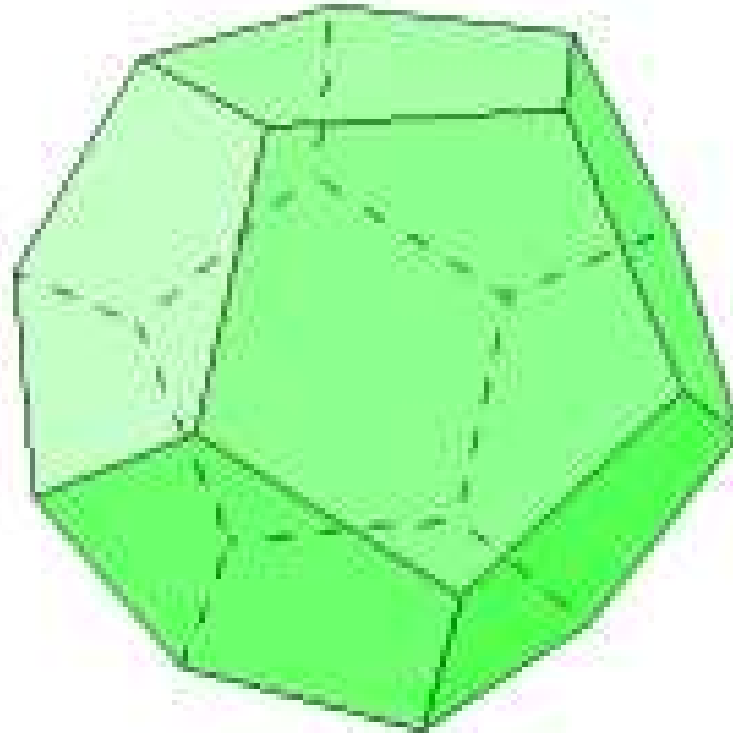
In this case the vectors have 4 components  $(f_0, f_1, f_2, f_3)$ .

## Problem 2:

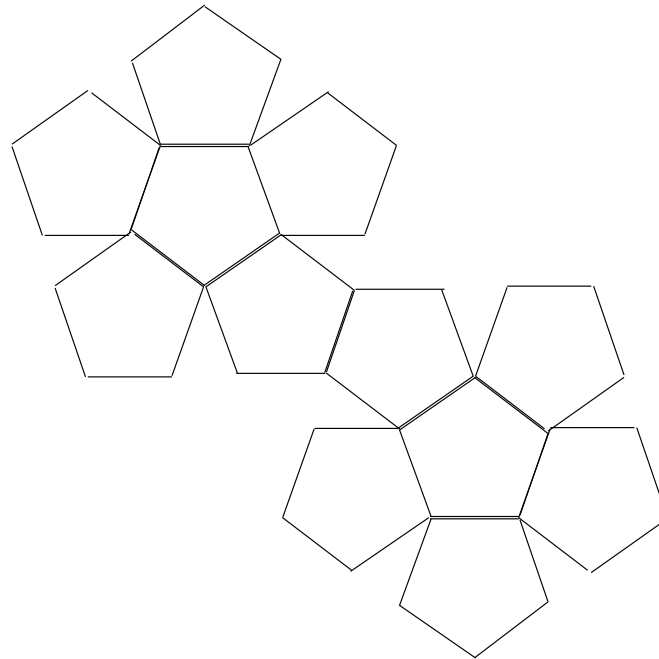
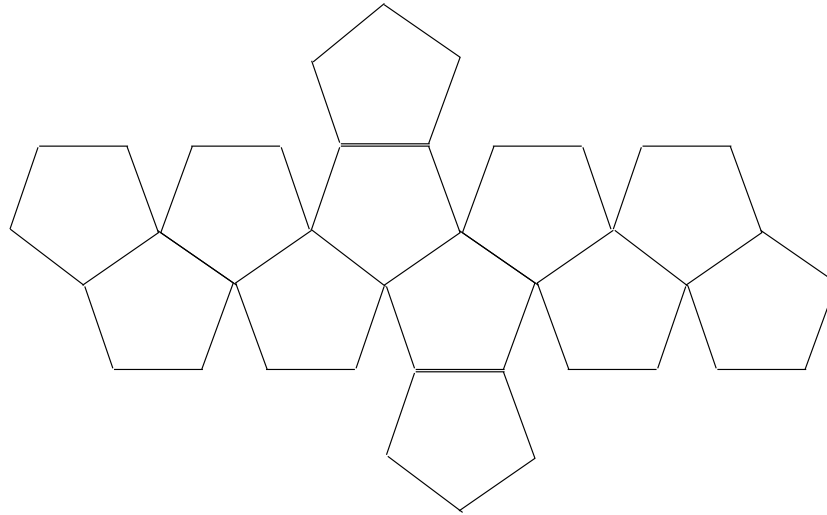
# UNFOLDING POLYTOPES

