

Abstract:

When does a candidate have the approval of a majority? How does the geometry of the political spectrum influence the outcome? When mathematical objects have a social interpretation, the associated results have social applications. We will show how generalizations of Helly's Theorem can be used to understand voting in "agreeable" societies. This talk also features research with undergraduates.

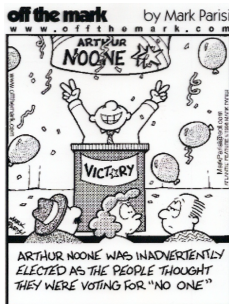
Voting in Agreeable Societies

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Claremont, California, USA

Geometric and Topological Combinatorics
MSRI Connections for Women 2017

thanks: NSF DMS-1002938



Featuring, among others, the work of...



Deborah Berg



Kathryn Nyman



Maria Klawe

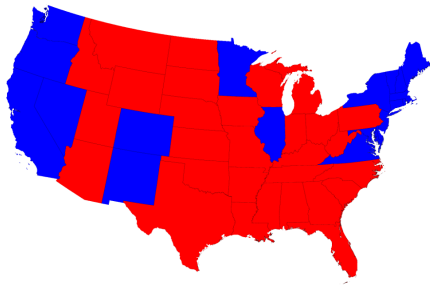
Overview

- ① Sets Modeling Preferences
- ② Intersections
- ③ The Agreeable Society Theorem
- ④ Three Generalizations to Explore

Sets Model Preferences in Some Space

2016 U.S. PRESIDENTIAL ELECTION

Space of states

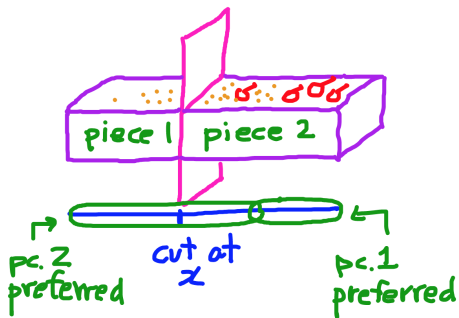


Clinton set, Trump set

Sets Model Preferences in Some Space

CAKE-CUTTING

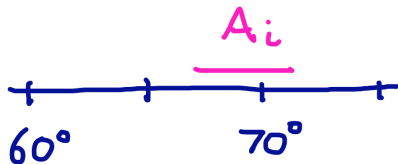
Space of divisions



Sets Model Preferences in Some Space

THERMOSTAT DISAGREEMENTS

Space of temperatures

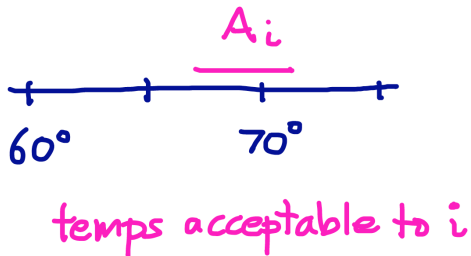


temps acceptable to i

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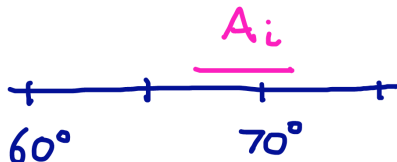


Important: **intersections** of these sets!

Sets Model Preferences in Some Space

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Important: intersections of these sets!

Space has geometry/topology, a notion of “closeness”

Restaurants along Shattuck Ave.

A B C D E F G H I J K L M N

- Where to go for dinner?
How popular is the most popular restaurant?

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Everyone pick five **CONSECUTIVE** restaurants.

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- I predict: one restaurant will get at least $\frac{1}{2}$ the votes!

Third parties in American Politics



graphic: *The Simpsons*

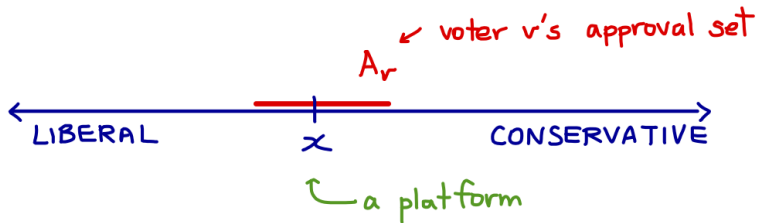
Approval Voting

Voters vote for as many options they “approve”



A Model

A POLITICAL SPECTRUM



Society of Voters



A *society* of voters consists of:

Society of Voters



A *society* of voters consists of:

- a *political spectrum*. Each point a *platform*.

