

# Day 4 Talk 3

Santosh Vempala

"Progress on the KLS Conjecture"

with slides

Joint w Yin Tat Lee

An exercise:

Lemma: isotropic log concave  $p$ :

$$\mathbb{E}_{x,y \sim p} \langle x, y \rangle^3 \leq n^{1.5}$$

Proof.

$$\mathbb{E}_x \mathbb{E}_y \langle x, y \rangle^3 \leq \mathbb{E}_x (\mathbb{E}_y \langle x, y \rangle^2)^{3/2}$$
$$= \mathbb{E}_{x \sim p} \|x\|^3$$

$\leq \dots$  in slides

isoperimetry, definition of isoperimetric constant.

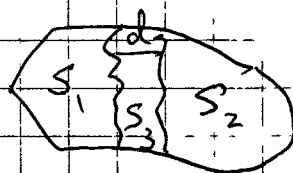
$$\gamma_p = \min_{S: p(S) \leq \frac{1}{2}} \frac{p(\partial S)}{p(S)}$$

Conjecture: halfspaces minimizers for isoperimetric constants for log-concave densities.

(LS, DF)

~~thm~~ thm:

$$p(S_3) \geq \frac{2 \cdot d}{D} \min(p(S_1), p(S_2))$$



Thm KLS95

~~PLA~~

$$P(S_3) \geq \frac{cd}{R} \min(p(S_1), p(S_2))$$

KLS Thm & Conjecture

Conj  $\psi_P \geq \frac{c}{\sqrt{V(A)}} \kappa_A$

Outline of talk:

- connections to other problems
- techniques
- best bounds
- open problems

Related to:

Algorithmic problem: Sampling

related

- min convex function
- compute volume
- complexity bounds for ball-walk depend on isoperimetric constant
- slicing conjecture
- Thin-shell conjecture

Bounds for

$\psi_P$  imply bounds in many other

problems:

- concentration
- Poincaré
- Thin-shell ...

KLS implies:  $L_p \approx G_p \approx \frac{1}{\sqrt{J_p}} \approx \frac{1}{\sqrt{r}}$

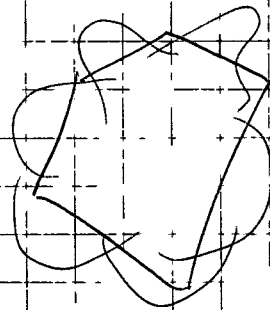
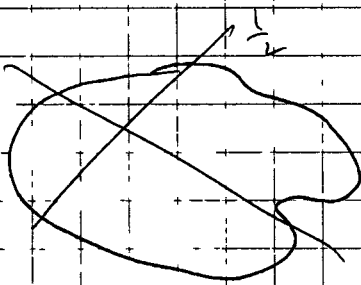
(see slide for details)

applications to metabolic networks

Techniques

- Localization & stochastic localization

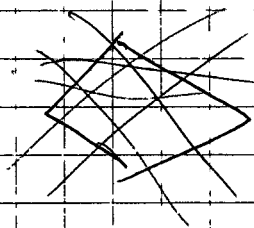
Pictures from talk



stochastic localization

see video from slides

→ gaussian



# The Kannan-Lovász-Simonovits Conjecture

Yin Tat Lee (U. Washington) and Santosh S. Vempala (Georgia Tech)



Thank You!

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## An exercise

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Lemma. For isotropic logconcave  $p$ :  $\mathbb{E}_{x,y \sim p} \langle x, y \rangle^3 \lesssim n^{1.5}$

Proof.

$$\begin{aligned} \mathbb{E}_{x \sim p} \mathbb{E}_{y \sim p} |\langle x, y \rangle|^3 &\lesssim \mathbb{E}_{x \sim p} \left( \mathbb{E}_{y \sim p} \langle x, y \rangle^2 \right)^{3/2} \\ &= \mathbb{E}_{x \sim p} \|x\|^3 \\ &\lesssim \left( \mathbb{E}_{x \sim p} \|x\|^2 \right)^{3/2} \\ &= n^{1.5}. \end{aligned}$$

Exercise. Prove a better bound.

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# Isoperimetry

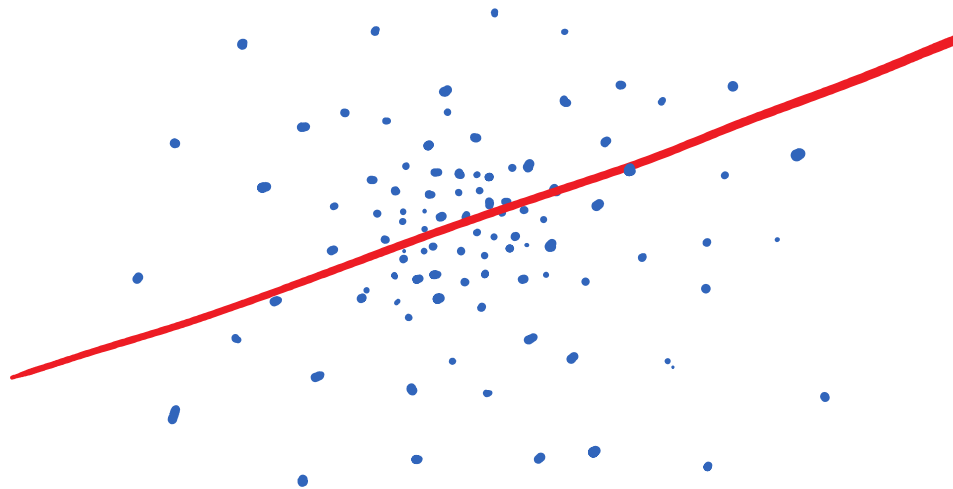
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Isoperimetric Ratio/Cheeger Constant/Expansion of a function  $p$ :

$$\psi_p = \min_{S:p(S)\leq\frac{1}{2}} \frac{p(\partial S)}{p(S)}$$

Q. What is the Cheeger constant of the Gaussian distribution?

A. The isoperimetric ratio of a halfspace through its centroid:  $\sqrt{\frac{2}{\pi}}$



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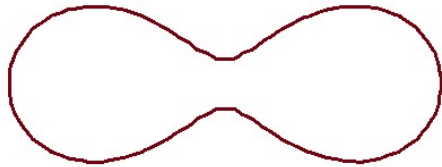
In fact for any  $0 < t < 1$ , the subset of measure  $t$  with minimum surface area is a halfspace!



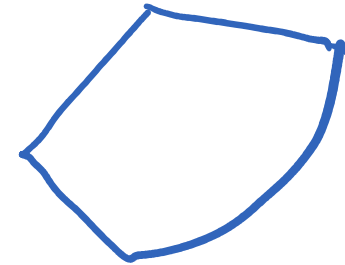
# Isoperimetry

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Isoperimetric Ratio/Cheeger Constant/Expansion:



$$\psi_p = \min_{S: p(S) \leq \frac{1}{2}} \frac{p(\partial S)}{p(S)}$$



Can be arbitrarily small...  
Structured distributions?

Logconcave function:  $f(\lambda x + (1 - \lambda)y) \geq f(x)^\lambda f(y)^{1-\lambda}$   
(nonnegative function whose logarithm is concave)

Common generalization of Gaussians and indicators of convex sets.



Halfspace cuts do not have to be the minimal ones, but...

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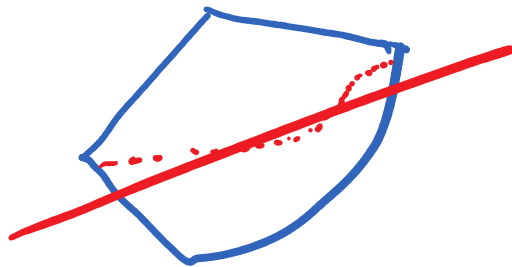
# The Conjecture

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Isoperimetric Ratio/Cheeger Constant/Expansion:

$$\psi_p = \min_{S:p(S)\leq\frac{1}{2}} \frac{p(\partial S)}{p(S)}$$

Logconcave function:  $f(\lambda x + (1 - \lambda)y) \geq f(x)^\lambda f(y)^{1-\lambda}$



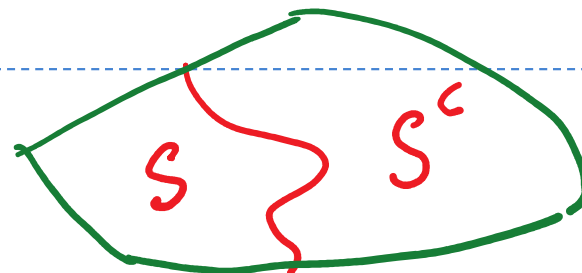
Conjecture: For any logconcave density in any dimension, halfspaces minimize the isoperimetric ratio up to an absolute universal constant.