## Unfolding Polyhedra

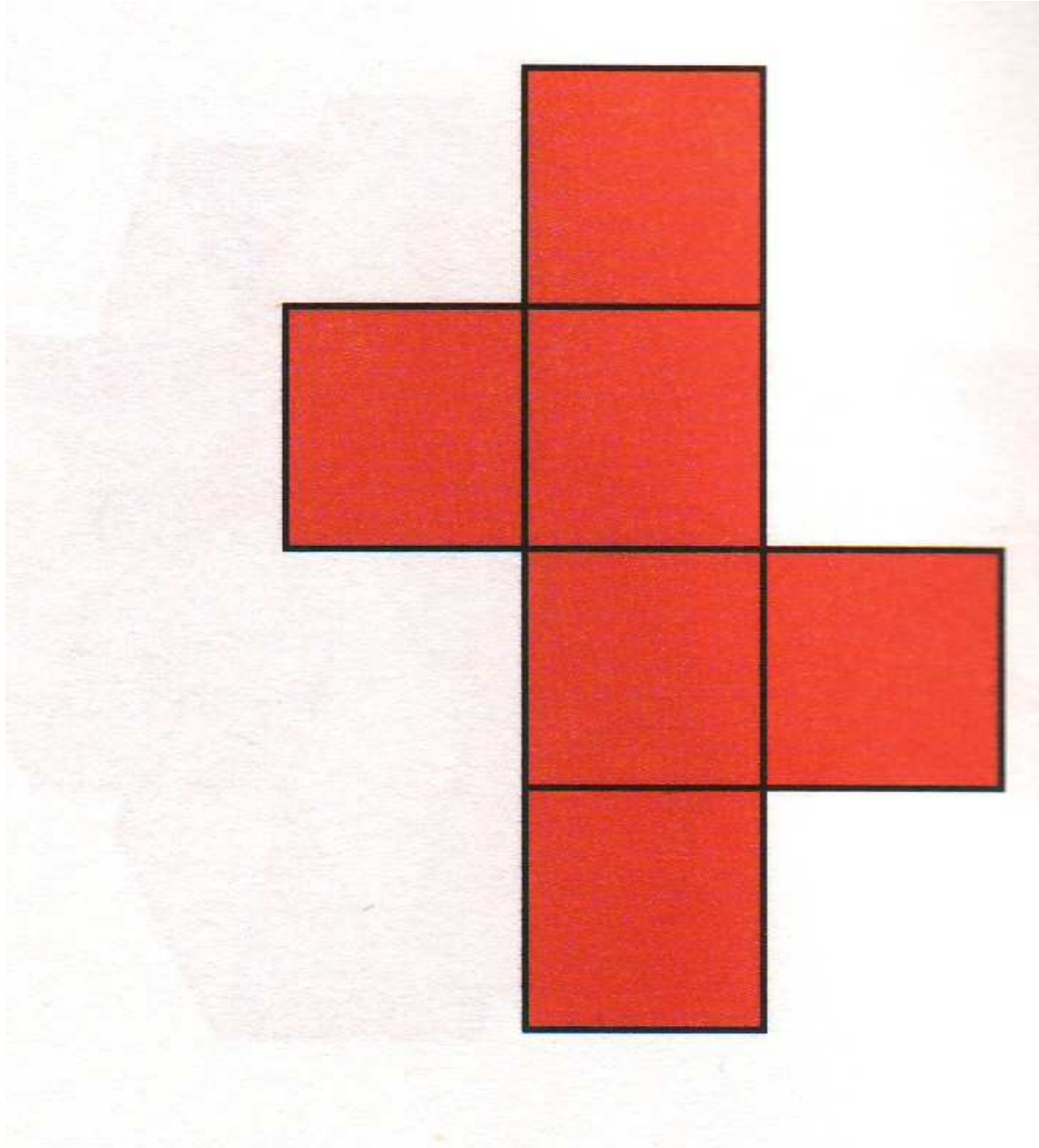
What happens if we use scissors and cut along the edges of a polyhedron?  
What happens to a dodecahedron?



