

Mogami triangulations

Bruno Benedetti

University of Miami

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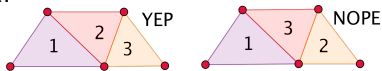
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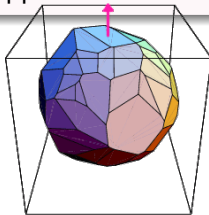
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Fact (Bruggesser–Mani 1971, also Schläfli 1850)

The boundaries of simplicial polytopes are shellable: “Lift off” from a facet, and moving along a generic line, record the facets in the order in which they appear at the horizon (“rocket shelling”).



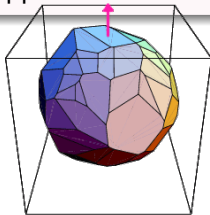
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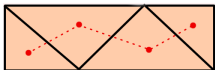


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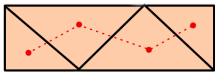
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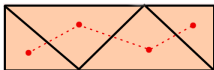
Trees of d -simplices



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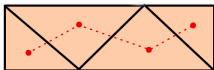
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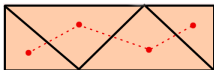


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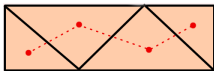
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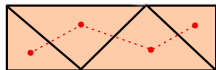




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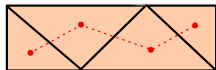


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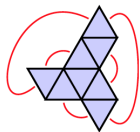
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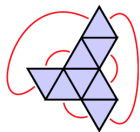


Figure: Gluing boundary edges according to red matching \rightsquigarrow octahedron

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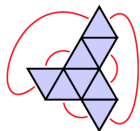


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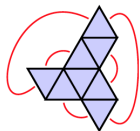


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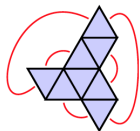


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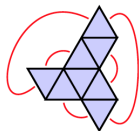


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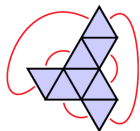


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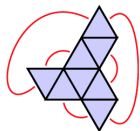


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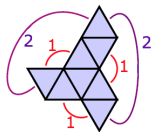
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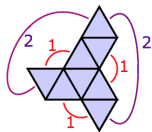
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From now on, $d \geq 2$.



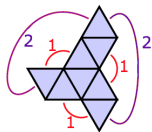
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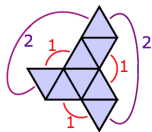
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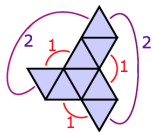
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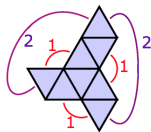
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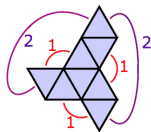
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In dimension $d \leq 3$, also Mogami d -manifolds are exp. many.

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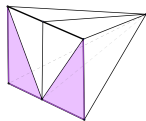
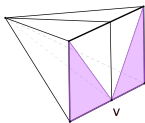
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So LC/Mogami triangulations = combinatorial way to capture simple connectedness (for manifolds).

Connection to shellability





Gluing Lemma

Let A, B be two d -pseudomanifolds. Let $C \subset \partial A$ be a pure $(d - 1)$ -dimensional complex combinatorially equivalent to a subcomplex $C' \subset \partial B$. Let $A \cup B$ be the complex obtained from the disjoint union $A \sqcup B$ by identifying $C \equiv C'$.



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Question (Gromov et al)

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If any is true, then 3-balls are exponentially many!

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... unfortunately, this characterization does not extend to Mogami. So Mogami's conjecture stayed open.

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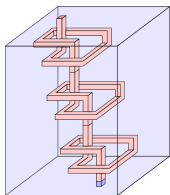
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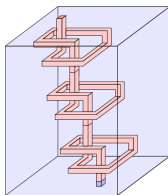
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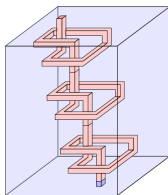


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Any (smooth) knot can be realized in some triangulated 3-ball B , as one interior edge $[x, y]$ plus a boundary path from x to y .