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# KLS Theorem

Thm [LS, DF]  $p(\partial S) \geq \frac{2}{D} \min(p(S), p(S^c))$



(special case of isoperimetry for Riemannian manifolds with nonnegative curvature)

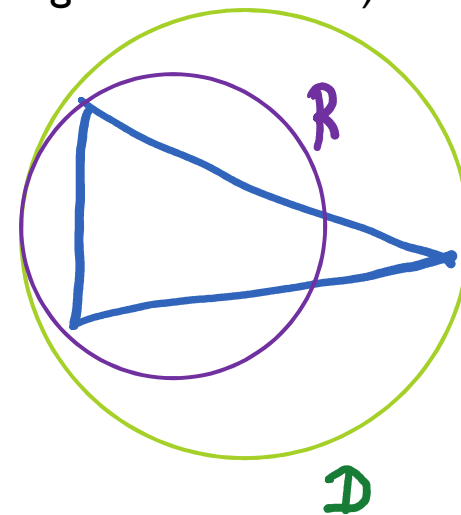
$A = \mathbb{E}_p((x - \bar{x})(x - \bar{x})^T)$  : covariance matrix of  $p$

$R^2 = \mathbb{E}_p(\|x - \bar{x}\|^2) = \text{Tr}(A) = \sum_i \lambda_i(A)$

Thm. [KLS95]. For any logconcave density,

$$p(\partial S) \geq \frac{c}{R} \min(p(S), p(S^c)).$$

(note: isotropic distribution has  $A = I$ . So,  $\psi_p \geq \frac{c}{\sqrt{n}}$  for isotropic  $p$ )



Slide 7

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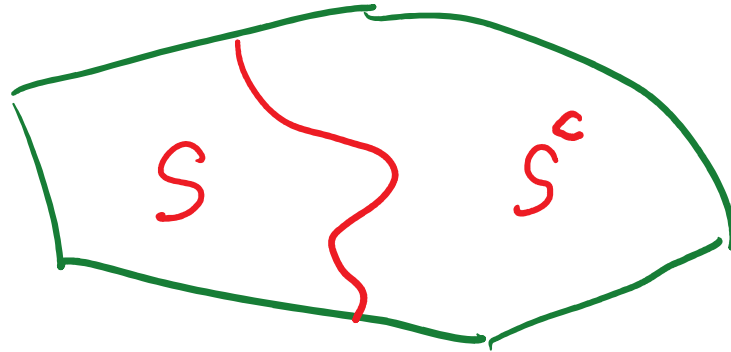
YTL1

Draw  $D$  and  $R^2$ ?

Yin Tat Lee, 11/14/2017

# KLS Theorem and Conjecture

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$A$  = covariance matrix of  $p$

$$R^2 = \mathbb{E}_p(\|x - \bar{x}\|^2) = \text{Tr}(A) = \sum_i \lambda_i(A)$$

Thm. [KLS95].  $\psi_p \geq \frac{c}{\sqrt{\text{Tr}(A)}} = \frac{c}{\sqrt{n}}$  for isotropic  $p$ .

**Conj.** [KLS95].  $\psi_p \geq \frac{c}{\sqrt{\lambda_1(A)}} = \Omega(1)$  for isotropic  $p$ .



# Outline

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Original Motivation

Connections

- ▶ Probability and Geometry
- ▶ Algorithms

Techniques

- ▶ Localization
- ▶ Stochastic Localization

Some open problems



# An algorithmic problem: Sampling

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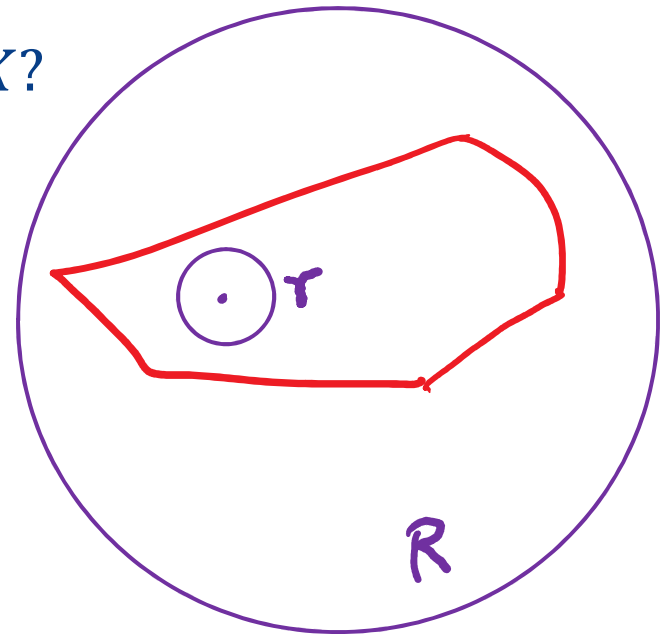
Given convex body  $K$ , generate uniform random point in  $K$ .

$K$  specified by a “well-guaranteed” membership oracle:

- ▶  $x_0, r, R: x_0 + rB_n \subseteq K \subseteq RB_n$
- ▶ An oracle that answers YES/NO to  $x \in K$ ?

Related problems we will see later:

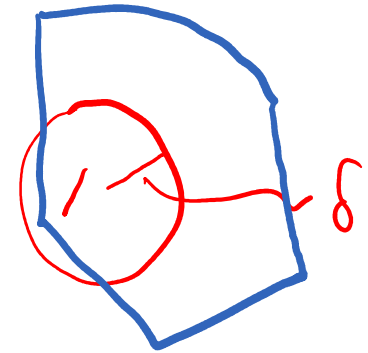
- ▶ Compute the volume of  $K$
- ▶ Minimize a convex function over  $K$



# Sampling with the ball walk

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At  $x$ , pick random  $y$  from  $x + \delta B_n$ ,  
if  $y$  is in  $K$ , go to  $y$ .



Approaches the uniform distribution over  $K$ .

Rate of convergence?

Cheeger constant of Markov chain...

$$\psi_K = \min_S \frac{\text{vol}(\partial S)}{\min(\text{vol}(S), \text{vol}(S^c))}$$

Thm. [KLS97]

Mixing time of the ball walk from a warm start is  $O^* \left( \frac{n^2}{\psi_K^2} \right)$ .

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# Connections I: Geometry and Probability

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The KLS conjecture has very interesting consequences, many of which were conjectured independently and earlier.

Slicing/small ball probability

Thin shell/Central Limit theorem

Poincaré/Lipschitz Concentration

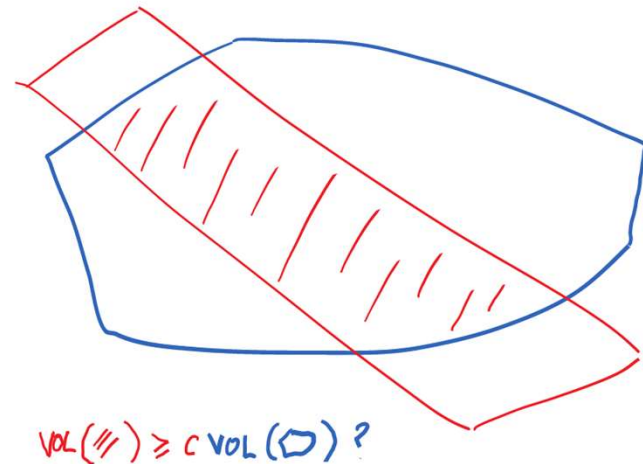




# The Slicing conjecture: anti-concentration

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Convex body of volume one has a hyperplane section of volume at least some constant.