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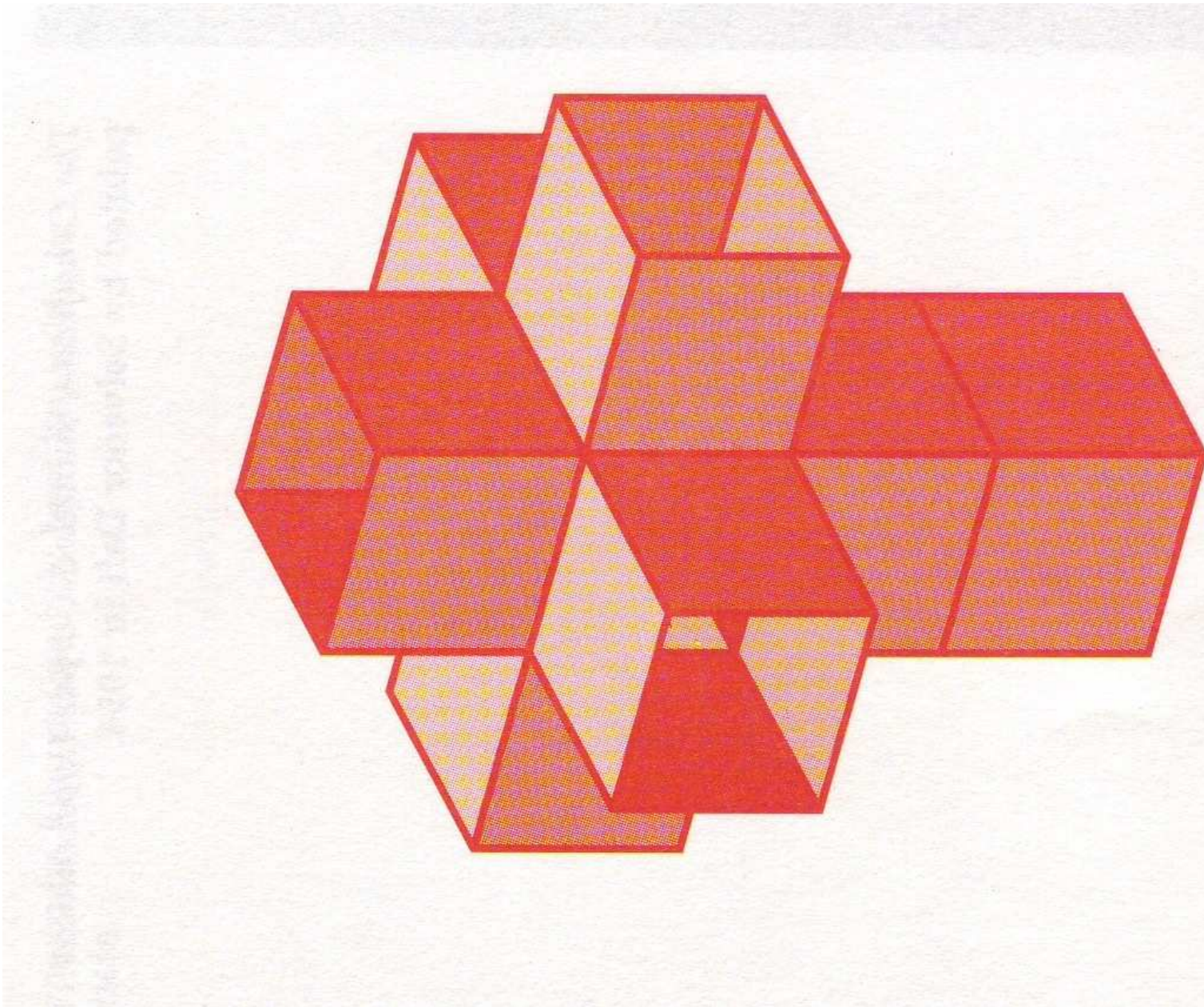