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Q. How popular is the most popular platform? Or candidate?

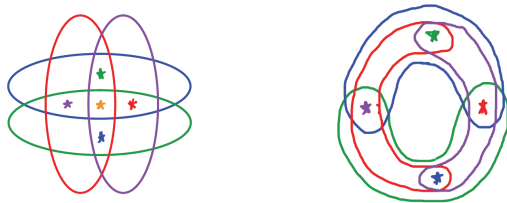
Intersection Theorems

If approval sets are intervals.... they're convex sets.

Intersection Theorems

Helly's Theorem, 1913

In a finite family of **convex** sets in d dimensions, if every $d + 1$ sets intersect mutually, then all sets have a point in common.

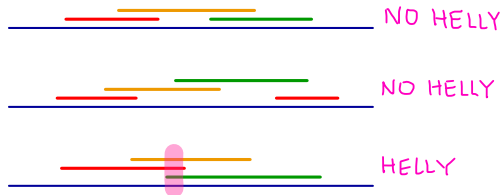


Note: true for infinite collections if sets are also compact.

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Helly-on-a-line

Any **pairwise-intersecting** family of intervals must have a point in all sets.

Setting Thermostats

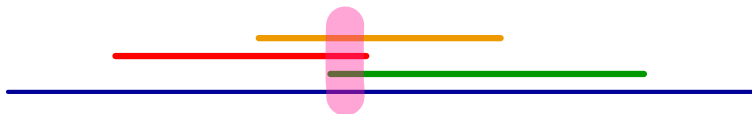
Helly-on-a-line rephrased:

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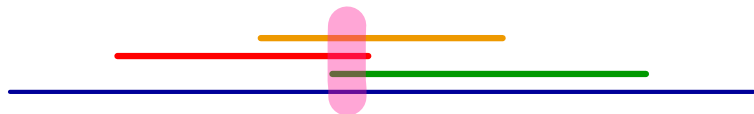


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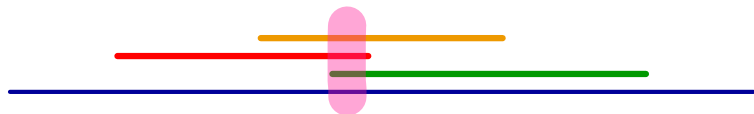


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- Since each pair A_i, A_j overlap, $L_i \leq H_j$.
- So $\max_i L_i \leq \min_j H_j$, any temp in between works. □

Super-agreeable Societies

- A *linear* society: spectrum is a line.



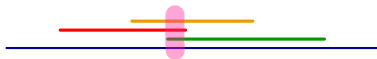
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Hypothesis too strong?

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We can say more...

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$(3, 4)$ -agreeable \implies some platform wins **2/3** the votes.

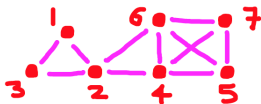
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Proof Sketch: Agreeable Society Theorem

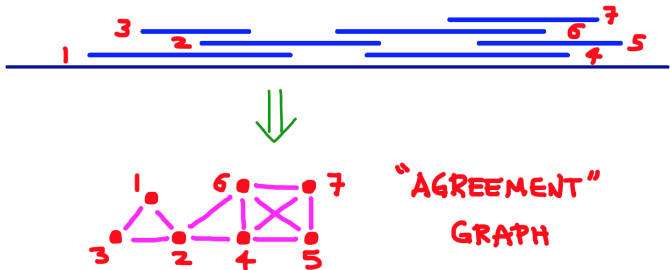
- Construct an **agreement graph** G :