Equivalently: for any isotropic logconcave density  $f$ ,

$$L_p = f(0)^{1/n} = O(1).$$

(Paouris) Slicing implies  $\Pr(\|x\|_2 \leq \epsilon\sqrt{n}) \leq (C\epsilon)^n$

(Ball) KLS  $\Rightarrow$  Slicing

(Klartag; Bourgain)  $L_p \lesssim n^{1/4}$

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# The Thin-shell conjecture: a CLT

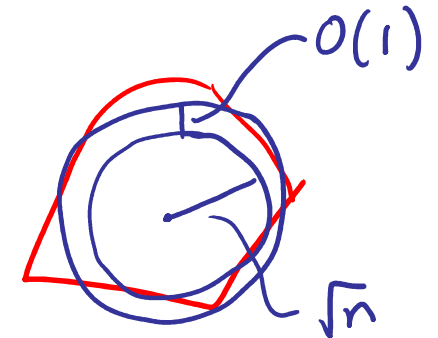
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For  $X$  from any isotropic logconcave distribution  $p$ ,

$$\text{Var}_p \left( \|X\|^2 \right) = O(n)$$

or

$$\sigma_p^2 = \mathbb{E}_p \left( (\|X\| - \sqrt{n})^2 \right) = O(1)$$



“Most of an isotropic logconcave distribution is contained in an annulus of constant thickness.”

KLS  $\Rightarrow$  thin-shell. In fact,  $\sigma_p \lesssim \frac{1}{\psi_p}$

Thm [Eldan-Klartag].  $L_p \lesssim \sigma_p$

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**Slide 14**

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**YTL2**

You never defined  $\sigma_p$  and  $L_p$

Yin Tat Lee, 11/14/2017

# The Thin-shell conjecture: a CLT

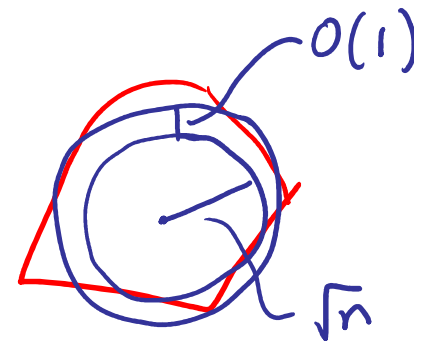
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$X$  from any isotropic logconcave distribution,

$$\mathbb{E} \left( \left( \|X\| - \sqrt{n} \right)^2 \right) = O(1)$$

$$\sigma_p \lesssim \frac{1}{\psi_p}$$

[Eldan-Klartag].  $L_p \lesssim \sup_p \sigma_p$



CLT: Most marginals are approximately Gaussian.

For an isotropic convex body  $K$ , let  $g_\theta(s) = \text{vol}(K \cap \{x: x^T \theta = s\})$

Then,

$$\Pr \left( \left\{ \theta \in S^{n-1}: \max_t \left| \int_{-\infty}^t g_\theta(s) - \int_{-\infty}^t \gamma(s) \right| \lesssim \delta + \frac{\psi_K}{\sqrt{n}} \right\} \right) \geq 1 - n \exp(-\Omega(\delta^2 n))$$



# Progress on the thin-shell bound

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Year/Authors	bound on $\sigma_p$
2006/Klartag	$\sqrt{\frac{n}{\log n}}$
2006/Fleury-Guédon-Paouris	$\sqrt{n} \frac{(\log \log n)^2}{\log^{1/6} n}$
2006/Klartag	$n^{2/5}$
2010/Fleury	$n^{3/8}$
2011/Guédon-E. Milman	$n^{1/3}$
2016/Lee-Vempala	$n^{1/4}$



# Poincaré Conjecture

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Smooth function  $g$

Isotropic logconcave density  $p$

There exists a universal constant  $c$  s.t.

$$\forall g: \zeta_p = \inf \frac{\mathbb{E}_p(\|\nabla g\|^2)}{\text{Var}_p(g)} \geq c$$

Thm. [Mazja, Cheeger; Buser; Ledoux]  $\zeta_p \approx \psi_p^2$



# Lipschitz concentration

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Lipschitz function  $g$  on the sphere

Levy's classical concentration:

$$\Pr(|g(x) - \mathbb{E}_{S^{n-1}}(g)| \geq t) \leq 2e^{-ct^2n}$$

Thm.[Gromov-Milman]  $L$ -Lipschitz function  $g$  in  $R^n$ , isotropic logconcave density  $p$ ,

$$\Pr_p(|g(x) - \mathbb{E}(g)| > L \cdot t) \leq e^{-\Omega(t\psi_p)}$$

Thm. [E. Milman] Concentration  $\Rightarrow$  KLS.

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## Entropy gaps and jumps

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$\text{Ent}(X) = -E(\log p(x))$ . How random is a distribution?

[Shannon-Stam] Entropy gap: equality only for Gaussian

$$X, Y \sim p: \quad \text{Ent}\left(\frac{X + Y}{\sqrt{2}}\right) \geq \text{Ent}(X)$$

[Ball, Nguyen] Entropy gap jump.  $X, Y \sim p, Z \sim N(0, I)$

$$\text{Ent}\left(\frac{X + Y}{\sqrt{2}}\right) - \text{Ent}(X) \gtrsim \psi_p^2(\text{Ent}(Z) - \text{Ent}(X))$$





# Connections: Geometry and Probability

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**Slicing conjecture:**  $L_p = p(0)^{1/n} = O(1)$

**Thin-Shell conjecture:**  $\sigma_p = \mathbb{E}(\|x\| - \sqrt{n})^2 = O(1)$

**Poincaré conjecture:**  $\zeta_p = \inf_g \frac{\mathbb{E}_p(\|\nabla g\|^2)}{\text{Var}_p(g)} = \Omega(1)$

**Generalized Levy concentration:**

Lipschitz  $f$  with  $\mathbb{E}f = 0$ ,  $\mathbb{P}(f(x) > t) = \exp(-\Omega(t))$ .

**KLS conjecture implies all:**  $L_p \lesssim \sigma_p \lesssim \frac{1}{\sqrt{\zeta_p}} \approx \frac{1}{\psi_p}$

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# Connections: Algorithms

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Sampling

Optimization

Volume Computation/Integration

Learning

What is the **complexity** of computational problems **as the dimension grows?**

Dimension = number of variables

Typically, size of input is a function of the dimension.



# Computational model

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## Well-guaranteed Membership oracle:

Compact set  $K$  is given by

a membership oracle: answers YES/NO to “ $x \in K?$ ”

a point  $x_0 \in K$

Numbers  $r, R$  s.t.  $x_0 + rB^n \subseteq K \subseteq RB^n$

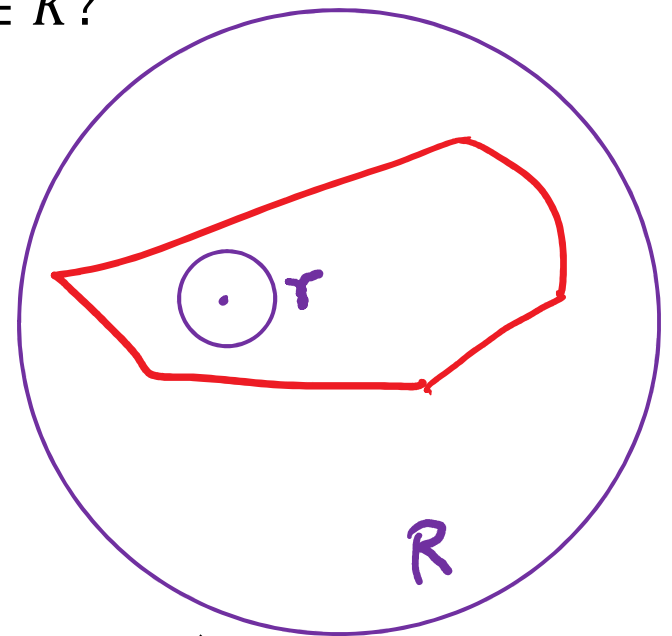
## Well-guaranteed Function oracle

An oracle that returns  $f(x)$  for any  $x \in R^n$

A point  $x_0$  with  $f(x_0) \geq \beta$

Numbers  $r, R$  s.t.

$$x_0 + rB^n \subset L_f\left(\frac{1}{8}\right) \quad \text{and} \quad R^2 = \mathbb{E}_f\left(\|X - \bar{X}\|^2\right)$$



# Problem 1: Sampling

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Input: function  $f: R^n \rightarrow R_+$ ,  $\int f < \infty$ , specified by an oracle, a point  $x$ , error parameter  $\varepsilon$ .

Output: A point  $y$  from a distribution within distance  $\varepsilon$  of distribution with density proportional to  $f$ .