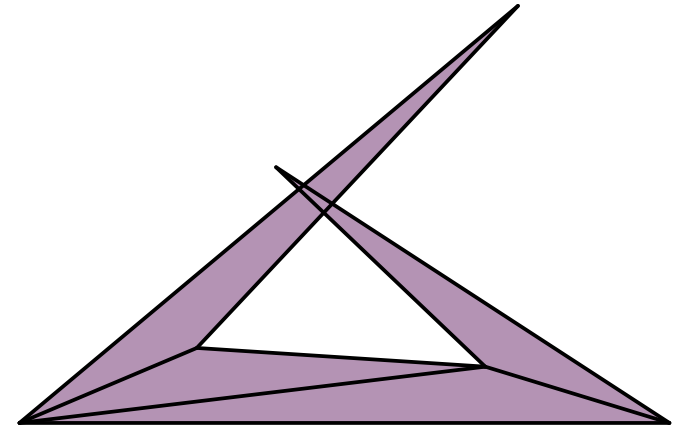
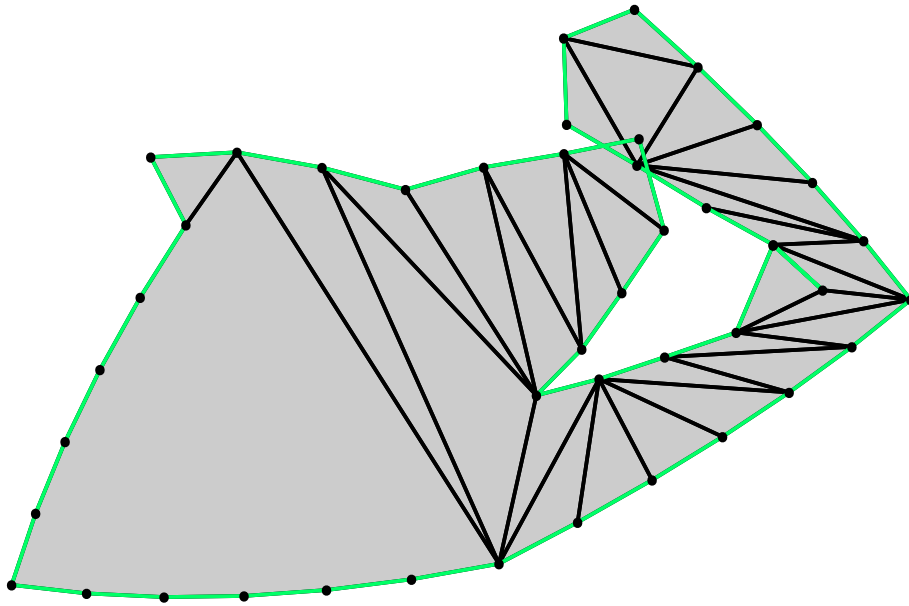
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**Open Problem 2:** Can one always find an unfolding that has no self-overlappings for all 3-polytopes?

# Problem 3:

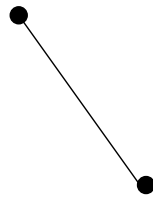
# TRIANGULATIONS

# Simplices

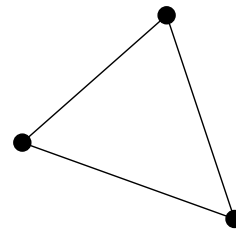
A **simplex** is any polytope of dimension  $d$  with  $d + 1$  vertices.



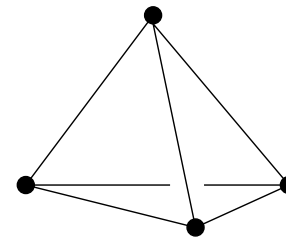
0-simplex



1-simplex



2-simplex



3-simplex

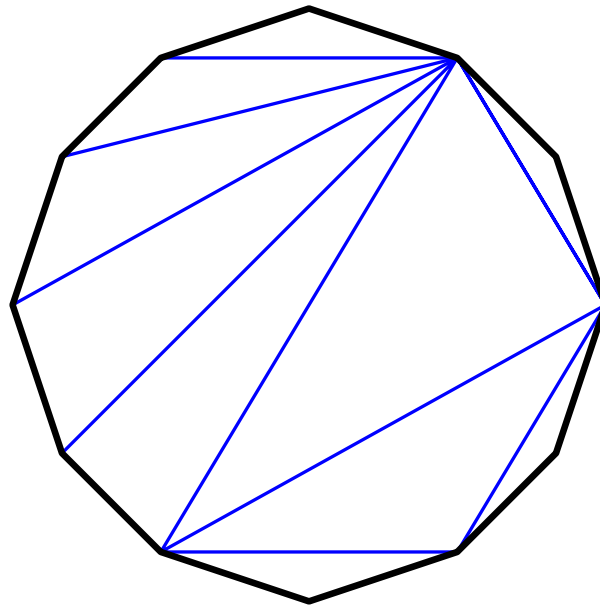
A  $d$ -simplex has exactly  $\binom{d+1}{i+1}$  faces of dimension  $i$ , ( $i = -1, 0, \dots, d$ ), which are themselves  $i$ -simplices.