"AGREEMENT"
GRAPH

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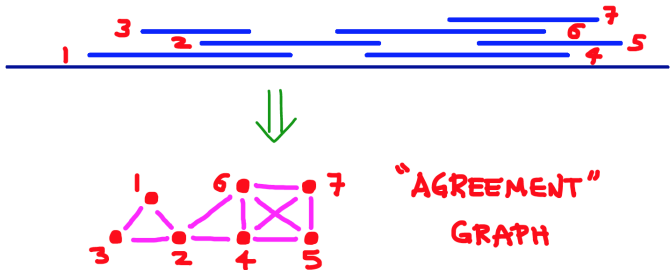
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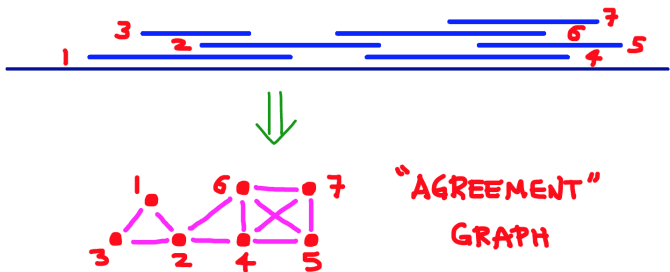
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Proof Sketch: Agreeable Society Theorem

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- **Cliques** in $G \iff$ **intersections** of approval sets (Helly)
- **Color** overlapping voters different colors.
- For linear societies, G is **perfect**: (**chromatic#** = **clique#**)
- (k, m) -condition \implies **color classes** aren't too big, so there have to be **many colors**!



Related

Hadwiger-Debrunner, 1957

Let \mathcal{A} be a collection of compact intervals on the line such that among any m intervals, some k mutually intersect. Then the sets can be *pierced* by a family of $m - k + 1$ points, i.e., every set contains one point in the family.

Related

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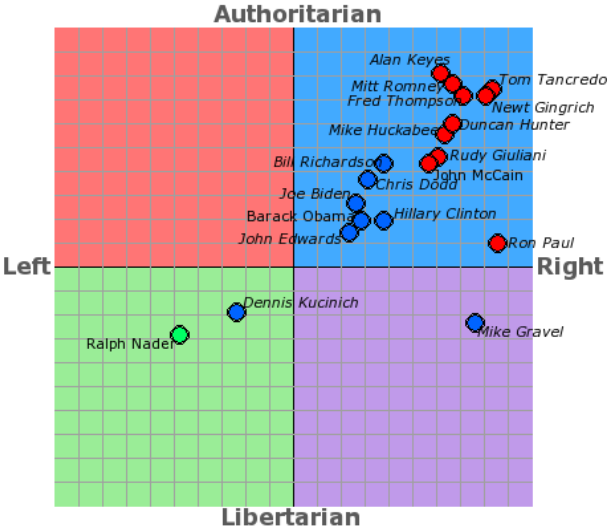
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Note: *piercing* $\#$ = min'l size of group representing all voters.

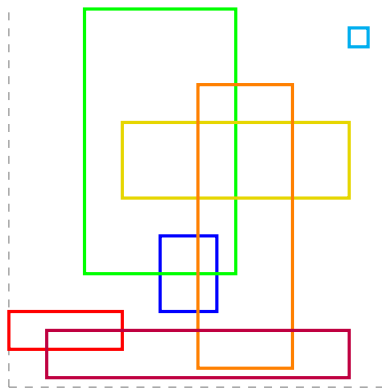
Q. What new questions does this interpretation suggest?

Generalization: change the spaces



Box societies in \mathbb{R}^d

- Political prefs: convex vs. boxes vs. box-convex
- Boxes have Helly property: If boxes pairwise intersect, then there's a point in all boxes



Box societies in \mathbb{R}^d

Theorem. (Berg et. al. 2010)

In an \mathbb{R}^d -convex (k, m) -agreeable society, the agreement proportion β satisfies

$$\beta \geq 1 - \left(1 - \frac{\binom{k}{d+1}}{\binom{m}{d+1}}\right)^{1/(d+1)}.$$

Uses Fractional Helly Theorem.

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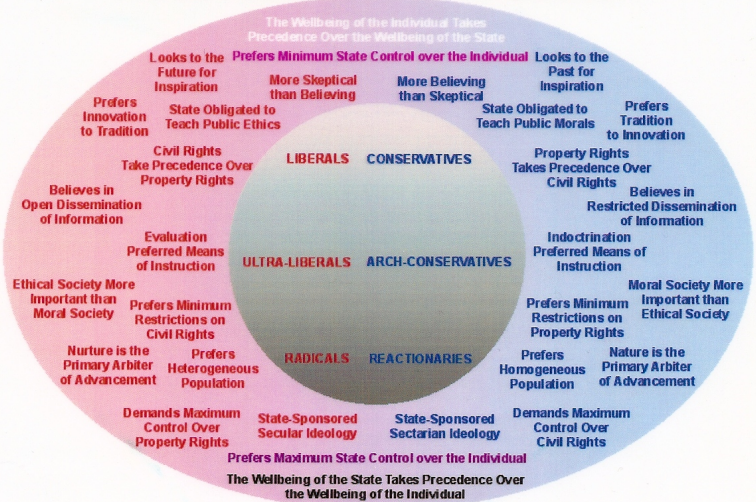
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Open Q: results for *box-convex* sets?