Examples:  $f(x) = 1_K(x)$ ,  $f(x) = e^{-a\|x\|} 1_K(x)$



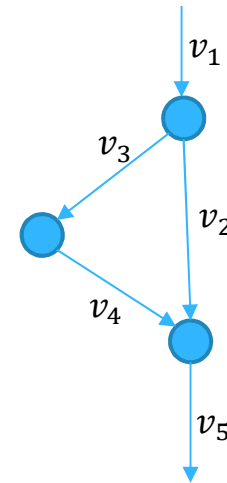
# Sampling metabolic networks

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Given a metabolic network  $S \in \mathbb{R}^{m \times n}$  on  $m$  metabolites and  $n$  reactions.

Find mass conserving flow  $v \in \mathbb{R}^n$  with bounds  $l, u \in \mathbb{R}^n$ :

$$Sv = b$$
$$l \leq v \leq u$$



Many possible  $v$ 's, which one to pick?

- ▶ Sample over all possible values of  $v$ !

Dimension: 5000-100,000.

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# Analysis of metabolic networks

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Sampling enables an unbiased study of all feasible metabolic flows

- ▶ Could optimize with respect to an objective function, but it is unclear if the human body acts like this

Can also compute the volume of the space!

- ▶ Price et al. (2004) analyze the human red blood cell metabolic network (dim=11)
- ▶ They observe the volume of diseased patient's networks is significantly lower.



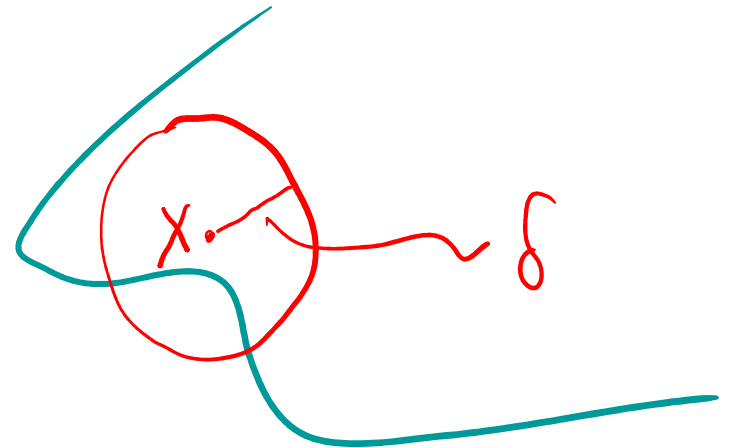
# How to Sample?

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Ball walk:

At  $x$ ,

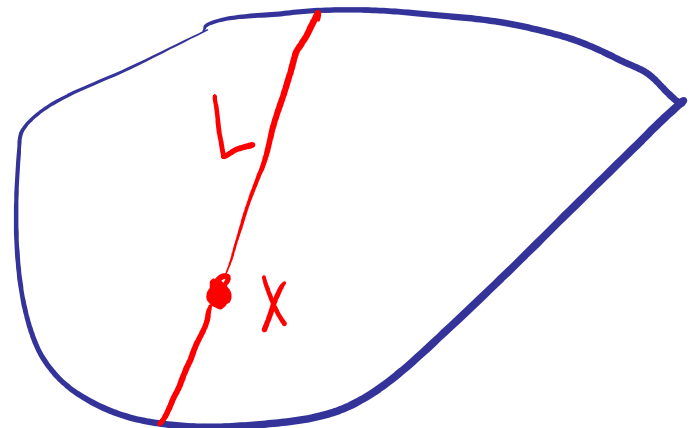
- pick random  $y$  from  $x + \delta B_n$
- if  $y$  is in  $K$ , go to  $y$



Hit-and-Run:

At  $x$ ,

- pick a random chord  $L$  through  $x$
- go to a random point  $y$  on  $L$



# Markov chains

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State space  $K$ , next step distribution  $P_u(\cdot)$  associated with each point  $u$  in  $K$ .

Stationary distribution  $Q$ , ergodic “flow” defined as

$$\Phi(A) = \int_A P_u(K \setminus A) dQ(u)$$

For a stationary distribution, we have  $\Phi(A) = \Phi(K \setminus A)$

Conductance:

$$\phi(A) = \frac{\int_A P_u(K \setminus A) dQ(u)}{\min Q(A), Q(K \setminus A)} \quad \phi = \inf \phi(A)$$

Thm. [LS93]  $Q_t$ : distribution after  $t$  steps

$$M = \sup_{A \subset K} \frac{Q_0(A)}{Q(A)} : d_{TV}(Q_t, Q) \leq \sqrt{M} \left(1 - \frac{\phi^2}{2}\right)^t$$

$$M = E_{Q_0} \left( \frac{Q_0(x)}{Q(x)} \right) : d_{TV}(Q_t, Q) \leq \epsilon + \sqrt{\frac{M}{\epsilon}} \left(1 - \frac{\phi^2}{2}\right)^t \quad \forall \epsilon > 0$$

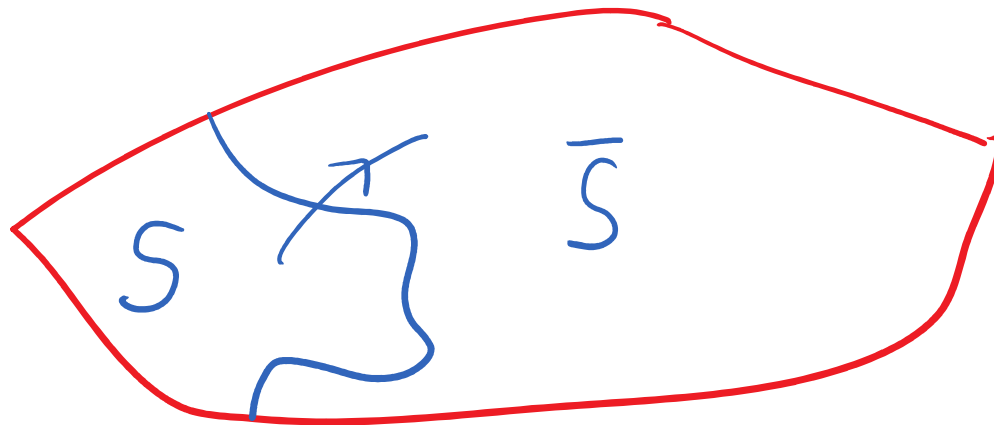




# Conductance

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Consider an arbitrary measurable subset  $S$ .



Need to show that the escape probability from  $S$  is large.

(Smoothness of 1-step distribution) Points that do not cross over are far from each other i.e., nearby points have large overlap in 1-step distributions

(Isoperimetry) Large subsets have large boundaries

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# Convergence of ball walk

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Theorem [KLS97]. The ball walk applied to an isotropic logconcave density  $p$ , from a warm start, converges in

$O^* \left( \frac{n^2}{\psi_p^2} \right)$  steps.

“Cheeger constant of this Markov chain is determined by Cheeger constant of its stationary distribution”



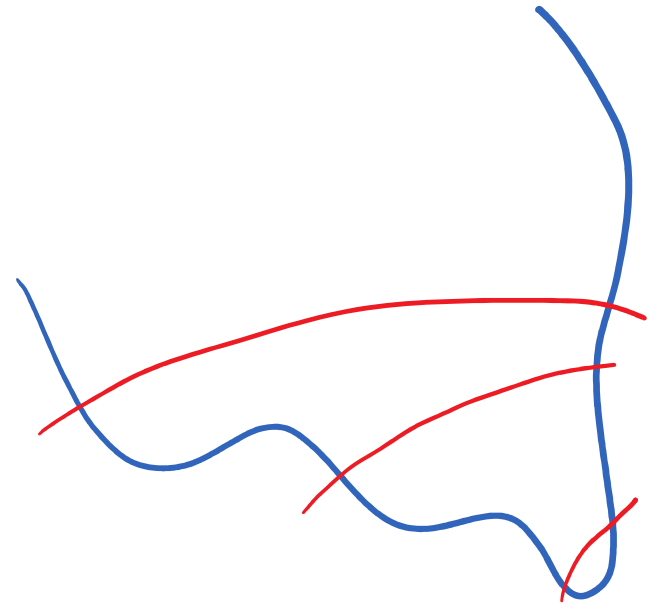
## Problem 2: Optimization

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Input: function  $f: R^n \rightarrow R$  specified by an oracle,  
point  $x$ , error parameter  $\epsilon$  .