**IMPORTANT:** Every polytope can be decomposed as a union of simplices.

## Triangulations

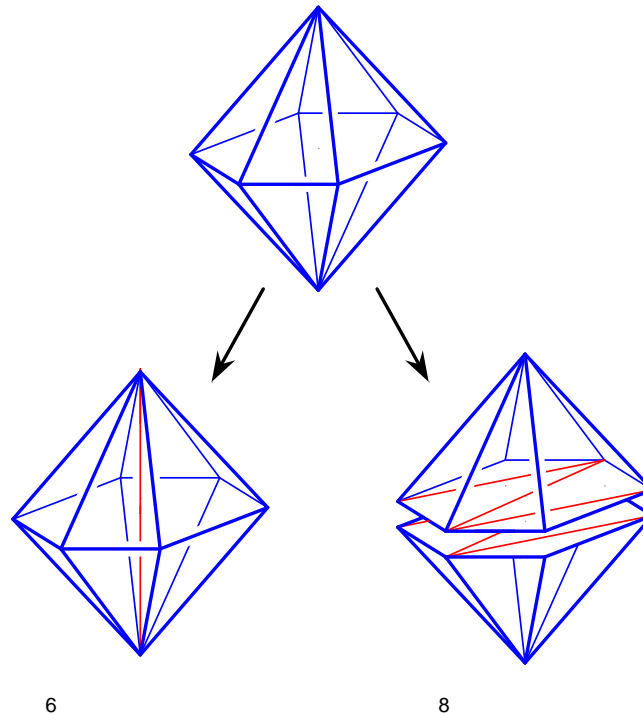
A **triangulation** of the polytope  $P$  is a partition of  $P$  into **simplices** such that

- (1) The union of all these simplices equals the polytope.
- (2) Any pair of simplices intersects in a (possibly empty) common face.



## The size of a triangulation

Triangulations of  $d$ -polytopes come in different sizes!!



**Open Problem 3:** If a polytope has triangulations of sizes  $k_1$ ,  $k_2$ , with  $k_2 > k_1$ . Does it have a triangulation of every size  $S$ , with  $k_1 < S < k_2$ ?

Polytopes are not just pretty faces!!

Polytopes are **VERY** useful!

# Problem 4:

# THE SIMPLEX METHOD



## Linear Programming: Optimal decisions!!!

Optimal decisions by solving **Linear Programming Problems**:

$$\text{maximize } C_1x_1 + C_2x_2 + \dots + C_dx_d$$

among all  $x_1, x_2, \dots, x_d$ , satisfying:

$$a_{1,1}x_1 + a_{1,2}x_2 + \dots + a_{1,d}x_d \leq b_1$$

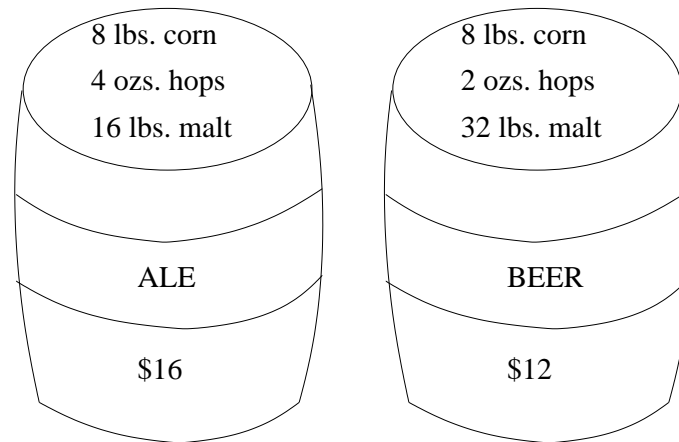
$$a_{2,1}x_1 + a_{2,2}x_2 + \dots + a_{2,d}x_d \leq b_2$$

⋮

$$a_{k,1}x_1 + a_{k,2}x_2 + \dots + a_{k,d}x_d \leq b_k$$

## The Brewery Problem

A brewery produces barrels of ale or beer, each needs fixed amount of corn, hops, and malt:



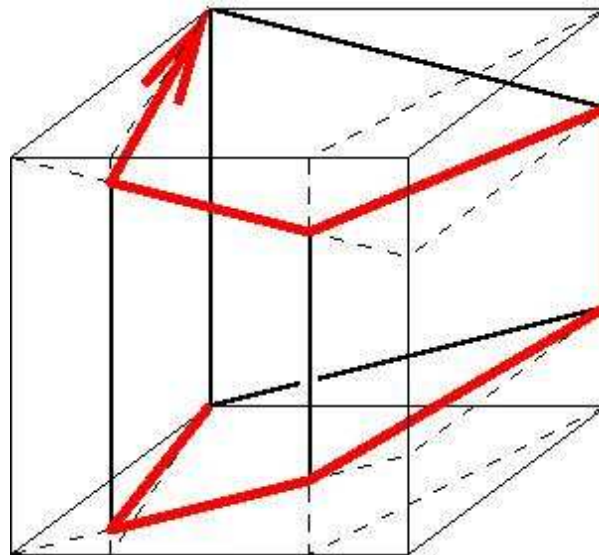
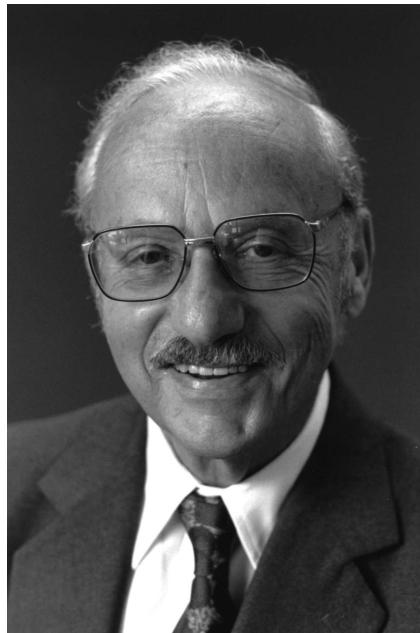
They only have 320 lbs. of corn, 130 ozs. of hops and 1120 lbs. of malt. Each barrel of ale yields a profit of \$16 and each barrel of beer yields a profit of \$12.

**How many barrels of ale beer should be produced to maximize total profit given limited resources?**

## The simplex method

**Lemma:** A vertex of the polytope is always an optimal solution for a linear program. We need to find a special vertex of the polytope!

The simplex method **walks** along the edges of the graph of the polytope, each time moving to a better and better cost vertex!



## Polynomial Hirsch Conjecture

Performance of the simplex method depends on the **diameter** of the graph of the polytope: largest distance between any pair of nodes.

The best UPPER bounds due to Barnette-Larman and Kalai-Kleitman are exponential:

$$(\#facets(P))^{\log(dim(P))+1} \text{ and } \frac{2^{dim(P)-2}}{3}(\#facets(P) - dim(P) + 5/2).$$

While the best LOWER bound on the diameter is linear! Due to Francisco Santos 2010.

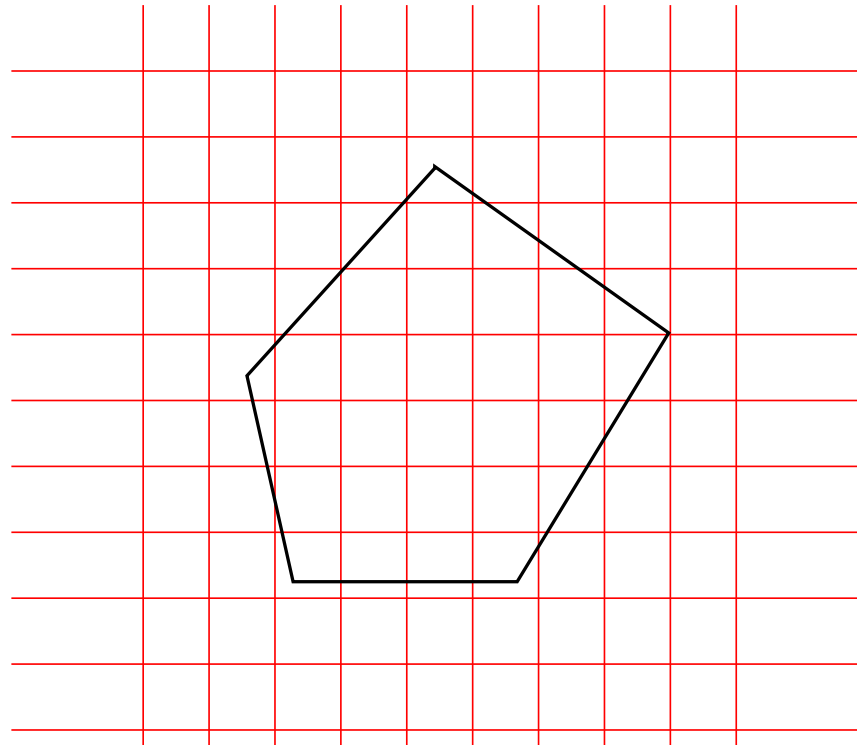
**OPEN QUESTION:** Is there a polynomial bound in terms of the number of facets and the dimension?