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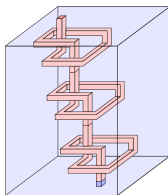


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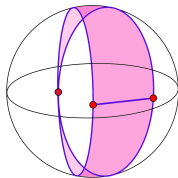
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↪ New idea: The Mogami construction of a 3-ball without interior vertices, could be spartan... 'Chic' gluings may cost interior vertices!

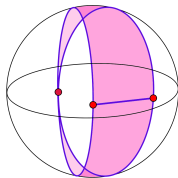
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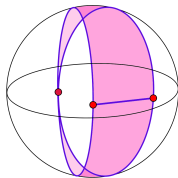
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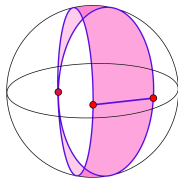
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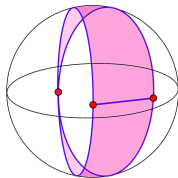


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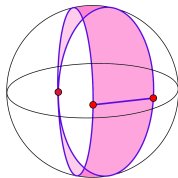
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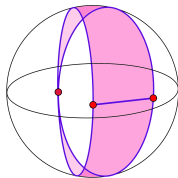
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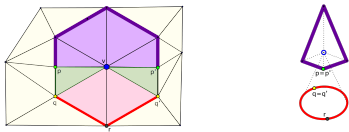
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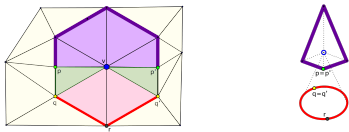
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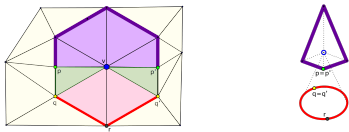
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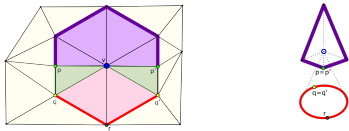
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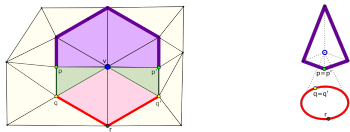
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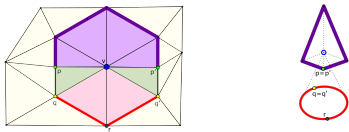
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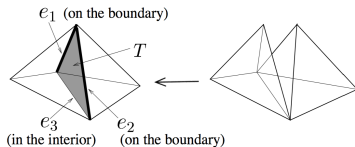
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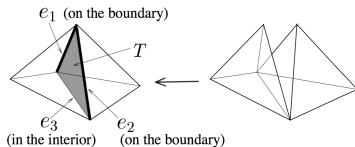
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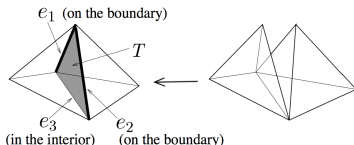
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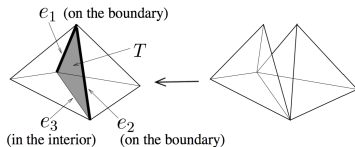
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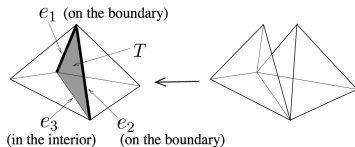
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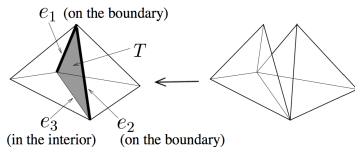
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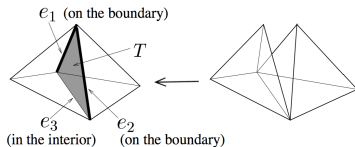
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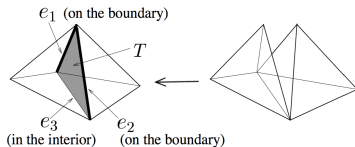
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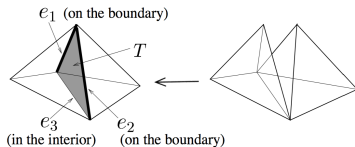
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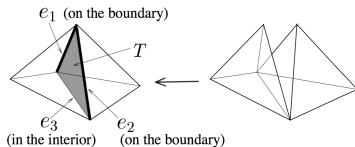
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So the cone over an annulus, say, is Mogami but not LC. (Annulus is not simply connected, so not Mogami, so not LC.)

