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Sam Raskin Speaker's Name:				
Talk Title:	The notion of singul	ar support in	DAG and its	applications I
Date:/_4	<u>19</u>	4 : 00 am / (	m circle one)	
Please summarize the lecture in 5 or fewer sentences:				

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## COHERENT SINGULAR SUPPORT

#### SAM RASKIN

#### 1. Meta-overview of mathematical research

- (1) The first step is to have something you want to be true.
- (2) The second is to calculate something, and usually you realize that your initial dream has complications.
- (3) The third step is to salvage what you can, which is where technical stuff happens.

The subject of coherent singular support is technical in nature, concentrated in step (3). But I want to start with (1).

### 2. What we want

Let k be a field of characteristic 0. We recall the dual numbers  $k[\epsilon] := k[\epsilon]/\epsilon^2$ .

**Lemma 2.1.** There is an equivalence between flat  $k[\epsilon]$ -modules and extensions

$$0 \to V \to \mathcal{E} \to V \to 0$$

where V is a vector space over k.

*Proof.* Given  $0 \to V \to \mathcal{E} \to V \to 0$  you get a  $k[\epsilon]$ -module  $\mathcal{E}$  where multiplication by  $\epsilon$  is the composition  $E \to V \hookrightarrow E$ .

Conversely, if M is a flat  $k[\epsilon]$ -module we get an extension

$$0 \to V = M/\epsilon \xrightarrow{\epsilon} M \to V \to 0.$$

We can try to extend this by working in derived categories, and dropping the word "flat".

The hope is then to get an equivalence of the form

$$k[\epsilon] - \text{mod} \cong \{V \in \text{Vec}, \eta \colon V \to V[1]\}.$$

Here everything is occurring in some suitable derived category.

However this is wrong, and for somewhat subtle reasons. In order to explain why, we have to give a digression about how to do these sorts of calculations. (It's possible to make a mistake and think that you proved this equivalence.)

Date: February 4, 2019.

#### SAM RASKIN

#### 3. Yoga of derived categories

Let C be a dg category. To first approximation, this means a category enriched over chain complexes. But we really want to view C as an object in the infinity category of dg categories, which means that for all practical purposes we cannot distinguish quasi-isomorphisms from chain complexes. (By contrast, when we look at chain complexes, it makes sense to make this distinction.)

Assume that  $\mathcal{C}$  has all (homotopy) colimits, (equivalently, all direct sums).

**Definition 3.1.** An object  $\mathcal{F} \in \mathcal{C}$  is *compact* if  $\text{Hom}(\mathcal{F}, -) \colon \mathcal{C} \to \text{Vect commutes}$  with all colimits (equivalently, all direct sums). (Here Vect is the derived category of vector spaces, i.e. chain complexes.)

**Example 3.2.** Think of compactness as a "smallness" condition, analogous to "finitely presented".

**Example 3.3.** Let A be a ring. Then  $A \in A - \text{mod}$  is compact.

**Definition 3.4.** A category C is *compactly generated* if  $\operatorname{Hom}_{\mathcal{C}}(\mathcal{G}, \mathcal{F}) = 0$  for all compact  $\mathcal{G} \in \mathcal{C}$  implies that  $\mathcal{F} = 0$ .

In this case, if  $\mathcal{C}^c \subset C$  is the subcategory of compact objects, then you can recover  $\mathcal{C}$  canonically via the "ind-category" construction:  $\mathcal{C} = \text{Ind}(\mathcal{C}^c)$ .

**Example 3.5.**  $(A - \text{mod})^c = \text{Perf}(A)$ , the smallest subcategory of A-mod containing A and closed under shifts, finite colimits, and direct summands.

**Definition 3.6.**  $\mathcal{F} \in \mathcal{C}$  is a compact generator if  $\mathcal{F} \in \mathcal{C}^c$  and  $\underline{\operatorname{Hom}}_{\mathcal{C}}(\mathcal{F}, \mathcal{G}) = 0 \Longrightarrow \mathcal{G} = 0$ .