Output: point  $y$  such that

$$f(y) \geq \max f - \epsilon$$



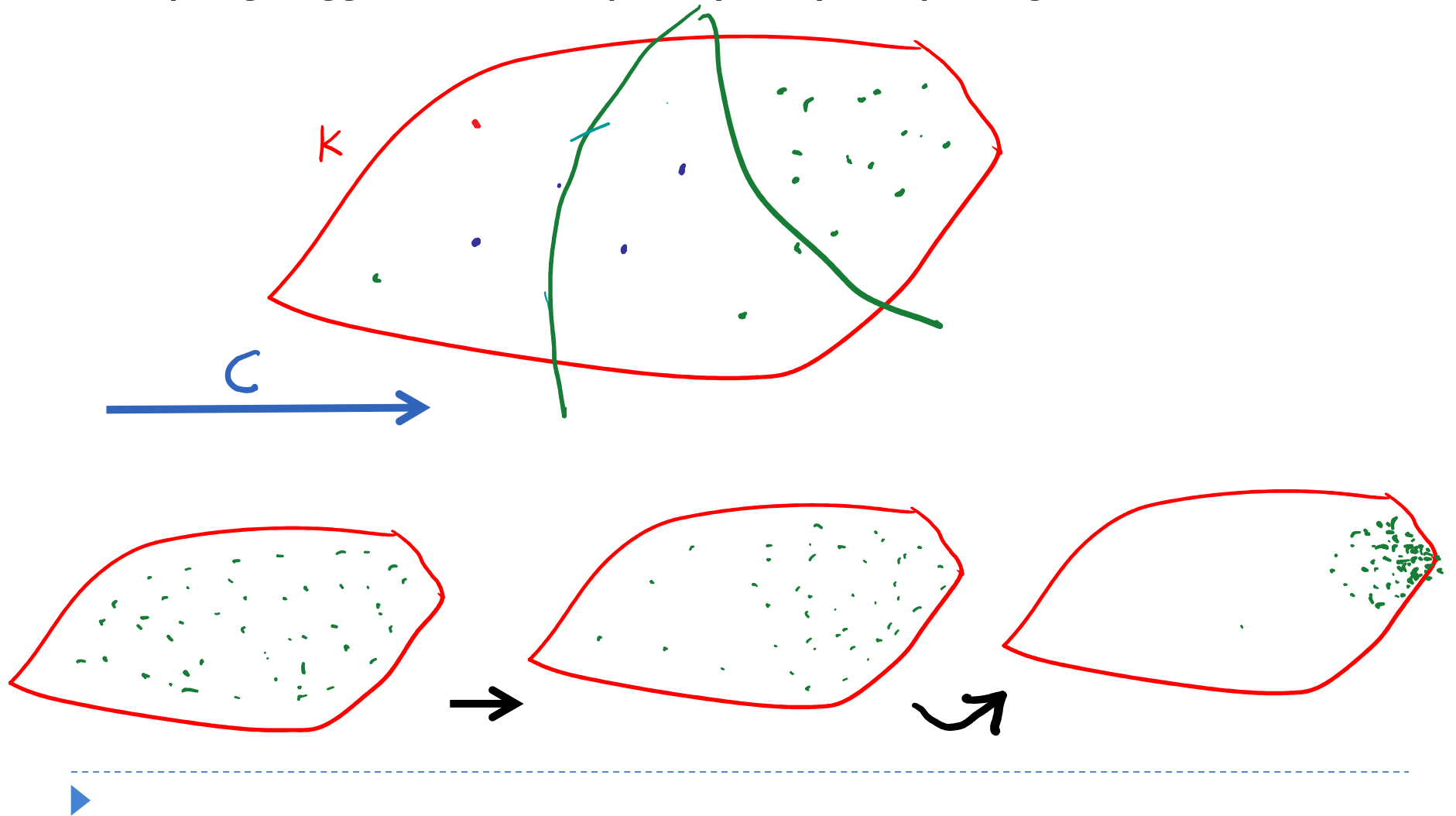
Examples:  $\max c \cdot x$  s.t.  $Ax \geq b$  ,  $\min ||x||$  s.t.  $x \in K$ .



# Optimization from membership

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Sampling suggests a conceptually very simple algorithm.



# Simulated Annealing [Kalai-V.04]

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To optimize  $f$  consider a sequence  $f_0, f_1, f_2, \dots$ ,  
with  $f_i$  more and more concentrated near the optimum.

$$f_i(x) = e^{-t_i \langle c, x \rangle}$$

Corresponding distributions:

$$P_{t_i}(x) = \frac{e^{-t_i \langle c, x \rangle}}{\int_K e^{-t_i \langle c, x \rangle} dx}$$

**Lemma.**  $E_{P_t}(c \cdot x) \leq \min c \cdot x + \frac{n}{t}$ .

So going up to  $t = \frac{n}{\epsilon}$  suffices to obtain an  $\epsilon$  approximation.

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# Volume Computation

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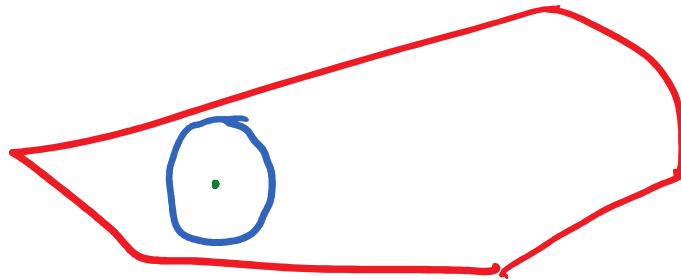
Given a measurable, compact set  $K$  in  $n$ -dimensional space and  $\epsilon > 0$ , find a number  $A$  such that:

$$(1 - \epsilon) \text{ volume}(K) \leq A \leq (1 + \epsilon) \text{ volume}(K)$$

$K$  is given by

a point  $x_0 \in K$ , s.t.  $x_0 + B_n \subseteq K \subseteq RB_n$

a membership oracle: answers YES/NO to " $x \in K$ ?"



## *Randomized Volume/Integration*

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[DFK89]. Polytime **randomized** algorithm that estimates volume to within relative error  $(1 + \epsilon)$  with probability at least  $1 - \delta$  in time  $\text{poly}(n, \frac{1}{\epsilon}, \log(\frac{1}{\delta}))$ .

[Applegate-K91]. Polytime randomized algorithm to estimate integral of any (Lipshitz) logconcave function.



# Progress on Volume Computation

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	Power	New aspects
Dyer-Frieze-Kannan 89	23	everything
Lovász-Simonovits 90	16	localization
Applegate-K 90	10	logconcave integration
L 90	10	ball walk
DF 91	8	error analysis
LS 93	7	multiple improvements
KLS 97	5	speedy walk, isotropy
LV 03,04	4	annealing, isoperimetry
LV 06	4	integration, local analysis
Cousins-V. 15 (well-rounded)	3	Gaussian cooling

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# Does it work?

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[Cousins-V.13] Matlab implementation of a new algorithm

- ▶ “volume computation matlab”
- ▶ <https://volumecomputation.wordpress.com/>
- ▶ Incorporated into the COBRA toolbox for Systems Biology



# Outline

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Original Motivation

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- ▶ Algorithms

**Techniques**

- ▶ Localization
- ▶ Stochastic Localization

Some open problems



# Localization

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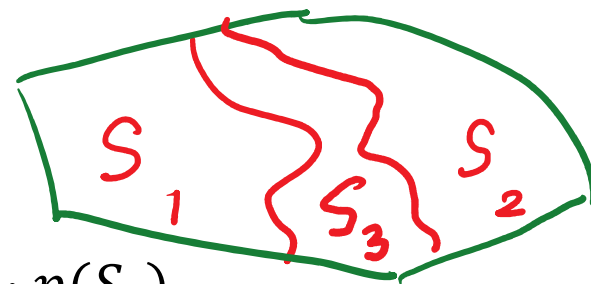
Idea: Reduce inequalities in high dimension to inequalities in one dimension.