# Problem 5:

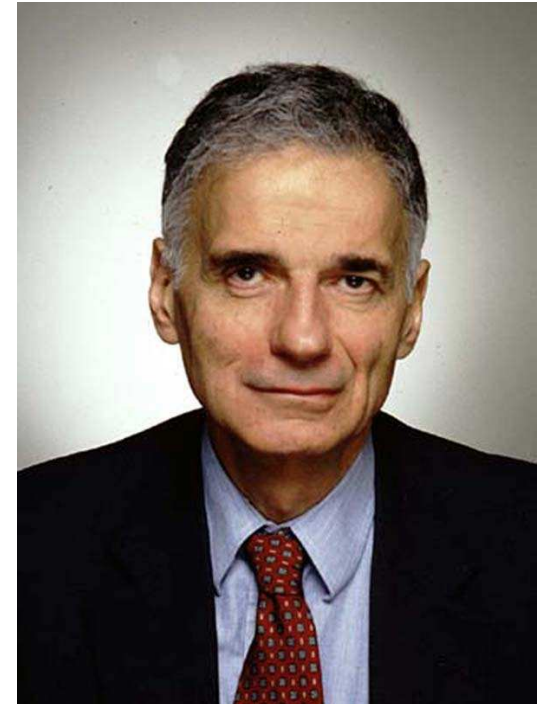
# COUNTING LATTICE POINTS

## Counting lattice points

Lattice points are those points with integer coordinates:  $\mathbb{Z}^n = \{(x_1, x_2, \dots, x_n) \mid x_i \text{ integer}\}$  We wish to count how many lie inside a given polytope!



# VOTING THEORY



- There are three (presidential) candidates  $a, b$  and  $c$ ....

- Let the preference orders of the  $n = \sum_{i=1}^6 n_i$  voters be

$$abc(n_1), acb(n_2), bac(n_3), bca(n_4), cab(n_5), cba(n_6)$$

- Here, there are  $n_1$  voters who rank candidate  $a$  as first,  $b$  second, and  $c$  third,  $n_2$  voters who rank  $b$  first,  $a$  second,  $c$  third, etc.
- Let us assume all 6 rankings or orderings are equally likely!
- Under **simple plurality voting**, the candidate with the most vote wins. (USA elections)
- In a **plurality runoff system**, if no candidate wins more than 50% of the vote, the two candidates with the highest vote count advance to a second voting round. (FRENCH elections)



- **Fact:** Different voting systems give different winners!!!
- **Question** What is the probability that the simple plurality and plurality runoff systems give different winners
- **POLYTOPES** are the answer:

Set up a system of equations that describes the situation.  $M$  denotes the total number of voters.

Say  $a$  wins by plurality but, using plurality runoff,  $b$  obtains higher score than  $c$  and a majority of voters then prefer  $b$  to  $a$ .

$$\begin{aligned}
 0 &< n_1 + n_2 - n_3 - n_4 \\
 0 &< n_3 + n_4 - n_5 - n_6 \\
 \frac{-M}{2} &< -n_1 - n_2 - n_5 \\
 M &= n_1 + n_2 + n_3 + n_4 + n_5 + n_6 \\
 0 &\leq n_i, i = 1, \dots, 6
 \end{aligned}$$

- The probability of this event equals

### Number of lattice points on polytope above

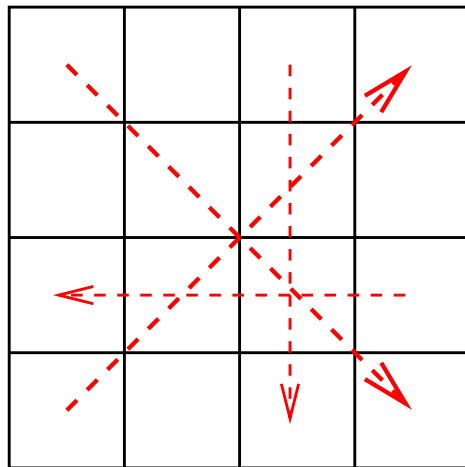
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number lattice pts in  $\{ (n_1, \dots, n_6) : n_1 + n_2 + \dots + n_6 = M, n_i \geq 0 \}$

- But all 6 possible voting rankings are possible. Multiply by 6 above number, because the plurality winner may be  $a$ ,  $b$  or  $c$  and the second position could be  $c$  not just  $b$ .
- **ANSWER:** For three candidates the probability the two voting systems give different winners for a large population is 12%.

## Combinatorics via Lattice points

Many objects can be counted as the lattice points in some polytope:  
 E.g., Sudoku configurations, matchings on graphs, and **MAGIC squares**:



12	0	5	7
0	12	7	5
7	5	0	12
5	7	12	0

5

**Open Problem** **HOW MANY**  $11 \times 11$  **magic squares with sum n are there?** Same as counting the points with integer coordinates inside the  $n$ -th dilation of a “magic square” polytope!