Circular Societies

Most Comfortable with Democratic Forms of Government

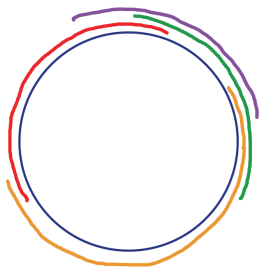


Set Intersections on a Circle

Q. Given a collection of **connected** subsets of circle, if each pair intersect, must there be a point in all the sets?

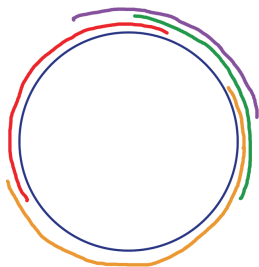
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A. No, but there is a point in a **strict majority of the sets!**

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Circular Society Theorem (Niedermaier-Rizzolo-S. 2014)

In a **circular** society, if any 2 voters agree on a candidate, then some candidate will win **strict majority** approval.

Uses a KKM theorem for trees and cycles.

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Theorem (Hardin 2010)

In a (k, m) -agreeable **circular** society, there is a candidate that has the approval of at least $\frac{k-1}{m}$ of the voters.

Generalization: change the hypotheses

Consider *fractional pairwise agreeability*: suppose α of all voter pairs can agree on some candidate.

What is the **agreement proportion** (of the most popular candidate)?

Pairwise Agreeability

Meta-Theorem

If α of all voter pairs agree,
then the agreement proportion must be at least β .

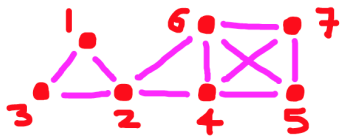
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Think about the agreement graph G .

If there are sufficiently many edges,
then there must be a large clique.



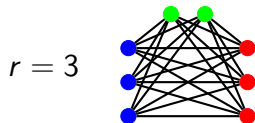
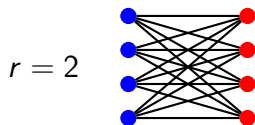
Pairwise Agreeability

Turán's Theorem 1941

Any graph G with n vertices that does not have an $(r + 1)$ -clique as a subgraph has at most $(1 - \frac{1}{r})\frac{n^2}{2}$ edges.

Bound achieved by complete r -partite graph, where size of each partite set varies by at most 1.

Examples for $n = 8$:



Interval Societies - fractional agreement

Theorem (Abbott-Katchalski 1979, re-interpreted)

If an interval graph has pairwise agreement α ,
then the agreement proportion is **at least** β , where

$$\alpha = \beta(2 - \beta).$$

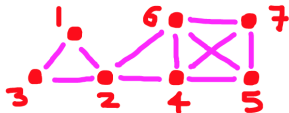
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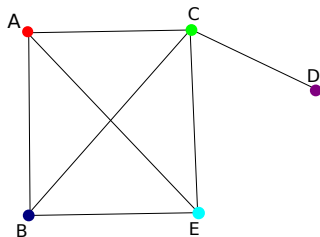
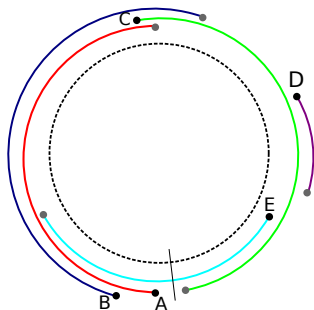
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Ex: So pairwise agreement of $\frac{3}{4} = \frac{1}{2}(2 - \frac{1}{2})$ of all voter pairs ensures an agreement proportion of **at least** $\frac{1}{2}$.



Circular Societies - fractional agreement

Big problem here: clique number k does not imply agreement number k .