# Isoperimetry via localization

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$$p(S_3) \geq \frac{2d(S_1, S_2)}{D} \min p(S_1), p(S_2)$$



Write as 2 inequalities:  $p(S_1) \leq p(S_2)$ ,  $p(S_3) \geq \psi \cdot p(S_1)$

Let  $g(x) = f(x)(1_{S_2}(x) - 1_{S_1}(x))$ ,  $h(x) = f(x)(\psi \cdot 1_{S_1}(x) - 1_{S_3}(x))$

Then, need to show:  $\int g \geq 0 \Rightarrow \int h \leq 0$ .

Suppose not, i.e.,  $\exists S_1, S_2, S_3: \int g \geq 0, \int h > 0$ .

Idea:

1. No such counterexample in one dimension
2. If such a counterexample exists in some dimension, then it also exists in 1 dimension.



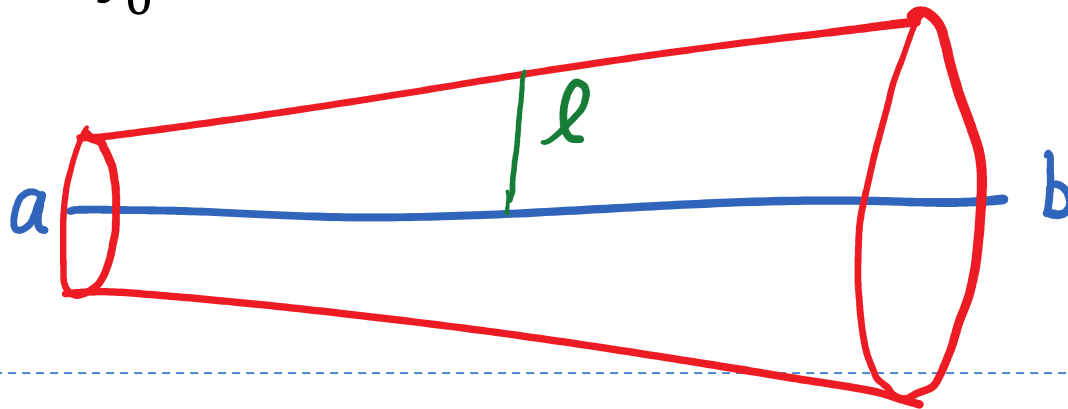
## Localization Lemma [LS, KLS]

---

Lemma. Let  $g, h: R^n \rightarrow R$  be integrable, lower semi-continuous functions. Suppose  $\int g, \int h > 0$ . Then, there exists an interval  $[a, b] \subset R^n$  and a linear function  $\ell: [0, 1] \rightarrow R_+$  s.t.

$$\int_0^1 g((1-t)a + tb)\ell(t)^{n-1} dt > 0$$

$$\int_0^1 h((1-t)a + tb)\ell(t)^{n-1} dt > 0.$$



# Localization lemma

---

$$g, h: R^n \rightarrow R, \int g, \int h > 0.$$

1. Find a bisecting halfspace for one function
2. Show support of limit of bisections is an interval or a point.
3. The limit function has a concave profile
4. Reduce to linear cross-sectional profile.

[Fradelizi-Guédon] Extremal characterization and generalization to multiple inequalities.

---



# Isoperimetry via localization

---

$$\text{vol}(S_1) \leq \text{vol}(S_2) \Rightarrow \text{vol}(S_3) \geq \psi \text{vol}(S_1)$$

$$g(x) = 1_{S_2}(x) - 1_{S_1}(x), \quad h(x) = \psi 1_{S_1}(x) - 1_{S_3}(x)$$

Need to show:  $\int g \geq 0 \Rightarrow \int h \leq 0$ .

Suppose not, i.e., for some partition,  $\int g \geq 0, \int h > 0$ .

Applying localization,

$$\int_0^1 g((1-t)a + tb) \ell(t)^{n-1} dt > 0, \quad \int_0^1 h((1-t)a + tb) \ell(t)^{n-1} dt > 0.$$

Let  $Z_i = \{t \in [0,1]: (1-t)a + tb \in S_i\}$ ,  $F(t) = f((1-t)a + tb) \ell(t)^{n-1}$ .

Then this means that

$$\int_{Z_1} F \leq \int_{Z_2} F, \quad \text{but} \quad \int_{Z_3} F < \psi \int_{Z_1} F, \quad \text{a 1-d counterexample must exist.}$$



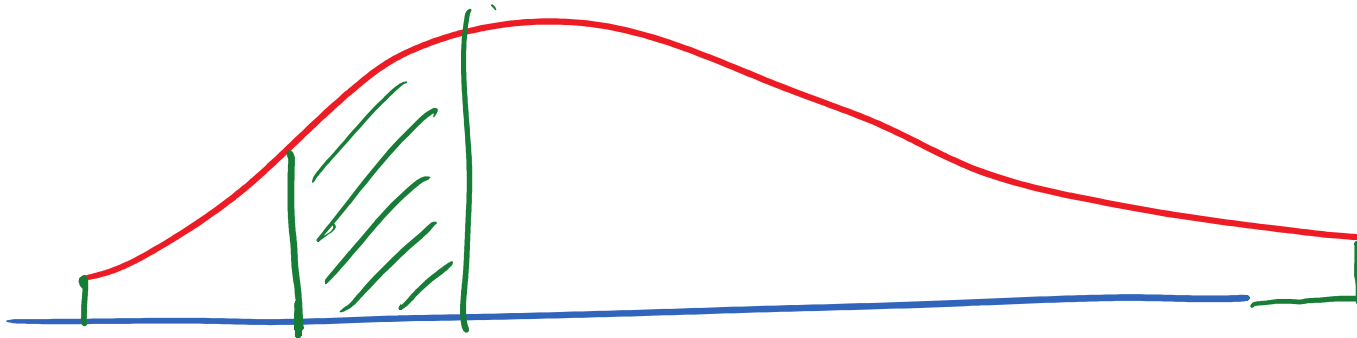
# One-dimensional isoperimetry

---

For any logconcave function:

$$\int_{S_3} F \geq \frac{2d(S_1, S_2)}{D} \min \int_{S_1} F, \int_{S_2} F$$

Suffices to show it for partition of into 3 intervals.



Without factor of 2, follows from unimodality!  
Therefore, same isoperimetric ratio holds in  $R^n$ .

---





# Localization: many applications

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Many other isoperimetric inequalities

E.g. the KLS conjecture holds for a Gaussian restricted by any logconcave function.

Thm. For a density  $h$  proportional to  $e^{-\frac{t}{2}\|x\|^2} f(x)$  for any logconcave function  $f$ , we have  $\psi_p \gtrsim \sqrt{t}$ .

Analysis of [Cousins-V.2015] algorithm:

Thm. Volume of a well-rounded convex body ( $B \subseteq K \subseteq \tilde{O}(\sqrt{n})$ ) can be computed in  $O^*(n^3)$  steps!

Carbery-Wright anti-concentration of polynomials

---



# Stochastic localization

---

Goal: Lower bound on expansion of subset  $S$  of measure  $1/2$ .

Idea: Apply hyperplane bisections randomly. Show measures of sets remain close to original. Prove isoperimetry for (hopefully) simpler distribution.

*(I tried this for years, still think it might work 😊)*

*(I also think there will be world peace, and I will stop eating junk food tomorrow)*

Meanwhile, **Eldan**: Apply infinitesimal linear reweighting in random direction to maintain expected value of density at each point.

[Eldan2012] used this to prove that the thin-shell conjecture implies the KLS constant up to a  $\log n$  factor.

---



# Eldan's Stochastic Localization

---

The process starts with  $p_0(x) = p(x)$  and maintains a density  $p_t$  at time  $t$  with mean  $\mu_t = E_{p_t}(x)$ .

The infinitesimal change is

$$dp_t(x) = (x - \mu_t)^T dW_t \cdot p_t(x)$$

where  $dW_t$  is an infinitesimal Gaussian (Wiener process).

We can imagine this discretely as

$$p_{t+h}(x) = p_t(x) \left( 1 + \sqrt{h}(x - \mu_t)^T W_t \right)$$

where  $W_t \sim N(0, I)$  is a standard Gaussian.

---



# Stochastic localization apps

---

Thm [Eldan12]. Let  $\sigma(n) = \sup_p \sigma_p$ . Then,

$$\psi_p \gtrsim \frac{\log n}{\sigma(n)}$$

Thm [LeeVI6]. For any logconcave density  $p$  with covariance  $A$ :

$$\psi_p \gtrsim \frac{1}{\text{Tr}(A^2)^{1/4}}$$

For isotropic logconcave  $p$ ,  $\psi_p \gtrsim n^{-1/4}$ .