**Example 3.7.**  $A \in A - \text{mod}$  is a compact generator.

Conversely, if  $\mathcal{F} \in \mathcal{C}$  is a compact generator then  $\mathcal{C} \cong A - \text{mod where } A = \underline{\text{End}}_{\mathcal{C}}(\mathcal{F})$ , via the functor  $\underline{\text{Hom}}(\mathcal{F}, -)$ .

## 4. The hope, revisited

We now reformulate our hope: let B the tensor algebra T(k[-1]), i.e. the free dg algebra on k[-1]. Then our hope is that

$$k[\epsilon] - \mod \cong B - \mod.$$

To prove this, it would be enough to find a compact generator inside  $k[\epsilon] - \text{mod}$  and then show that its endomorphisms are B.

Maybe the first thing to try is  $k[\epsilon]$ , but this will just give the tautological identification  $k[\epsilon] - \mod \cong k[\epsilon] - \mod$ .

The interesting thing to try is  $\mathcal{F} = k$ , where  $\epsilon$  acts by 0. A quick calculation shows that  $\underline{\operatorname{End}}_{k[\epsilon]}(k) = B$ . To compute this, use your favorite resolution of k as a  $k[\epsilon]$ -module:

 $\dots \to k[\epsilon] \xrightarrow{\epsilon} k[\epsilon] \xrightarrow{\epsilon} k[\epsilon] \to 0$ 

The  $\underline{\operatorname{End}}_{k[\epsilon]}(k)$  will be a *B*-module, so any element gives a map  $B \to \underline{\operatorname{End}}_{k[\epsilon]}(k)$ , and then you need to check that this induces an isomorphism on cohomology.

**Exercise 4.1.** Check that this corresponds to what we did in the beginning.

However there is a problem:  $k[\epsilon]$  is not compact in  $k[\epsilon] - \text{mod.}$ 

*Proof.* Define  $\mathcal{F}_n$  to be the naive truncation

$$\mathcal{F}_n = 0 \to k[\epsilon] \xrightarrow{\epsilon} \dots \xrightarrow{\epsilon} k[\epsilon] \to 0$$

Then  $k = \underline{\operatorname{colim}}_n \mathcal{F}_n$ , but if k were compact it would be a summand of some  $\mathcal{F}_n$ , and this would contradict the computation of its self-Ext. In other words, any perfect complex has the property that Hom out of it vanishes in sufficiently high homological degree.

#### 5. Rescuing the hope

Now we have to do something technical.

We define  $\operatorname{Coh}(k[\epsilon]) \subset k[\epsilon] - \mod$  to be the full subcategory of bounded complexes with finite-dimensional cohomology. This contains  $\operatorname{Perf}(k[\epsilon])$  strictly, since for example it contains k.

We now claim that

$$\operatorname{Ind}(\operatorname{Coh}(k[\epsilon])) \cong B - \operatorname{mod}$$

*Proof.* Apply the previous argument (using that  $\operatorname{End}_{\operatorname{Coh}}(k) = \operatorname{End}_{k[\epsilon]-\operatorname{mod}}(k) = B$ ), using that k is compact in  $\operatorname{Ind}\operatorname{Coh}(k[\epsilon])$  by fiat, and it generates.  $\Box$ 

We have an embedding

$$\operatorname{Perf}(k[\epsilon]) \hookrightarrow \operatorname{Coh}(k[\epsilon])$$

which induces a fully faithful functor (on applying Ind)

$$k[\epsilon] - \operatorname{mod} \hookrightarrow \operatorname{IndCoh}(k[\epsilon]) \cong B - \operatorname{mod}.$$

**Lemma 5.1.**  $k[\epsilon] - \mod$  corresponds to the full subcategory  $(B - \mod)_{loc.\ nilp.}$  of  $\mathcal{G} \in B - \mod$  where

$$\underbrace{\operatorname{colim}}_{\mathcal{G}}(\mathcal{G} \to \mathcal{G}[1] \to \mathcal{G}[2] \to \ldots) = 0.$$

*Proof.* The functor  $k[\epsilon] \to k$  induces the 0 map  $k \to k[1]$ . Since  $k[\epsilon] - \text{mod}$  is closed under colimits, everything lies in this subcategory.