Circular Societies - fractional agreement

Theorem (Carlson-Flood-O'Neill-S. 2011)

If an circular arc graph has pairwise agreement α , and the min/max agreement ratio is γ , then the agreement proportion is **at least β** ,

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$$\alpha = \begin{cases} \beta(2 - (1 - \gamma)^2\beta) & 0 \leq \beta \leq \frac{1}{2} \\ \beta(\gamma + 1)(2 - \beta(\gamma + 1)) & \frac{1}{2} \leq \beta \leq \frac{1}{1+\gamma}. \end{cases}$$

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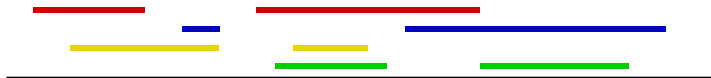
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Ex: In a circular society, if every candidate has at least $\gamma = 20\%$ of the votes of the most popular candidate, and **if there is $\alpha = 46\%$ pairwise agreement**, then **the most popular candidate will have at least $\beta = 25\%$ approval**.

Generalization: Change the set geometry

Double-interval societies

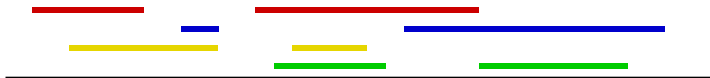
with Kathryn Nyman, Maria Klawe, Jacob Scott:



Generalization: Change the set geometry

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with Kathryn Nyman, Maria Klawe, Jacob Scott:



Theorem. (Klawe-Nyman-Scott-S. 2014)

The approval ratio of any n -voter **pairwise-intersecting, double-interval society** is at least

$$0.268 + \frac{0.789}{n} - \frac{1.732}{24n^2}.$$

Conjecture: tight bound is $1/3$.

Double-N Strings: Smallest Agreement Proportion?

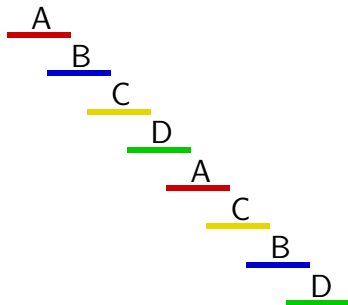


Figure: The double-4 string: ABCDACBD.
Has $\text{diam}=1$, since any 2 letters are 1 apart.

Double-N Strings: Smallest Agreement Proportion?

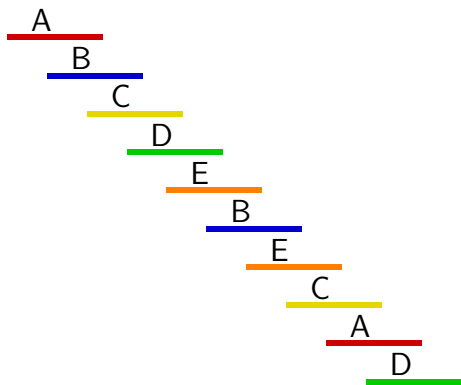


Figure: Double 5-string: ABCDEBECAD has $\text{diam}=2$.

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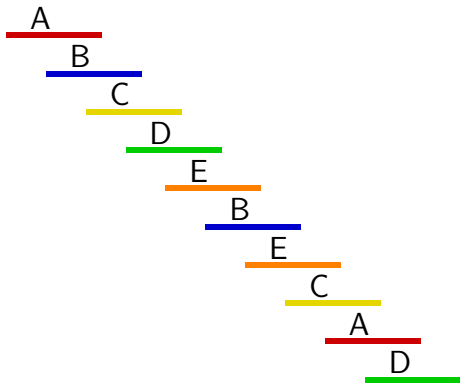


Figure: Double 5-string: ABCDEBECAD has $\text{diam}=2$.

Surprise: asymptotic agreement prop. for double n -strings
 $\geq 8/23 > 1/3$.

Double Interval Societies

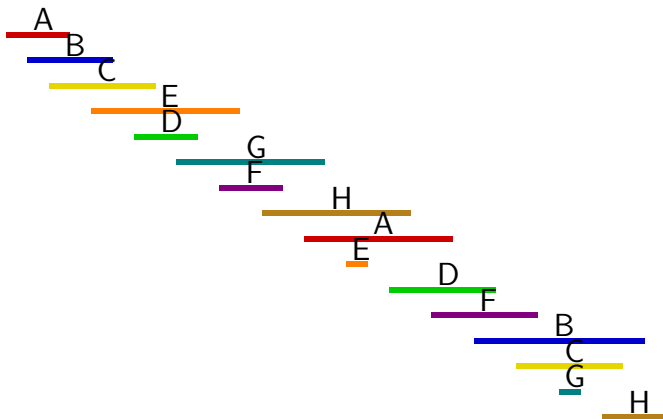


Figure: A society of size 8 with approval number 3.
Derived from double 8-string ABCDEFGHEADFCGBH.