Thm [LeeVI7]. Log-Sobolev constant of isotropic logconcave  $p$  with support of diameter  $D$  is  $[\text{nexttalk}]$ . This bound is tight.

---



# The KLS constant

---

KLS theorem.

$$\psi_p \gtrsim \frac{1}{\text{Tr}(A)^{1/2}} = \frac{1}{(\sum_i \lambda_i(A))^{1/2}}$$

LeeV theorem.

$$\psi_p \gtrsim \frac{1}{\text{Tr}(A^2)^{1/4}} = \frac{1}{(\sum_i \lambda_i(A)^2)^{1/4}}$$

KLS conjecture.

$$\psi_p \gtrsim \frac{1}{\|A\|_{op}^{1/2}} = \frac{1}{\lambda_1(A)^{1/2}}$$

---



# Proof Strategy

---

$p_t$  is a martingale:

$$\mathbb{E}p_t = p_0$$

Suffices to prove the theorem for  $p_t$ :

$$\psi_p = \Omega(\psi_{p_T}) \text{ for "large" } T.$$

$p_t$  has large Cheeger constant:

$$\psi_{p_T} \gtrsim \sqrt{T}$$



# Why does $p_T$ have good expansion?

---

Localization

Stochastic localization

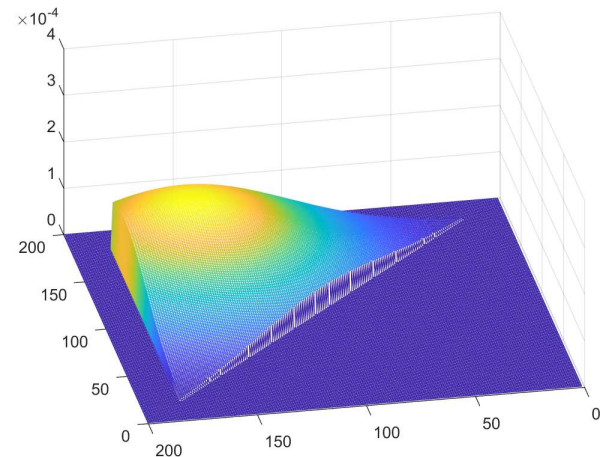
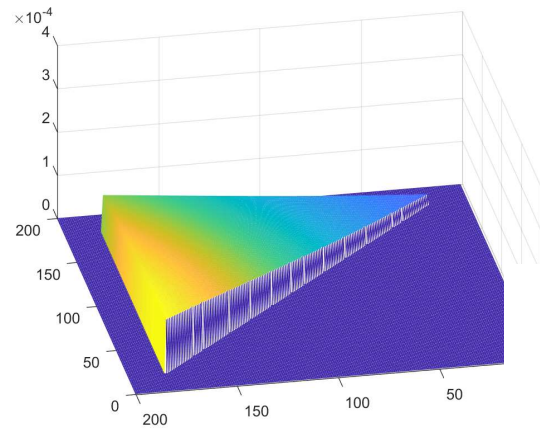
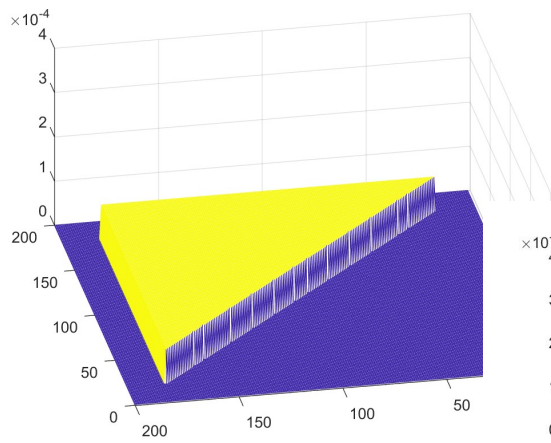
But's let's see it for real...

---



# What is happening?

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Really?

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# Emergence of Gaussian factor

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$$dp_t(x) = (x - \mu_t)^T dW_t \cdot p_t(x)$$

We will see that:

$$d \log p_t(x) = (x - \mu_t)^T dW_t - \frac{1}{2} \|x - \mu_t\|^2 dt$$

Itô's lemma:  $dX_t = \mu_t dt + \sigma_t dW_t \Rightarrow df(X_t) = \frac{df(X_t)}{dX} dX_t + \frac{1}{2} \frac{d^2f(X_t)}{dx^2} \sigma_t^2 dt$

This is by Taylor expansion and noting that  $(dW_t)^2 = dt$ .

Applying this,

$$\begin{aligned} d \log p_t(x) &= \frac{dp_t(x)}{p_t(x)} - \frac{1}{2} \frac{(\|x - \mu_t\| \cdot p_t(x))^2}{p_t(x)^2} dt \\ &= (x - \mu_t)^T dW_t - \frac{1}{2} \|x - \mu_t\|^2 dt \\ &= x^T (\mu_t dt + dW_t) - \frac{1}{2} \|x\|^2 dt + g(t) \end{aligned}$$

where the last term does not depend on  $x$ .

Therefore,

$$p_t(x) \propto e^{x^T c_t - \frac{t}{2} \|x\|^2} p(x) \quad \text{for } dc_t = \mu_t dt + dW_t$$



# KLS is easy for Gaussianic distribution

---

Thm [Bakry-Ledoux96, also Bobkov2000, Cousins-V. 2013 by localization]

For a density  $h$  proportional to  $e^{-\frac{t}{2}\|x\|^2} f(x)$  for any logconcave function  $f$ , we have  $\psi_p \gtrsim \sqrt{t}$ .

Proof.

Apply localization lemma.

The resulting statement in 1-d is implied by the following Brascamp-Lieb inequality: the variance of a density given by a Gaussian times a logconcave function (in one dimension) is at most the variance of the Gaussian.

Recall that  $p_t(x) \propto e^{x^T c_t - \frac{t}{2}\|x\|^2} p(x)$  has such a Gaussian part!

Hence,  $\psi_{p_t} = \Omega(\sqrt{t})$ .

---



# How long can we go?

---

$p_t$  is a martingale:  $dp_t(x) = (x - \mu_t)^T dW_t \cdot p_t(x)$

Let  $A_t$  be the covariance of  $p_t$ .

For any measurable subset  $E$ ,

$$\frac{d}{dt} \int_E p_t(x) dx = \int_E (x - \mu_t)^T dW_t p_t(x) dx = \left( \int_E (x - \mu_t) p_t(x) dx \right)^T dW_t$$

Hence, the measure of  $E$  (or any subset) is also a martingale.

$$\begin{aligned} \text{Var} \left( \frac{d}{dt} \int_E p_t(x) dx \right) &= \left\| \int_E (x - \mu_t) p_t(x) dx \right\|_2^2 \\ &\leq \max_{\|\zeta\|_2 \leq 1} \left( \int_E (x - \mu_t)^T \zeta \cdot p_t(x) dx \right)^2 \\ &\leq \max_{\|\zeta\|_2 \leq 1} \int_E ((x - \mu_t)^T \zeta)^2 \cdot p_t(x) dx \\ &= \|A_t\|_{op} \end{aligned}$$

As long as  $\|A_t\|_{op}$  is bounded, any set is approximately preserved.

---



# Suffices to bound $\psi$ for $p_t$

---

$$dp_t(x) = (x - \mu_t)^T dW_t \cdot p_t(x)$$

Thm [E. Milman09] To bound  $\psi$  it suffices to consider subsets of measure  $1/2$ .

Suppose  $\int_0^T \|A_t\|_o dt \leq 0.01$  with constant probability.

Since  $\mathbb{E}p_T = p_0$ , we have that

$$\begin{aligned} \frac{p(\partial S)}{p(S)} &= 2\mathbb{E}p_T(\partial S) \\ &\geq 2\psi_{p_T} \mathbb{E}(\min p_T(S), p_T(S^c)) \\ &\geq \frac{2\psi_{p_T}}{4} \Pr\left(\frac{1}{4} < p_T(S) < \frac{3}{4}\right) \\ &= \Omega(\psi_{p_T}) = \Omega(\sqrt{T}). \end{aligned}$$

Need to keep the spectral norm of covariance small for as long as possible...

---



# Bounding $\|A_t\|_{op}$

---

Back where we started?!

The stochastic process will give us some control.

