

Day 4 Talk 5

Van Handel

"the dimension-free structure of non-homogeneous random matrices."

Simple setting:

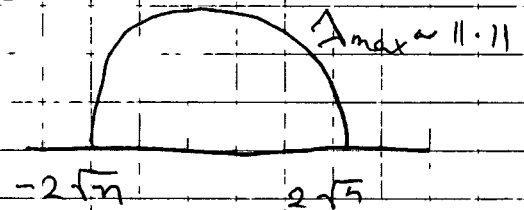
• X $n \times n$ symmetric

• $X_{ij} \sim N(0, b_{ij}^2)$ indep.

Q. What can we say about spectrum?

Classical Wigner-matrix $b_{ij} = 1 \forall i, j$

distribution of eigenvalues



- what if b_{ij} very different?
- maximum / minimum

E.g. When does X ($\infty \times \infty$) random matrix define a bdd operator ℓ_2 ?
(if $b_{ij} = 1$ oper norm $\rightarrow \infty$)

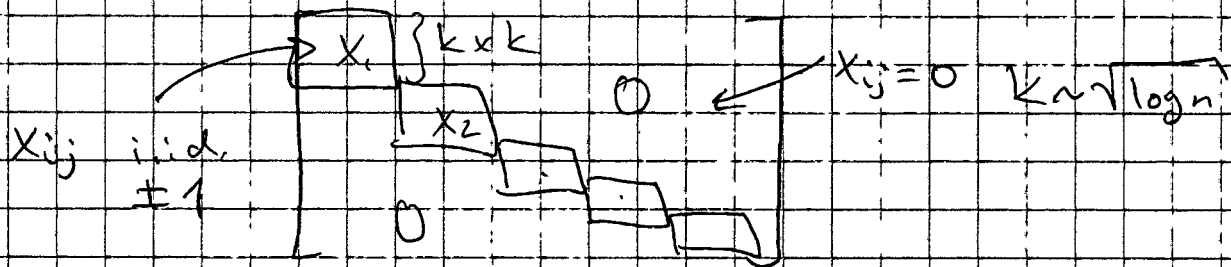
Conjecture (Latafala ~ 15 years ago)

$$\mathbb{E} \|x\| \asymp \mathbb{E} \max_i \sqrt{\sum_j x_{ij}^2}$$

- Would answer previous problems.

- What if x_{ij} not Gaussian

= Can not be true for non-Gaussian random matrices: Counter-example (Seginer)



Now:

$$\bullet \mathbb{E} \max_i \sqrt{\sum_j x_{ij}^2} = \sqrt{k} = (\log n)^{\frac{1}{4}}$$

$$\bullet \mathbb{E} \|x\| = \mathbb{E} \max_i \|x_i\|$$

Wigner

By universality x_i behave like Gaussians.

$$\sim \sqrt{\log n}$$

Q Is it true for Gaussians? why?

Thm (L y. w.) Yes

$$\mathbb{E} \|x\| \asymp \mathbb{E} \max_i \sqrt{\sum_j x_{ij}^2} \asymp \max_i \sqrt{\sum_j b_{ij}^2} + \max_{i,j} b_{ij} \sqrt{\log d}$$

if $\max b_{ij}$ decreasing rearrangement

Similar methods give

In fact can do more

constant indep
of p

$$\mathbb{E} \|x\|_{S_p} \lesssim \mathbb{E} \left[\left(\sum_i \left(\sum_j x_{ij}^2 \right)^{p/2} \right)^{1/p} \right] \approx \text{explicit expression}$$

l_p -norm of eigenvalues

In fact

$$\mathbb{E} \left(\Phi(\|x\|_{l_p(l_2)}) \right) \leq \mathbb{E} \Phi(\|x\|_{S_p})$$

$$\leq \mathbb{E} \Phi(C \|x\|_{l_p(l_2)})$$

$\forall 2 \leq p \leq \infty \quad \forall \text{ inc. convex } \Phi$

Rk.

Focus of this talk: 1st estimate

Some history:

- "Wigner bound" $\mathbb{E} \|x\| \lesssim \max_{ij} b_{ij} \sqrt{n}$ increasing variance to be maximal

- Non-commutative Khintchine (Lust, Piquard, Pisier)

"Matrix concentration"

(also called Matrix Chernoff)

could get

$$\mathbb{E} \|x\| \lesssim \sqrt{\log n} \max_i \sqrt{\sum_j b_{ij}^2}$$

- Ledata:

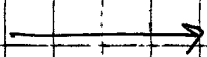
$$\mathbb{E} \|x\| \lesssim \max_i \sqrt{\sum_j b_{ij}^2} + \|(b_{ij})\|_1$$