Tantalizing

DOUBLE-INTERVAL SOCIETIES

$a(S)$	n	Approval Ratio	Observed n	Observed Approval Ratio
2	≤ 4	≥ 0.500	4	0.500
3	≤ 8	≥ 0.375	8	0.375
4	≤ 12	≥ 0.333	12	0.333
5	≤ 15	≥ 0.333	15	0.333
6	≤ 19	≥ 0.316	18	0.333
7	≤ 23	≥ 0.304	21	0.333
8	≤ 26	≥ 0.308	24	0.333
9	≤ 30	≥ 0.300	27	0.333
10	≤ 34	≥ 0.294	30	0.333
11	≤ 38	≥ 0.289	33	0.333
12	≤ 41	≥ 0.293	36	0.333

TABLE 1. On the left, this table shows for a given approval number the largest n that is given by inequality (4.3) as well as the resulting bound on the approval ratio derived from inequality (4.1). On the right, this table shows, for a given approval number, known examples of the largest n that has this approval number and the observed approval ratio in that case, obtained by a modification of a double- n string construction.

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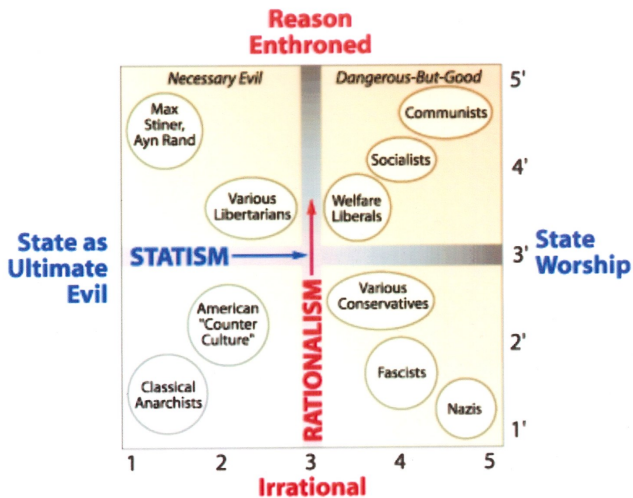
Questions: For double- N strings, is there systematic way to construct strings of smallest diameter?

Beyond double- N strings, is there general construction yielding societies with lowest agreement ratios?

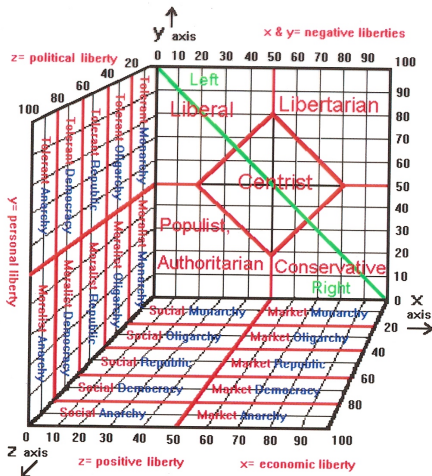
Further questions...

- Agreement proportion for other spaces? How does topology affect agreement?
- Non-convex approval sets?
- Probabilistic voting results?
- Statistics: estimates from samples?
- Other “social” applications?

Two Dimensional Space of Platforms



Three Dimensional Space of Platforms



29-Dimensional Approval Sets?

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29 Dimensions® of Compatibility

Most people know that the key to success in a long-term relationship is compatibility. But what does that mean? If you and your new mate both like foreign movies and mocha ice cream, will you still feel the magic in 25 years?

eHarmony matches singles based on a deeper level of compatibility, not likes and dislikes, but true compatibility. Do you and your potential mate resolve conflict in a similar fashion? Are you both romantics at heart? And we are the only online dating web site that matches singles based on these 29 Dimensions®. To help you better understand these dimensions, we've grouped them into Core Traits and Vital Attributes.

Core Traits are defining aspects of who you are that remain largely unchanged throughout your adult life. Vital Attributes are based on learning experience, and are more likely to change based on life events and decisions you make as an adult.



Take Aways

- Math can model/answer questions in the social sciences
- Social problems can motivate new mathematical questions