We use the potential  $\phi_t = \text{Tr}A_t^2$ .

We will see that  $\text{Tr}A_t^2 \leq 2\phi_0$  when  $t \lesssim 1/\sqrt{\phi_0}$  with high probability.

Therefore, if  $T \leq 0.001/\sqrt{\phi_0}$ , we have  $\int_0^T \|A_t\|_{op} dt \leq 0.01$ .

This gives  $\psi_p \gtrsim \sqrt{T} \gtrsim (\text{Tr}A^2)^{-1/4}$ .

- ▶ For isotropic  $p$ , this is  $\psi_p \gtrsim n^{-1/4}$ .



# Bounding the largest eigenvalue

---

$$\phi_t = \text{Tr} A_t^2$$

Let's Itô it!

$$\begin{aligned} d\phi_t &= 2E(x - \mu_t)^T A_t (x - \mu_t) (x - \mu_t)^T dW_t - 2\text{Tr} A_t^3 dt + E((x - \mu_t)^T (y - \mu_t))^3 dt \\ &= \delta_t dt + v_t^T dW_t \end{aligned}$$

**Lemma.** For a logconcave density  $p$ ,

- $\mathbb{E}(\|x\|^k) \leq (2k)^k \mathbb{E}(\|x\|^2)^{k/2}$
- $\mathbb{E} |(x - \mu)^T (y - \mu)|^3 \lesssim (\text{Tr} A^2)^{3/2}$
- $\|\mathbb{E}(x - \mu)(x - \mu)^T A(x - \mu)\| \lesssim \|A\|_{op}^{1/2} \cdot \text{Tr} A^2$

Using this,  $|\delta_t| \lesssim \phi_t^{3/2}$  and  $\|v_t\| \lesssim \phi_t^{5/4}$

Therefore,  $d\phi_t \lesssim \phi_t^{3/2} dt + \phi_t^{5/4} dW_t$

---



# Bounding the largest eigenvalue

---

$$\phi_t = \text{Tr} A_t^2$$

After Itôing,

$$d\phi_t \lesssim \phi_t^{3/2} dt + \phi_t^{5/4} dW_t$$

Or

$$\phi_t \lesssim \phi_t^{3/2} t + \phi_t^{5/4} \sqrt{t}$$

So  $\phi_t \leq 2\phi_0$  for  $t \lesssim 1/\sqrt{\phi_0}$

And for  $T = c/\sqrt{\phi_0}$ , we get:

$$\int_0^T \|A_t\|_o dt \leq \int_0^T \sqrt{\phi_t} dt \leq \sqrt{2\phi_0} \cdot T \leq 0.01.$$

So the measure of the subset stays balanced up to time  $T = c/\sqrt{\phi_0}$  and the lower bound on expansion is

$$\Omega(\sqrt{T}) \gtrsim (\text{Tr} A^2)^{-1/4}$$

(This is  $\gtrsim n^{-1/4}$  for isotropic  $p$ .)

---





# An improved concentration inequality

---

Thm. [Paouris2006]. For an isotropic logconcave  $p$ ,

$$\Pr(\|x\| - \sqrt{n} > c \cdot t) \leq e^{-t}$$

Best possible when  $t \gtrsim \sqrt{n}$ .

[Guédon-E. Milman]: RHS is  $\exp(-\min(\frac{t^3}{n}, t))$ .

[Lee-V. 2017] RHS is  $\exp(-\min(\frac{t^2}{\sqrt{n}}, t))$

Thm. For any isotropic logconcave  $p$ , and any Lipschitz function  $g$ ,

$$\Pr(|g(x) - \bar{g}| > c \cdot t) \leq 2e^{-\frac{t^2}{t+\sqrt{n}}}$$



# Outline

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Original Motivation

Connections

- ▶ Probability and Geometry
- ▶ Algorithms

Techniques

- ▶ Localization
- ▶ Stochastic Localization

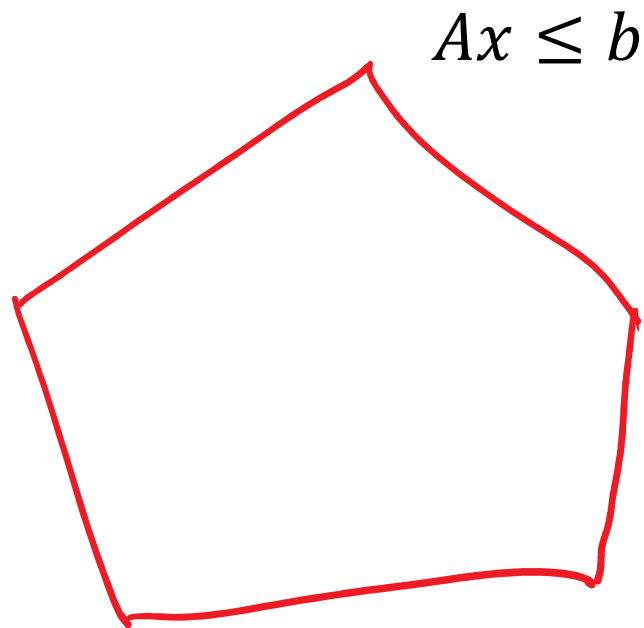
Some open problems



# Deterministic Polytope Volume

---

Can we estimate the volume of an explicit polytope in *deterministic* polynomial time?



## Lower bound for Sampling

---

KLS says complexity of sampling from a warm start is  $n^2$ .

Is this the best possible?



# Faster isotropy and sampling

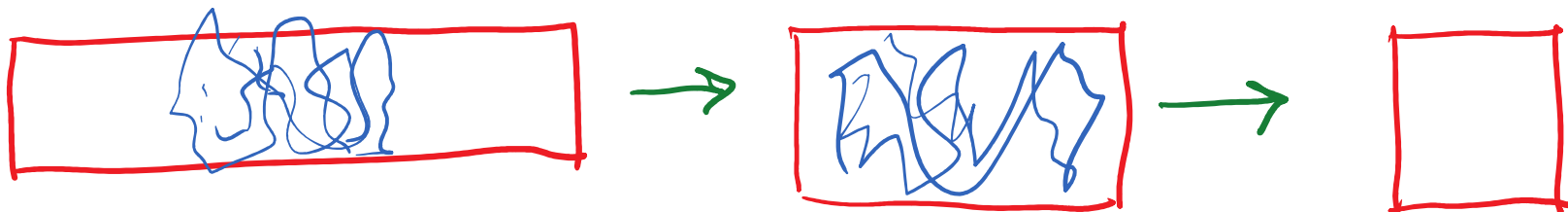
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Isotropic transformation/rounding is the bottleneck for faster general volume computation/sampling.

Candidate algorithm:

Repeat:

1. Estimate the covariance of the standard Gaussian density restricted to the current convex body.
2. If the covariance has eigenvalues smaller than some constant, apply a transformation to make it identity.



Conjecture [Cousins-V.]. This algorithm terminates in  $O(\log n)$  iterations with a well-rounded body.

---



# Faster isotropy and sampling

---

Per-step arithmetic complexity:  $n^2$ .

Coordinate Hit-and-Run. Could be faster by a factor of  $n$  in the per-step complexity.

But is it rapidly mixing?  $n^3$ ?!  

---



# Manifold KLS

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Thm. [Lee-V. 2017] Let  $K$  be a convex body and  $\phi: K \rightarrow \mathbb{R}$  be a convex function with a convex Hessian. Let  $d$  be the distance in the Riemannian metric defined by the Hessian. Then, for any partition of  $K$  into subsets  $S_1, S_2, S_3$ ,

$$\frac{\int_{S_3} e^{-\alpha\phi(x)} dx}{\min \int_{S_1} e^{-\alpha\phi(x)} dx, \int_{S_2} e^{-\alpha\phi(x)} dx} \gtrsim \sqrt{\alpha} d(S_1, S_2)$$

In other words, this Gibbs distribution satisfies a manifold KLS!  
 $\phi(x) = \|x\|^2$  and Euclidean metric  $d$  is the special case of KLS for a Gaussian restricted to a convex body.

What are further generalizations?

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# Needle decompositions

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Used by [Bobkov]; also [Chandrasekaran-Dadush-V.]

Apply hyperplane cuts to get a needle decomposition

Maintain relative measure of subset  $S$ .

Show that a positive fraction of needles have bounded variance.

Conclude KLS!