## Indeed, we can describe it by linear constraints!

The possible magic squares are non-negative integer solutions of a system of equations and inequalities: Ten equations, one for each row sum, column sum, and diagonal sum. For example,

$$x_{11} + x_{12} + x_{13} + x_{14} = 24, \text{ first row}$$

$$x_{13} + x_{23} + x_{33} + x_{43} = 24, \text{ third column, and of course } x_{ij} \geq 0$$

**We can do the same with tables of statistical data!!**

For example, we can ask questions about gender, race, and income as done U.S. census every ten years.

**Real Life Challenge:** Develop fast algorithms to count how many possible tables have particular **margin statistics** for big sizes!

# LIMITATION OF INFORMATION DISCLOSURE.

Gender = Male

Income Level

Race	$\leq$ \$10,000	$>$ \$10000 and $\leq$ \$25000	$>$ \$25000	Total
White	96	72	161	329
Black	10	7	6	23
Chinese	1	1	2	4
<b>Total</b>	<b>107</b>	<b>80</b>	<b>169</b>	<b>356</b>

Gender = Female

Income Level

Race	$\leq$ \$10,000	$>$ \$10000 and $\leq$ \$25000	$>$ \$25000	Total
White	186	127	51	364
Black	11	7	3	21
Chinese	0	1	0	1
<b>Total</b>	<b>197</b>	<b>135</b>	<b>54</b>	<b>386</b>

## Cyber-thieves care about this!

Counting lattice points can be used to factor integer numbers of shape  $n = pq$ ! Consider the 4-dimensional convex ball (not a polytope!):

$$B(n) = \{x \in R^4 : x_1^2 + x_2^2 + x_3^2 + x_4^2 \leq n\}$$

Let  $|B(n)|$  = the number of lattice points in  $B(n)$ , we have

$$|B(n)| - |B(n-1)| = 8 \sum_{m|n, m \neq 0 \pmod{4}} m$$

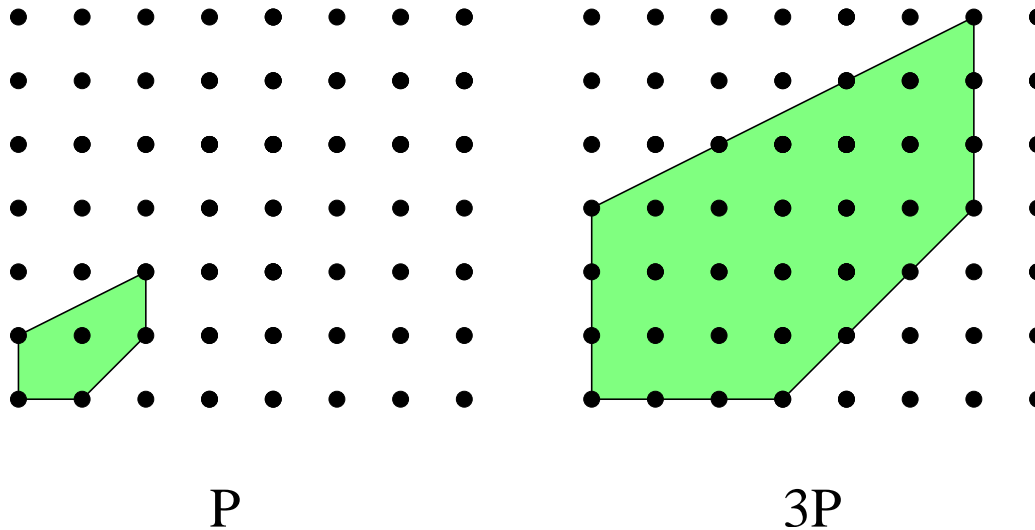
But since  $n = pq$ , with  $p, q$  primes, we have  $|B(n)| - |B(n-1)| = 8(1 + p + q + n)$

This, together with the fact that  $n = pq$  is sufficient to know  $p, q$ !  
**CREDIT CARDS AT RISK!**

## We can approximate the volume!

Let  $P$  be a convex polytope in  $\mathbb{R}^d$ . For each integer  $n \geq 1$ , let

$$nP = \{nq \mid q \in P\}$$



## Counting function approximates volume

For  $P$  a  $d$ -polytope, let

$$i(P, n) = \#(nP \cap \mathbb{Z}^d) = \#\{q \in P \mid nq \in \mathbb{Z}^d\}$$

This is the **number of lattice points in the dilation  $nP$** .

$$\text{Volume of } P = \lim_{n \rightarrow \infty} \frac{i(P, n)}{n^d}$$

At each dilation we can approximate the volume by placing a small unit cube centered at each lattice point:



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Thank you! Muchas Gracias!