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Thm. 1 (Bandeira-Rott)

$$\mathbb{E} \|x\| \leq \max_i \sqrt{\sum_j b_{ij}^2} + \max_{ij} b_{ij} \sqrt{\log n}$$

Almost tight



$$\left(\mathbb{E} \|x\|^{\log n} \right)^{\frac{1}{\log n}}$$

Proof of lower bound:

$$\left[\mathbb{E} \|x\|^{\log n} \right]^{\frac{1}{\log n}} \geq \left[\mathbb{E} \|x\|^2 \right]^{\frac{1}{2}}$$

$$\geq \| \mathbb{E} x^2 \|^{1/2} = \max_i \sqrt{\sum_j b_{ij}^2}$$

$$\left[\mathbb{E} \|x\|^{\log n} \right]^{\frac{1}{\log n}} \geq \left[\mathbb{E} |x_{ij}|^{\log n} \right]^{\frac{1}{\log n}} \approx \sqrt{\log n} b_{ij}$$

Moment method:

- $\|x\| = \max_i |x_i|$
- $\tau_r [x^{2p}]^{1/2p} = \left[\sum_i |x_i|^{2p} \right]^{\frac{1}{2p}} \Rightarrow p \sim \log n$
- $\|x\| \leq \tau_r [x^{2p}]^{1/2p} \lesssim n^{\frac{1}{2p}} \|x\|$

$$\left[\mathbb{E} \|x\|^p \right]^{1/p} \sim \left(\mathbb{E} \text{Tr} [X^2]^p \right)^{1/2p} \quad p \sim \log n$$

- ↓
- Compare to a much smaller Wigner matrix
 - + Switch to a combinatorial quantity

- Must use $p \sim \log n$ - moment

Discard ^{this} approach → Try something different!

$$\|x\| = \sup_{v \in S^{n-1}} | \langle v, xv \rangle |$$

- Gaussian variable on the sphere

- Many tools: Majorizing measures?

Thm 2 (RvH '12)

$$\mathbb{E} \|x\| \leq \max_i \sqrt{\sum_j b_{ij}^2} + \max_{ij} b_{ij} \log i$$

Can't get beyond $\log i$ to with this method

Now.

Combine the approaches of the
two prior sub-optimal results.

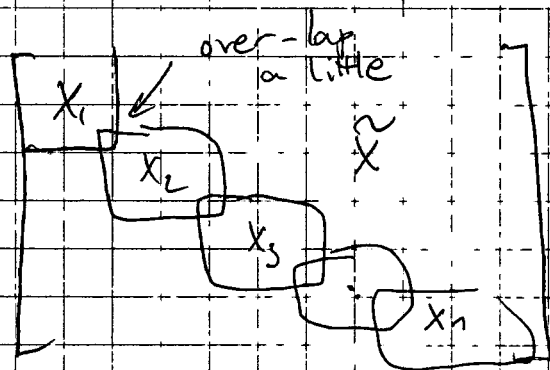
Give names

$$a = \max_i \sqrt{\sum_j b_{ij}^2}$$

$$b = \max_{ij} b_{ij} \sqrt{\log i}$$

To show: $\|x\| \leq a + b$

Lemma every matrix looks like this
(up to reordering rows/columns)



$n \times n$ (n may be ∞)

$$X = U X_i$$

* $\dim(x_i) \sim 2^{2^i}$

* $\max_j b_{ij} \uparrow_{i \in \mathbb{N}} \sim \frac{b}{\sqrt{\log(i+1)}}$ technically...

* $\max_j b_{ij} \uparrow_{i \in \mathbb{N}} \leq \frac{a}{\sqrt{i}}$

"nearly block-diagonal matrix"

Sketch of pf of Thm (using lemma)

In the core x_i :

by Thm 1 (dimension dep bound)

$$\mathbb{E} \|x_i\| \leq \underbrace{\max_i \sqrt{\sum_j b_{ij}^2}}_a + \sqrt{\log(2^i)} \cdot \underbrace{b}_{\sqrt{\log 2^{2^{i-1}}}}$$

Bound that I want!

Combinatorial
methods

$$\mathbb{E} \|x\| \approx \mathbb{E} \max_i \|x_i\|$$

$$\leq \max_i \mathbb{E} \|x_i\| + \sum_i \underbrace{\text{Var}[\|x_i\|]}_b^{1/2}$$

by Gaussian Concentration $\leq \frac{b}{\sqrt{\log 2^{2^{i-1}}}}$

$$\leq a + b + b \leq a + b$$

Outside the core: x

use Thm 2 (dimension free bound)

Gaussian
process
methods

$$\mathbb{E} \|x\| \leq \max_i \sqrt{\sum_j b_{ij}^2} + \max_{ij} \frac{a}{\sqrt{i}} \log i$$

$$\leq a$$

~~QED~~

Need both
methods