**Example 5.2.** There are two versions of "k" in  $\operatorname{IndCoh}(k[\epsilon])$ . One is gotten from  $k \in \operatorname{Coh}(k[\epsilon])$ , an the other is the colimit of the  $\mathcal{F}_n$ . They are different!

Moral: perfect complexes correspond to some kind of "local nilpotency" condition.

**Remark 5.3.** There is an alternative path we could have taken. If you take the category of flat  $k[\epsilon]$ -modules as a dg category, it is equivalent to  $\text{IndCoh}(k[\epsilon])$ .

Generalization of this example: let Y be smooth over k. Let  $f: Y \to \mathbf{A}^1$ . Consider  $X := \{f = 0\} = Y \times_{\mathbf{A}^1} 0$ . (This is a derived scheme if f is not flat.) Let  $\mathcal{F} \in \operatorname{QCoh}(X)$ . There exists a canonical map  $\mathcal{F} \to \mathcal{F}[2]$  naturally in  $\mathcal{F}$ , constructed as follows.

Notation: if  $g \colon S \to T$  is a map of (suitably finite) k-schemes, there is a pullback functor

$$g^* \colon \operatorname{QCoh}(T) \to \operatorname{QCoh}(S)$$

and

$$g_* \colon \operatorname{QCoh}(S) \to \operatorname{QCoh}(T).$$

We have a triangle

$$\mathcal{F}[1] \to i^* i_* \mathcal{F} \xrightarrow{\lambda} \mathcal{F}.$$

(From the Koszul complex one can at least see that this is consistent with the size of  $i^*i_*\mathcal{F}$ .) The usual yoga then gives a map  $\mathcal{F} \to \mathcal{F}[2]$ .

We define  $\operatorname{Coh}(X) \subset \operatorname{QCoh}(X)$  to be the full subcategory of bounded complexes with (locally) finitely generated cohomology groups. If X is a dg scheme with  $\mathcal{O}_X$ bounded below (not always satisfied but true in our situation), then  $\operatorname{Perf}(X) \subset$  $\operatorname{Coh}(X)$ . This then induces an embedding  $\operatorname{QCoh}(X) \subset \operatorname{IndCoh}(X)$ .

For formal reasons,  $\eta$  extends to IndCoh. For formal reasons,  $\eta$  extends to IndCoh, giving  $\mathcal{F} \to \mathcal{F}[2]$ .

**Proposition 5.4.** We can identify  $\operatorname{QCoh}(X) \subset \operatorname{IndCoh}(X)$  as the full subcategory of  $\mathcal{F}$  where  $\eta$  acts nilpotently, i.e.  $\{\mathcal{F} \mid \operatorname{colim} \mathcal{F}[2n] = 0\}$ .

(In the example,  $\eta$  can be thought of as the obstruction to extending  $\mathcal{F}$  to a first-order deformation.)

*Proof.* Let's show that if  $\eta$  acts nilpotently, then  $\mathcal{F}$  comes from  $\operatorname{QCoh}(X)$ .

Claim: for all  $\mathcal{F} \in \text{IndCoh}(X)$ , ker  $\eta \in \text{QCoh}(X)$ .

Proof: it suffices to study  $\mathcal{F} \in \operatorname{Coh}(X)$ . In this case ker  $\eta = i^*i_*\mathcal{F}$ , which is always in  $\operatorname{Perf}(X)$ . This is because  $i_*\mathcal{F} \in \operatorname{Coh}(Y) = \operatorname{Perf}(Y)$  by Serre's theorem (since Y is smooth over k). By induction we get that ker $(\eta^n) = 0$  for all n > 0. Then  $\mathcal{F} \cong \operatorname{ker}(\eta^{\infty} \colon \mathcal{F} \to \operatorname{Colim} \mathcal{F}[2n]) \in \operatorname{QCoh}$ .  $\Box$