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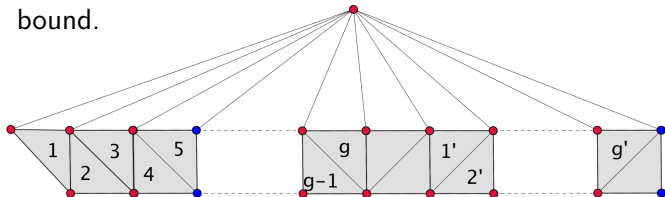
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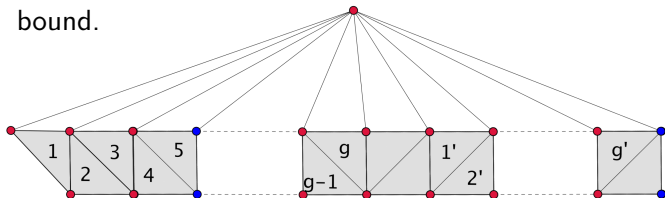
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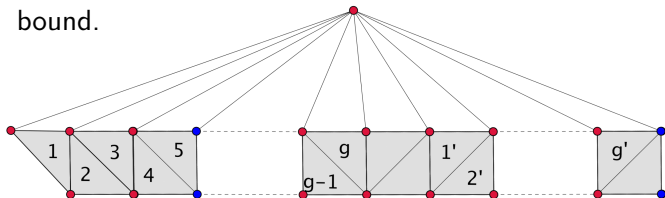
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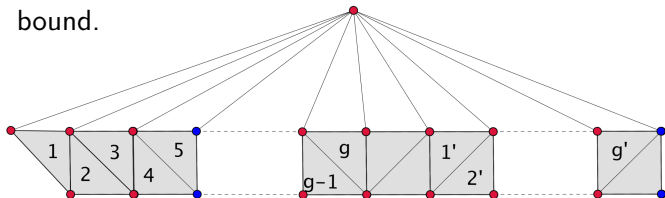
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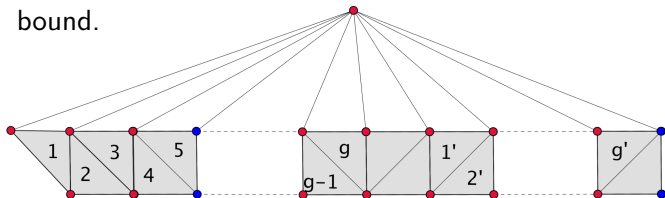
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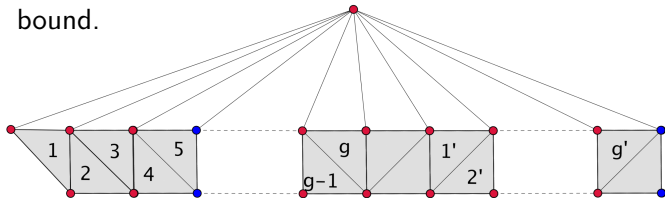
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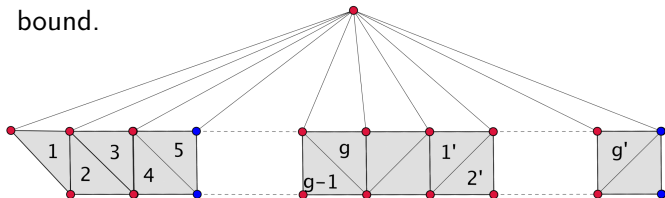
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